



## Some results related to consensus formation in graphon dynamics

(in collaboration with N. Pouradier Duteil and M. Sigalotti)

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**Conference “Round Meanfield: Crowd-opinion-cell”**  
**Ypatia Laboratory of Mathematical Sciences**

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# Outline of the talk

Multi-agent systems – From microscopic to macroscopic models

Review of consensus methods for microscopic cooperative systems

Consensus analysis in the context of graphon dynamics

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# Multi-agent systems – *Finite dimensional setting*

**Multi-agents** dynamics can be described by **systems of ODEs**

$$\dot{x}_i(t) = v_i(t, \mathbf{x}(t), x_i(t)), \quad i \in \{1, \dots, N\},$$

where

- ◇  $\mathbf{x} = (x_1, \dots, x_N) \in (\mathbb{R}^d)^N$  encodes the **states** of the agents,
- ◇  $v_i : [0, T] \times (\mathbb{R}^d)^N \times \mathbb{R}^d \rightarrow \mathbb{R}^d$  are **non-local velocity fields**.

Breadcrumb trail example (Time-dependent cooperative dynamics)

$$v_i(t, \mathbf{x}, x_i) = \frac{1}{N} \sum_{j=1}^N a_{ij}(t) \phi(|x_i - x_j|)(x_j - x_i).$$

Central observation (pattern formation)

Simple **microscopic** interactions  $\rightsquigarrow$  rich **macroscopic** structures.

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# Multi-agent systems – *Formation of global patterns*

Example (Classical patterns arising in multi-agent systems)

- ◇ **Consensus** (everybody goes at the same place)
- ◇ **Flocking** (everybody goes in the same direction)
- ◇ **Synchronisation** (periodic motions arise in the system)



Macroscopic approximations (Main motivations)

- ◇ Interest for **global** patterns, i.e. that involve **many** agents,
- ◇  $N$  is usually **very large**  $\rightsquigarrow$  **numerical** issues

**Today:** Consensus for **micro** and **macro** cooperative systems.

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# Multi-agent systems – General cooperative dynamics

We consider the **cooperative** dynamics

$$\dot{\boldsymbol{x}}_i(t) = \frac{1}{N} \sum_{j=1}^N a_{ij}(t) \phi(|\boldsymbol{x}_i(t) - \boldsymbol{x}_j(t)|) (\boldsymbol{x}_j(t) - \boldsymbol{x}_i(t)),$$

where

- ◇  $\phi \in \text{Lip}(\mathbb{R}_+, \mathbb{R}_+^*)$  encodes **distance-based** interactions,
- ◇  $a_{ij}(\cdot) \in L^\infty(\mathbb{R}_+, [0, 1])$  represent **communication** links.

**Definition (Graph-Laplacian operators)**

The **graph-Laplacian**  $L_N(t, \boldsymbol{x}) : (\mathbb{R}^d)^N \rightarrow (\mathbb{R}^d)^N$  is defined by

$$L_N(t, \boldsymbol{x}) \boldsymbol{y} = \left( \frac{1}{N} \sum_{i=1}^N a_{ij}(t) \phi(|\boldsymbol{x}_i - \boldsymbol{x}_j|) (\boldsymbol{y}_j - \boldsymbol{y}_i) \right)_{1 \leq i \leq N}.$$

↔ Semilinear reformulation of the dynamics

$$\dot{\boldsymbol{x}}(t) = -L_N(t, \boldsymbol{x}(t)) \boldsymbol{x}(t). \quad (\text{CS})$$

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# Multi-agent systems – Reformulation as graphons (1)

**Idea:** Study consensus for the **mean-field approximation**

$$\partial_t \mu_N(t) + \operatorname{div}_x \left( \Phi(t) \star \mu_N(t) \mu_N(t) \right) = 0$$

[Ha&Liu'09, Carrillo,Fornasier,Rosado&Toscani'10, Piccoli,Rossi&Trélat'15].

**Problem:** Mean-field needs **indistinguishability**, i.e.  $a_{ij}(t) = 1$ .

**Definition (Graph limit)**[LS'07,M'14]

Given a solution  $x(\cdot)$  of (CS), define the **piecewise constant** maps

$$i \in I \mapsto x_N(t, i) := \sum_{k=1}^N x_k(t) \mathbb{1}_{\left[\frac{k-1}{N}, \frac{k}{N}\right)}(i)$$

and

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and denote by  $I := [0, 1]$  the (continuum of) **indices**.



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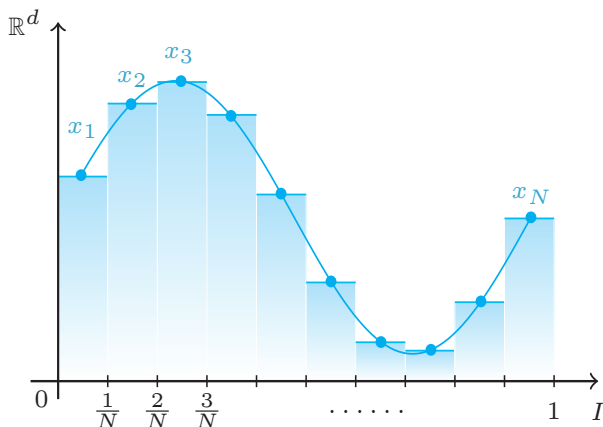
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## Multi-agent systems – Reformulation as graphons (2)

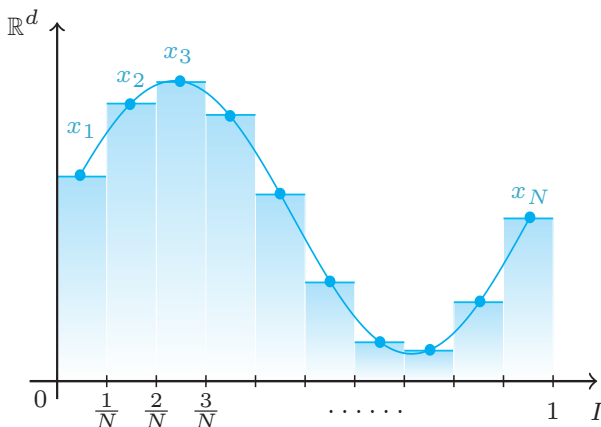


Graphon reformulation of (CS)  $\rightsquigarrow$  infinite dimensional ODEs

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## Consensus – Definition and main ideas

### Definition (Asymptotic consensus formation)

A solution  $\mathbf{x}(\cdot)$  of (CS) converges to **consensus** if

$$\lim_{t \rightarrow +\infty} |x_i(t) - x^\infty| = 0,$$

for all  $i \in \{1, \dots, N\}$  and some  $x^\infty \in \mathbb{R}^d$ .

**Idea:** Quantitative convergence results  $\rightsquigarrow$  **Lyapunov** methods!

### Definition (Candidate energy functionals)

We define the **variance functional**

$$\mathcal{V}(\mathbf{x}) := \frac{1}{2N} \sum_{i=1}^N |x_i - \bar{x}|^2 \quad (\ell_2\text{-convergence}),$$

and the **diameter**

$$\mathcal{D}(\mathbf{x}) := \max_{i,j \in \{1, \dots, N\}} |x_i - x_j| \quad (\ell_\infty\text{-convergence}),$$

$\Leftrightarrow$  Decay prescribed by two **intrinsic scalar** quantities.

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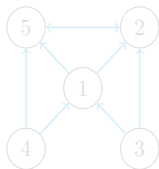
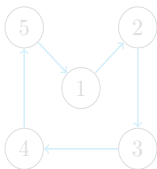
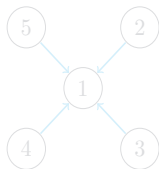
# Consensus – Scrambling coefficient and diameter estimates

Definition (Scrambling coefficient)[Seneta'79]

The **scrambling** of a graph  $\mathbf{A}_N = (a_{ij})_{i,j=1}^N$  satisfying  $a_{ii} = 1$  is

$$\eta(\mathbf{A}_N) := \min_{1 \leq i, j \leq N} \frac{1}{N} \left( \sum_{k=1, k \neq i, j}^N \min \{a_{ik}, a_{jk}\} + a_{ij} + a_{ji} \right)$$

$\hookrightarrow$  Positive if each  $(i, j)$  either **interact** or **follow** the same  $k$ .



Theorem (Quantitative diameter decay)[Motsch&Tadmor'14]

For each  $\mathbf{x}^0 \in (\mathbb{R}^d)^N$ , there exists  $\phi_0 > 0$  such that

$$\mathcal{D}(\mathbf{x}(t)) \leq \mathcal{D}(\mathbf{x}^0) \exp \left( -\phi_0 \int_0^t \eta(\mathbf{A}_N(s)) ds \right).$$

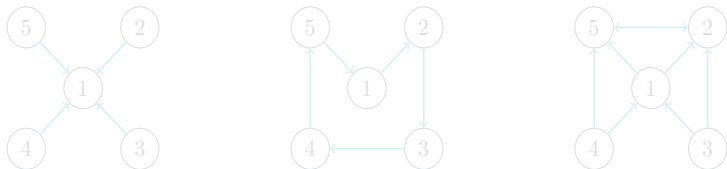
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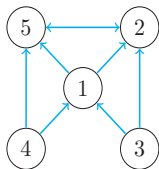
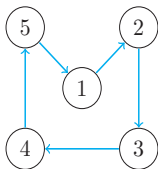
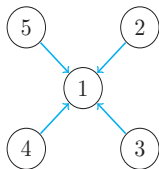
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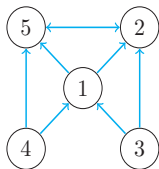
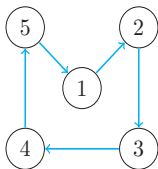
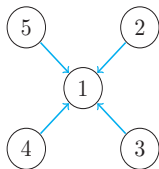
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## Consensus – Fiedler number and variance estimates

Definition (Algebraic connectivity of a graph)[Fiedler'73, Mohar'91]

The **Fiedler number** of a symmetric graph  $\mathbf{A}_N = (a_{ij})_{i=1}^N$  is

$$\lambda_2(\mathbf{A}_N) = \inf_{\mathbf{x} \in \mathcal{C}_N^\perp} \frac{\langle \mathbf{L}_N \mathbf{x}, \mathbf{x} \rangle_N}{|\mathbf{x}|_N^2},$$

where  $(\mathbf{L}_N \mathbf{x})_i = \frac{1}{N} \sum_{j=1}^N a_{ij} (x_i - x_j)$ ,  $\langle \mathbf{x}, \mathbf{y} \rangle_N := \frac{1}{N} \sum_{i=1}^N \langle x_i, y_i \rangle$  and

$$\mathcal{C}_N := \left\{ \mathbf{x} \in (\mathbb{R}^d)^N \text{ s.t. } x_1 = \cdots = x_N \right\}$$

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Suppose that  $\mathbf{A}_N(t) = (a_{ij}(t))_{i,j=1}^N$  is symmetric for a.e.  $t \in \mathbb{R}_+$ . Then for each  $\mathbf{x}^0 \in (\mathbb{R}^d)^N$ , there exists  $\phi_0 > 0$  such that

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$$\begin{aligned}\frac{d}{dt} \mathcal{V}(\mathbf{x}(t)) &= \frac{1}{N} \sum_{i=1}^N \langle \dot{\mathbf{x}}_i(t), \mathbf{x}_i(t) - \bar{\mathbf{x}} \rangle \\ &\quad \Downarrow \quad L_N(t) \bar{\mathbf{x}} = 0 \\ &= \frac{1}{N} \sum_{i=1}^N \left\langle (L_N(t)(\mathbf{x}(t) - \bar{\mathbf{x}}))_i, (\mathbf{x}_i(t) - \bar{\mathbf{x}}) \right\rangle \\ &\quad \Downarrow \quad \text{Def. of } \langle \cdot, \cdot \rangle_N \\ &= -\langle L_N(t)(\mathbf{x}(t) - \bar{\mathbf{x}}), (\mathbf{x}(t) - \bar{\mathbf{x}}) \rangle_N \\ &\quad \Downarrow \quad \text{Def. of } \lambda_2(\mathbf{A}_N(t)) \\ &\leq -\lambda_2(\mathbf{A}_N(t)) \|\mathbf{x}(t) - \bar{\mathbf{x}}\|_N^2 \\ &\quad \Downarrow \quad \text{Def. of } \mathcal{V}(\mathbf{x}(t)) \\ &= -\lambda_2(\mathbf{A}_N(t)) \mathcal{V}(\mathbf{x}(t))\end{aligned}$$

↪ Grönwall lemma and we're done!



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## Consensus – About algebraic and graph connectivity (1)

Theorem (Characterisation of graph connectivity)[Mohar'91]

A symmetric graph  $\mathbf{A}_N = (a_{ij})_{i,j=1}^N$  is **strongly connected**, i.e.

for all  $i, j$  there exists  $i = k_1, \dots, k_m = j$  s.t.  $a_{k_l k_{l+1}} > 0$

if and only if  $\lambda_2(\mathbf{A}_N) > 0$ .

**Question:** What happens when  $\mathbf{A}_N$  is not symmetric ?

Theorem (Characterisation of graph connectivity)[Wu'05]

A graph  $\mathbf{A}_N$  is a **disjoint union of str. connected components** (“DUSCC”) if and only if there exists  $(v_1, \dots, v_N) \in (\mathbb{R}_+^*)^N$  s.t.

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**Remark:**  $\mathbf{L}_N^*(1, \dots, 1) = 0$  if  $\mathbf{A}_N$  is symmetric  $\rightsquigarrow$  always DUSCC!

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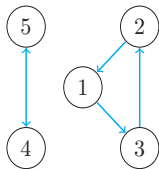
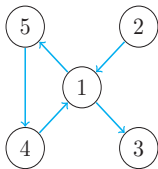
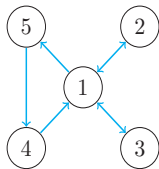
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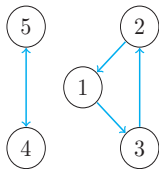
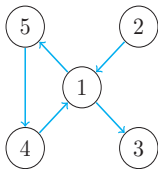
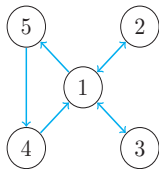
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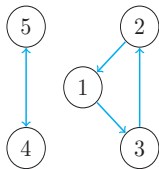
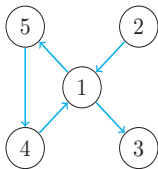
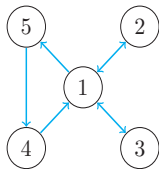
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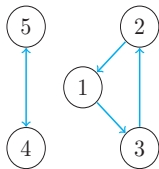
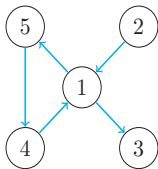
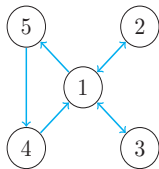
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# Outline of the talk

Multi-agent systems – From microscopic to macroscopic models

Review of consensus methods for microscopic cooperative systems

Consensus analysis in the context of graphon dynamics



## Graphon dynamics – *Adjacency and graph-Laplacian*

We consider the graphon dynamics

$$\partial_t x(t, i) = \int_I a(t, i, j) \phi(|x(t, i) - x(t, j)|) (x(t, j) - x(t, i)) dj$$

where  $a(t) \in L^\infty(I \times I, [0, 1])$  represents the **communications**.

Definition (Adjacency, degree and graph-Laplacian operators)

We define the **adjacency** operator by

$$\mathcal{A}(t, x) y : i \in I \mapsto \int_I a(t, i, j) \phi(|x(i) - x(j)|) y(j) dj,$$

as well as the **graph-Laplacian**  $\mathbb{L}(t, x) : L^2(I, \mathbb{R}^d) \rightarrow L^2(I, \mathbb{R}^d)$

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$\hookrightarrow$  Semilinear reformulation of the dynamics

$$\dot{x}(t) = -\mathbb{L}(t, x(t))x(t). \quad (\text{GD})$$

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$$\mathcal{A}(t, \mathbf{x}) \mathbf{y} : i \in I \mapsto \int_I \mathbf{a}(t, i, j) \phi(|\mathbf{x}(i) - \mathbf{x}(j)|) \mathbf{y}(j) \mathrm{d}j,$$

as well as the **graph-Laplacian**  $\mathbb{L}(t, \mathbf{x}) : L^2(I, \mathbb{R}^d) \rightarrow L^2(I, \mathbb{R}^d)$

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$\hookrightarrow$  Semilinear reformulation of the dynamics

$$\dot{\mathbf{x}}(t) = -\mathbb{L}(t, \mathbf{x}(t)) \mathbf{x}(t). \quad (\text{GD})$$

## Graphon dynamics – *Adjacency and graph-Laplacian*

We consider the graphon dynamics

$$\partial_t \mathbf{x}(t, i) = \int_I \mathbf{a}(t, i, j) \phi(|\mathbf{x}(t, i) - \mathbf{x}(t, j)|) (\mathbf{x}(t, j) - \mathbf{x}(t, i)) \mathrm{d}j$$

where  $\mathbf{a}(t) \in L^\infty(I \times I, [0, 1])$  represents the **communications**.

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# Graphon dynamics – *Scrambling and diameter estimate*

Definition (Scrambling coefficient and diameter)[BPDS'22]

We define the **scrambling coefficient** of a graphon  $\mathcal{A}$  by

$$\eta(\mathcal{A}) := \inf_{i,j \in I} \int_I \min\{a(i,k), a(j,k)\} dk,$$

as well as the **diameter** of a map  $x \in L^\infty(I, \mathbb{R}^d)$

$$\mathcal{D}(x) := \sup_{i,j \in I} |x(i) - x(j)|.$$

Theorem (Quantitative diameter decay)[BPDS'22]

For each  $x^0 \in L^\infty(I, \mathbb{R}^d)$ , there exists  $\phi_0 > 0$  such that

$$\mathcal{D}(x(t)) \leq \mathcal{D}(x^0) \exp \left( -\phi_0 \int_0^t \eta(\mathcal{A}(s)) ds \right).$$

Two technical novelties

- ◇ No **stochastic normalisation** trick  $\rightsquigarrow$  **Geometric** argument.
- ◇  $t \mapsto \mathcal{D}(x(t))$  not diff.  $\rightsquigarrow$  approx. with **Scorza-Dragoni**.

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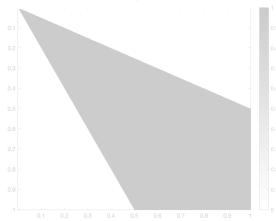
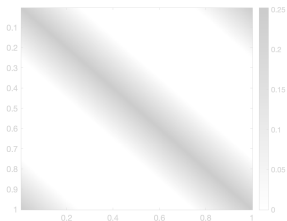
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Definition (Graphon connectivity)[Boudin,Salvarini&Trélat'21]

A graphon  $\mathcal{A}$  is **strongly connected** if the following holds.

- (i) (**Connectivity**) For  $\mathcal{L}^1$ -almost every  $i, j \in I$ , there exists  $i = k_1, \dots, k_m = j$  such that  $k_{l+1} \in \text{supp}(a(k_l, \cdot))$ .
- (ii) (**Degree lower-bound**)  $\inf_{i \in I} \int_I a(i, j) dj \geq \delta > 0$ .



Theorem (Canonical kernel of  $\mathbb{L}^*$ )[Boudin,Salvarini&Trélat'21]

If  $\mathcal{A}$  is strongly connected, there exists a unique  $v \in L^2(I, \mathbb{R}_+^*)$  s.t.

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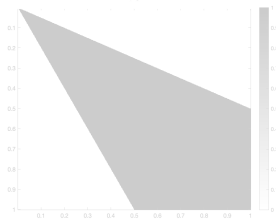
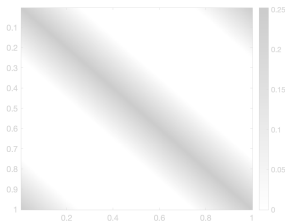


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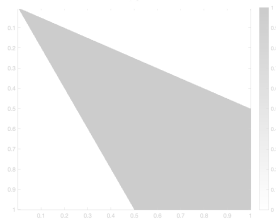
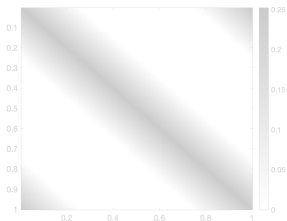
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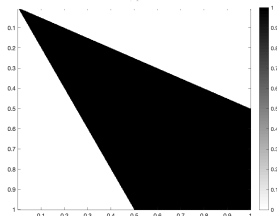
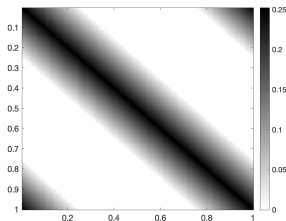
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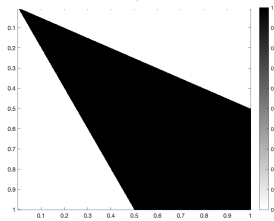
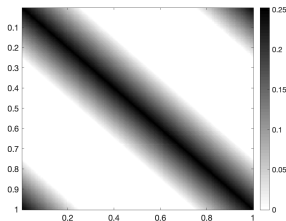
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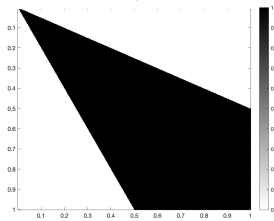
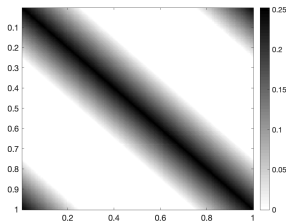
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# Graphon dynamics – Algebraic and graphon connectivity

## Definition (Generalised algebraic connectivity)

We define the **algebraic connectivity** of a DCUSCC graphon  $\mathcal{A}$  by

$$\lambda_2(\mathcal{A}) := \inf_{x \in \mathcal{C}^\perp} \frac{\langle \mathbb{L}_v x, x \rangle_{L^2(I)}}{\|x\|_{L^2(I)}^2}$$

where

- ◇  $\mathcal{C} := \{x \in L^2(I, \mathbb{R}^d) \text{ constant}\}$  is the **consensus manifold**,
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## Theorem (On algebraic and graphon connectivity)[BPDS'22]

For a graphon  $\mathcal{A}$ , the following connectivity characterisations hold.

- ◇ If  $\mathcal{A}$  is **symmetric**, strong connectedness  $\iff \lambda_2(\mathcal{A}) > 0$ .
- ◇ If  $\mathcal{A}$  is **DCUSCC**, strong connectedness  $\iff \lambda_2(\mathcal{A}) > 0$ .

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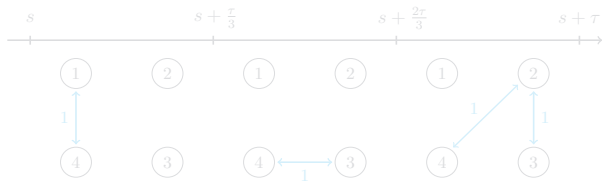


# Graphon dynamics – Variance decay and connectivity (1)

Theorem (Variance decay for symmetric graphons)[BPDS'22,BF'21]

Suppose that  $\mathcal{A}(t)$  is symmetric for a.e.  $t \in \mathbb{R}_+$ . Then for each  $x^0 \in L^\infty(I, \mathbb{R}^d)$  and every  $\tau > 0$ , there exist  $\phi_0, \alpha_\tau, \gamma_\tau > 0$  s.t.

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Time averaging

$$\mathcal{V}(x) := \frac{1}{2} \int_I \int_I |x(i) - x(j)|^2 dj di$$

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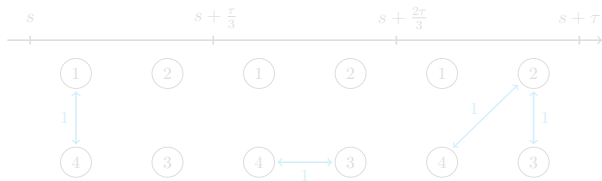
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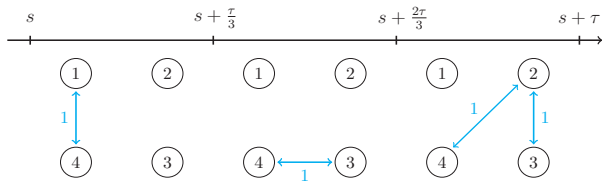
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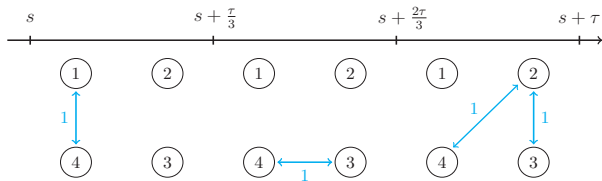
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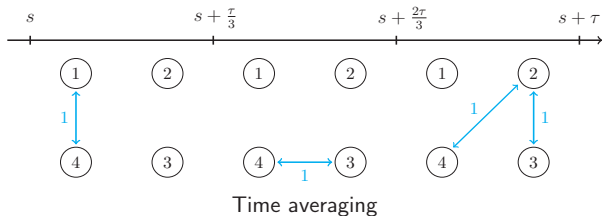
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$\hookrightarrow$  Equality between the **in-degree** and **out-degree** at each node.

### Theorem (Variance decay for balanced graphons)[BPDS'22]

Suppose that  $\mathcal{A}(t)$  is balanced for  $\mathcal{L}^1$ -almost every  $t \in [0, T]$  and that  $\phi(\cdot) \equiv 1$ . Then for each  $x^0 \in L^\infty(I, \mathbb{R}^d)$ , it holds that

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**Open problem:** Average condition like in the symmetric case ?

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Suppose that  $\mathcal{A}(t)$  is balanced for  $\mathcal{L}^1$ -almost every  $t \in [0, T]$  and that  $\phi(\cdot) \equiv 1$ . Then for each  $x^0 \in L^\infty(I, \mathbb{R}^d)$ , it holds that

$$\mathcal{V}(x(t)) \leq \mathcal{V}(x^0) \exp \left( - \int_0^t \lambda_2(\mathcal{A}(s)) \mathrm{d}s \right).$$

**Open problem:** Average condition like in the symmetric case ?

## Graphon dynamics – Variance decay and connectivity (2)

Definition (Balanced interaction topology)

A graphon  $\mathcal{A}$  is said to be **balanced** if  $\mathbb{L}^*1 = 0$ , namely

$$\int_I a(i, j) \mathrm{d}j = \int_I a(j, i) \mathrm{d}j.$$

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**Issue:** If  $\mathcal{A}(t)$  is DCUSCC  $\rightsquigarrow \mathcal{V}(\cdot)$  **not Lyapunov** anymore!

Definition (Weighted variance)

If a graphon  $\mathcal{A}$  is DCUSCC, we define the **weighted variance** by

$$\mathcal{V}_v(x) := \frac{1}{2} \int_I \int_I v(i)v(j) |x(i) - x(j)|^2 \mathrm{d}j \mathrm{d}i.$$

Theorem (Variance decay for DCUSCC dwelling graphons)[BPDS'22]

Suppose that  $\mathcal{A}(t)$  is DCUSCC for  $\mathcal{L}^1$ -a.e.  $t \in \mathbb{R}_+$  with

$$\nu \leq v(t, i) \leq \frac{1}{\nu} \quad \text{for } \mathcal{L}^1\text{-a.e. } i \in I,$$

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## Graphon dynamics – *Link between $L^2$ - and $L^\infty$ -consensus*

**Observation:** Under the **sufficient condition** for  $L^2$ -consensus

$$\lambda_2\left(\frac{1}{\tau} \int_t^{t+\tau} \mathcal{A}(s) ds\right) \geq \mu \quad \text{or} \quad \frac{1}{\tau} \int_t^{t+\tau} \lambda_2(\mathcal{A}(s)) ds \geq \mu$$

we **numerically** observed  $L^\infty$ -consensus  $\leadsto$  true in general ?

**Theorem (Equivalence between  $L^2$ - and  $L^\infty$ -consensus)**[BPDS'22]

Suppose that there exist constants  $(\tau, \mu) \in \mathbb{R}_+^* \times (0, 1]$  s.t.

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$$\|x(t) - x^\infty\|_{L^2(I)} \xrightarrow[t \rightarrow +\infty]{} 0$$

for some  $x^\infty \in \mathbb{R}^d$  **if and only if**

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### Wrap-up (Summary of presentation)

1. Convergence to **consensus** in **micro** and **macro** dynamics.
2. Generalisation of the **scrambling** and **Fiedler** numbers.
3. Some interesting **open problems** to investigate!



Thank you for your attention !

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