

Simplex Geometry of Graphs

Piet Van Mieghem

in collaboration with Karel Devriendt

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Google matrix: fundamentals, applications and beyond (GOMAX)
IHES, October 15-18, 2018



Outline

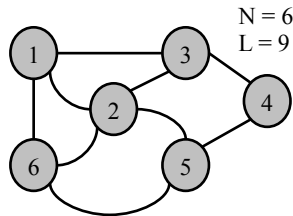


Background:
Electrical matrix equations

Geometry of a graph



Adjacency matrix A



$$A_{N \times N} = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$

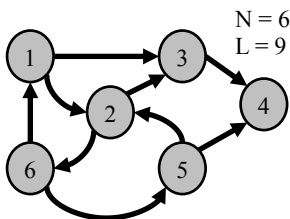
For an undirected graph: $A = A^T$ is symmetric

Number of neighbors of node i is the degree: $d_i = \sum_{k=1}^N a_{ik}$

if there is a link between node i and j , then $a_{ij} = 1$
else $a_{ij} = 0$



Incidence matrix B



- Label links (e.g.: $l_1 = (1,2)$, $l_2 = (1,3)$, $l_3 = (1,6)$, $l_4 = (2,3)$, $l_5 = (2,5)$, $l_6 = (2,6)$, $l_7 = (3,4)$, $l_8 = (4,5)$, $l_9 = (5,6)$)
- Col k for link $l_k = (i,j)$ is zero, except:
source node $i = 1 \rightarrow b_{ik} = 1$
destination node $j = -1 \rightarrow b_{jk} = -1$

$$B_{N \times L} = \begin{bmatrix} 1 & 1 & -1 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & 0 & 0 & 1 & -1 & 1 & 0 & 0 & 0 \\ 0 & -1 & 0 & -1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & -1 & -1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & -1 \\ 0 & 0 & 1 & 0 & 0 & -1 & 0 & 0 & 1 \end{bmatrix}$$

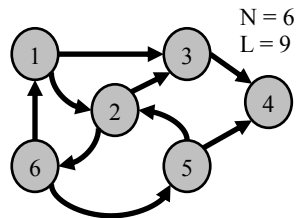
Col sum B is zero: $u^T B = 0$

where the all-one vector $u = (1,1,\dots,1)$

B specifies the directions of links



Laplacian matrix Q



$$Q_{N \times N} = \begin{bmatrix} 3 & -1 & -1 & 0 & 0 & -1 \\ -1 & 4 & -1 & 0 & -1 & -1 \\ -1 & -1 & 3 & -1 & 0 & 0 \\ 0 & 0 & -1 & 2 & -1 & 0 \\ 0 & -1 & 0 & -1 & 3 & -1 \\ -1 & -1 & 0 & 0 & -1 & 3 \end{bmatrix}$$

$$Q = BB^T = \Delta - A$$

$$\Delta = \text{diag}(d_1 \quad d_2 \quad \dots \quad d_N)$$

Since BB^T is symmetric, so are A and Q . Although B specifies directions, A and Q lost this info here.

Basic property: $Qu = 0$

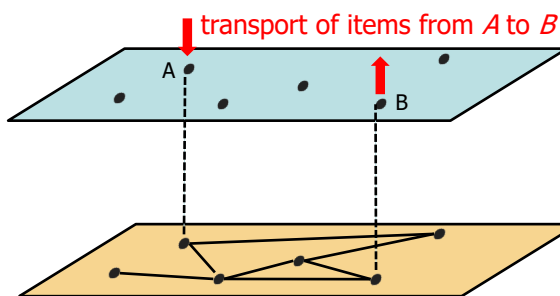
u is an eigenvector of Q
Belonging to eigenvalue $\mu = 0$

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$$Qu = BB^T u = 0 \quad \text{because} \quad 0 = u^T B = B^T u$$



Network: service(s) + topology



Service (function)

software, algorithms

Topology (graph)

hardware, structure

Service and topology

- own specifications
- both are, generally, time-variant
- service is often designed independently of the topology
- often more than 1 service on a same topology

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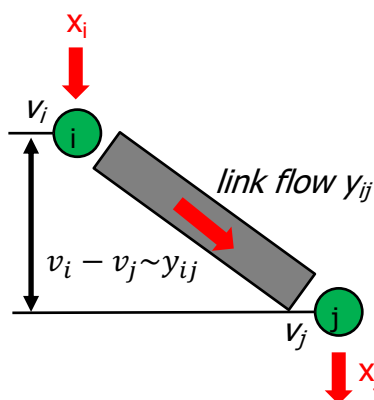
Function of network

- Usually, the function of a network is related to the *transport of items over its underlying graph*
- In man-made infrastructures: two major types of transport
 - Item is a **flow** (e.g. electrical current, water, gas,...)
 - Item is a **packet** (e.g. IP packet, car, container, postal letter,...)
- **Flow equations (physical laws)** determine transport (Maxwell equations (Kirchhoff & Ohm), hydrodynamics, Navier-Stokes equation (turbulent, laminar flow equations, etc.)
- **Protocols** determine transport of packets (IP protocols and IETF standards, car traffic rules, etc.)

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Linear dynamics on networks

Linear dynamic process: "proportional to" (\sim) graph of network



Examples:

- water (or gas) flow \sim pressure
- displacement (in spring) \sim force
- heat flow \sim temperature
- **electrical current \sim voltage**

$$\mathbf{x} = \mathbf{Q} \cdot \mathbf{v}$$

injected nodal current vector	=	\mathbf{Q}	.	\mathbf{v}	
		weighted Laplacian of the graph		nodal potential vector	

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Pseudoinverse of the Laplacian (review)

The inverse of the current-voltage relation $\mathbf{x} = \mathbf{Q}\mathbf{v}$
 is the voltage-current relation $\mathbf{v} = \mathbf{Q}^\dagger \mathbf{x}$
 subject to $\mathbf{u}^T \mathbf{x} = 0$ and $\mathbf{u}^T \mathbf{v} = 0$

The spectral decomposition

$$\tilde{\mathbf{Q}} = \sum_{k=1}^{N-1} \tilde{\mu}_k \mathbf{z}_k \mathbf{z}_k^T$$

allows us to compute the pseudoinverse (or Moore-Penrose inverse)

$$\mathbf{Q}^\dagger = \sum_{k=1}^{N-1} \frac{1}{\tilde{\mu}_k} \mathbf{z}_k \mathbf{z}_k^T$$

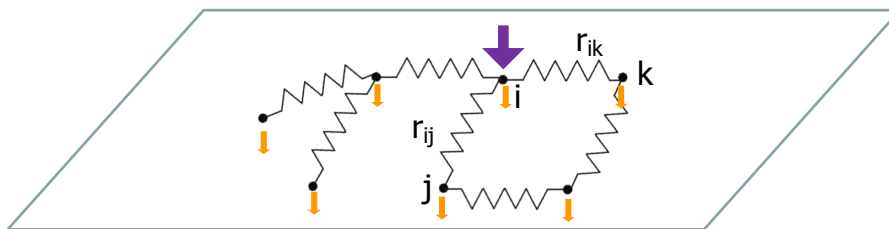
The effective resistance $N \times N$ matrix is $\tilde{\Omega} = \mathbf{u} \zeta^T + \zeta \mathbf{u}^T - 2\mathbf{Q}^\dagger$,
 where the $N \times 1$ vector $\zeta = (Q_{11}^\dagger, Q_{22}^\dagger, \dots, Q_{NN}^\dagger)$

An interesting graph metric is the effective graph resistance

$$R_G = N \mathbf{u}^T \zeta = N \text{trace}(\mathbf{Q}^\dagger) = N \sum_{k=1}^{N-1} \frac{1}{\mu_k}$$

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P. Van Mieghem, K. Devriendt and H. Cetinay, 2017, "Pseudo-inverse of the Laplacian and best spreader node in a network", Physical Review E, vol. 96, No. 3, p 032311.



Inverses: $\mathbf{x} = \mathbf{Q}\mathbf{v} \leftrightarrow \mathbf{v} = \mathbf{Q}^\dagger \mathbf{x}$ with voltage reference $\mathbf{u}^T \mathbf{v} = 0$

\mathbf{Q}^\dagger : pseudoinverse of the weighted Laplacian obeying $\mathbf{Q}\mathbf{Q}^\dagger = \mathbf{Q}^\dagger \mathbf{Q} = \mathbf{I} - \frac{1}{N} \mathbf{J}$
 $\mathbf{J} = \mathbf{u}\mathbf{u}^T$: all-one matrix \mathbf{u} : all-one vector

Unit current injected in node i
 $\mathbf{x} = \mathbf{e}_i - \frac{1}{N} \mathbf{u}$

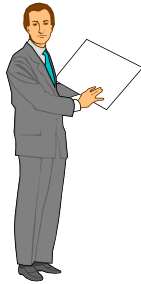


nodal potential of i
 $v_i = Q_{ii}^\dagger$

→ The best spreader is node k with $Q_{kk}^\dagger \leq Q_{ii}^\dagger$ for $1 \leq i \leq N$



Outline



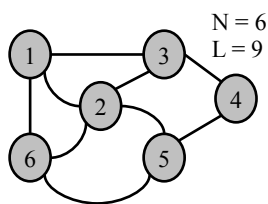
Background:
Electrical matrix equations

Geometry of a graph



Three representations of a graph

Topology domain



$$A_{N \times N} = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$

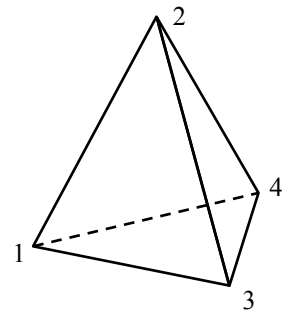
Spectral domain

$$A = A^T = X \Lambda X^T$$

$X_{N \times N}$: orthogonal
eigenvector matrix

$\Lambda_{N \times N}$: diagonal
eigenvalue matrix

Geometric domain



Undirected graph on
 N nodes
= **simplex** in Euclidean
 $(N-1)$ -dimensional space

Devriendt, K. and P. Van Mieghem, 2018, "The Simplex Geometry of Graphs",
Delft University of Technology, report20180717.
(<http://arxiv.org/abs/1807.06475>).



Encyclopedia of Mathematics and Its Applications 139

MATRICES AND GRAPHS IN GEOMETRY


Miroslav Fiedler

CAMBRIDGE

more information - www.cambridge.org/9780521461931

appeared in 2011

Miroslav Fiedler (1926-2015)




Father of "algebraic connectivity"
His 1972 paper: > 3400 citations

"This book comprises, in addition to auxiliary material, the research on which I have worked for over 50 years."


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TUDelft

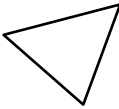
What is a simplex?



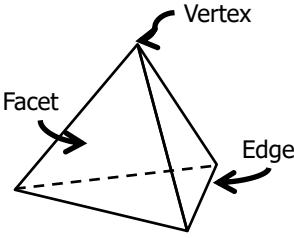
Point



Line Segment



Triangle



Tetrahedron

Roughly : a simplex is generalization of a triangle to N dimensions

Potential : Euclidean geometry is the oldest, mathematical theory

T. L. Heath, *The Thirteen Books of Euclid's Elements*, Vol. 1-3, Cambridge University Press, 1926

TUDelft

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Spectral decomposition weighted Laplacian (1)

Spectral decomposition: $Q = ZMZ^T$

where $M = \text{diag}(\mu_1, \mu_2, \dots, \mu_{N-1}, 0)$, because $Q \mathbf{u} = 0$

and the eigenvector matrix Z obeys $Z^T Z = Z Z^T = I$ with structure

$$\text{node } Z = \begin{bmatrix} (z_1)_1 & (z_2)_1 & \cdots & (z_N)_1 \\ (z_1)_2 & (z_2)_2 & \cdots & (z_N)_2 \\ \vdots & \vdots & \ddots & \vdots \\ (z_1)_N & (z_2)_N & \cdots & (z_N)_N \end{bmatrix} = \begin{bmatrix} (z_1)_1 & (z_2)_1 & \cdots & 1/\sqrt{N} \\ (z_1)_2 & (z_2)_2 & \cdots & 1/\sqrt{N} \\ \vdots & \vdots & \ddots & \vdots \\ (z_1)_N & (z_2)_N & \cdots & 1/\sqrt{N} \end{bmatrix}$$

↑
frequencies
(eigenvalues)

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$$Q = \sum_{k=1}^{N-1} \mu_k z_k z_k^T$$



Spectral decomposition weighted Laplacian (2)

Only for a positive semi-definite matrix, it holds that

$$Q = ZMZ^T = (Z\sqrt{M})(Z\sqrt{M})^T$$

The matrix $S = (Z\sqrt{M})^T$ obeys $Q = S^T S$ and has rank $N-1$
(row $N = 0$ due to $\mu_N = 0$)

$$S = \begin{bmatrix} \sqrt{\mu_1}(z_1)_1 & \sqrt{\mu_1}(z_1)_2 & \cdots & \sqrt{\mu_1}(z_1)_N \\ \sqrt{\mu_2}(z_2)_1 & \sqrt{\mu_2}(z_2)_2 & \cdots & \sqrt{\mu_2}(z_2)_N \\ \vdots & \vdots & \ddots & \vdots \\ \sqrt{\mu_{N-1}}(z_{N-1})_1 & \sqrt{\mu_{N-1}}(z_{N-1})_2 & \cdots & \sqrt{\mu_{N-1}}(z_{N-1})_N \\ 0 & 0 & \cdots & 0 \end{bmatrix}$$

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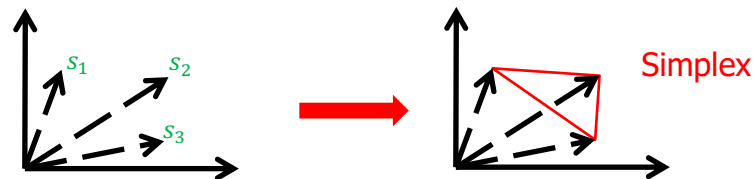
$$Q = \sum_{k=1}^{N-1} \mu_k z_k z_k^T$$



Geometrical representation of a graph

$$S = \begin{bmatrix} \sqrt{\mu_1}(z_1)_1 & \sqrt{\mu_1}(z_1)_2 & \cdots & \sqrt{\mu_1}(z_1)_N \\ \sqrt{\mu_2}(z_2)_1 & \sqrt{\mu_2}(z_2)_2 & \cdots & \sqrt{\mu_2}(z_2)_N \\ \vdots & \vdots & \ddots & \vdots \\ \sqrt{\mu_{N-1}}(z_{N-1})_1 & \sqrt{\mu_{N-1}}(z_{N-1})_2 & \cdots & \sqrt{\mu_{N-1}}(z_{N-1})_N \\ \hline 0 & 0 & \cdots & 0 \end{bmatrix}$$

The i -th column vector $s_i = (\sqrt{\mu_1}(z_1)_i, \sqrt{\mu_2}(z_2)_i, \dots, \sqrt{\mu_N}(z_N)_i = 0)$ represents a point p_i in $(N-1)$ -dim space (because S has rank $N-1$)



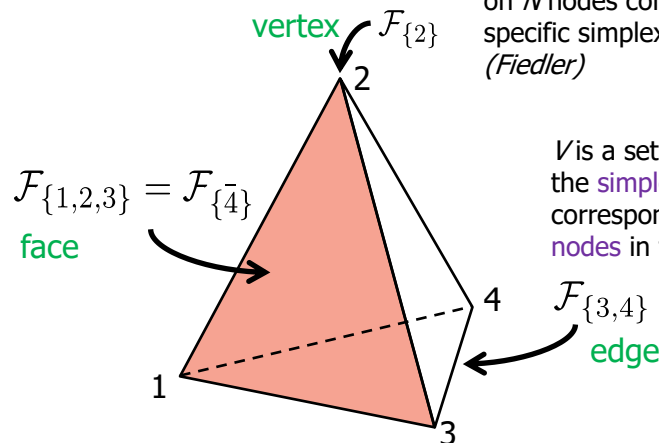
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Simplex geometry: omit zero row, $S_{N \times N} \rightarrow S_{(N-1) \times N}$



Faces of a simplex

Each connected, undirected graph on N nodes corresponds to 1 specific simplex in $N-1$ dimensions (Fiedler)



V is a set of vertices of the simplex in \mathbb{R}^{N-1} , corresponding to a set of nodes in the graph G

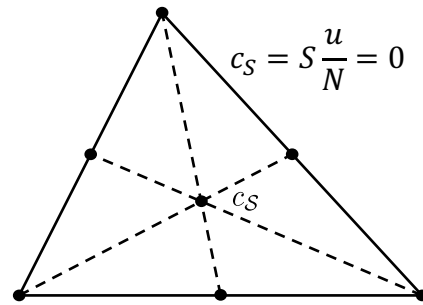
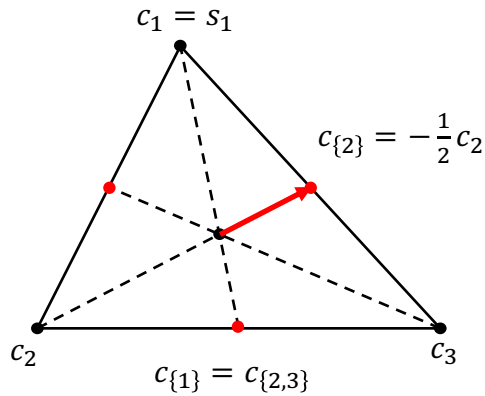
A face $F_V = \{p \in \mathbb{R}^{N-1} | p = Sx_V \text{ with } (x_V)_i \geq 0 \text{ and } u^T x_V = 1\}$

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The vector $x_V \in \mathbb{R}^N$ is a **barycentric coordinate** with $\begin{cases} (x_V)_i \in \mathbb{R} & \text{if } i \in V \\ (x_V)_i = 0 & \text{if } i \notin V \end{cases}$

Centroids

$c_V = S \frac{u_V}{|V|}$ is the centroid of face F_V with $(u_V)_i = 1_{i \in V}$



a centroid of a face is a vector

centroid of simplex is origin

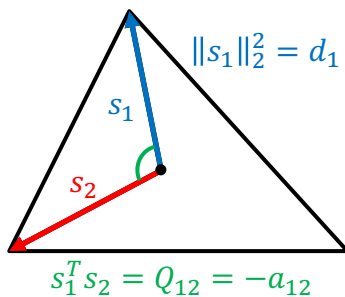
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$$u_V = u - u_V \rightarrow |V|c_V = S(u - u_V) = -(N - V)c_V$$



Geometric representation of a graph

$$\begin{aligned} \|s_i\|_2^2 &= d_i & \|s_i - s_j\|_2^2 &= (s_i - s_j)^T (s_i - s_j) = s_i^T s_i + s_j^T s_j - 2s_i^T s_j \\ & & &= Q_{ii} + Q_{jj} - 2Q_{ij} \\ & & &= d_i + d_j + 2a_{ij} \text{ for } i \neq j, \text{ else zero} \end{aligned}$$



The matrix with off-diagonal elements $d_i + d_j + 2a_{ij}$ is a **distance matrix** (if the graph G is connected)

The geometric graph representation is not unique (node relabeling changes Z)

$$s_i^T s_j = \sum_{k=1}^{N-1} \sqrt{\mu_k} (z_k)_i \sqrt{\mu_k} (z_k)_j = \sum_{k=1}^{N-1} \mu_k (z_k z_k^T)_{ij} = Q_{ij}$$

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$$Q = \sum_{k=1}^{N-1} \mu_k z_k z_k^T \text{ and } Q = S^T S$$



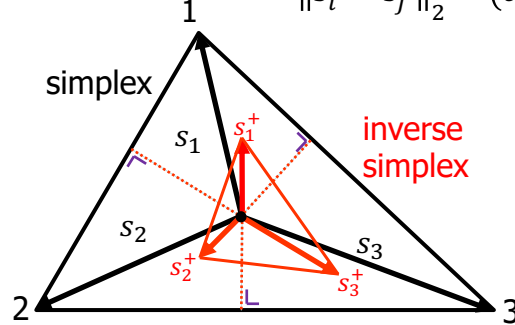
Geometry of a graph (dual representation)

Spectral decomposition: $Q^+ = ZM^+Z^T = (Z\sqrt{M^+})(Z\sqrt{M^+})^T$

The matrix $S^+ = (Z\sqrt{M^+})^T$ has rank $N-1$ and $Q^+ = (S^+)^T S^+$

The i -th column vector s_i^+ obeys

$$\|s_i^+ - s_j^+\|_2^2 = (e_i - e_j)^T Q^+ (e_i - e_j) = \omega_{ij}$$



From $S^T S^+ = I - \frac{uu^T}{N}$:

$$\begin{cases} s_j^T s_i^+ = -\frac{1}{N} \\ s_i^T s_i^+ = 1 - \frac{1}{N} \end{cases}$$

$$\Rightarrow (s_k - s_j)^T s_i^+ = 0$$

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ω_{ij} is the effective resistance between node i and j



Volume of simplex and inverse simplex of a graph

Volume of the simplex

$$V_G = \frac{N}{(N-1)!} \sqrt{\xi}$$

where the number of (weighted) spanning trees ξ is $\xi = \frac{1}{N} \prod_{k=1}^{N-1} \mu_k$

Volume of the inverse simplex

$$V_G^+ = \frac{1}{(N-1)! \sqrt{\xi}}$$

Hence:

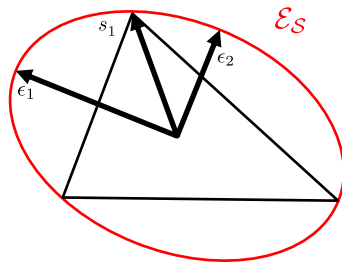
$$\frac{V_G}{V_G^+} = N\xi = \prod_{k=1}^{N-1} \mu_k$$

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K. Menger, "New foundation of Euclidean geometry",
American Journal of Mathematics, 53(4):721-745, 1931



Steiner ellipsoid of simplex



projection $s_1^T \epsilon_2 = \mu_2(z_2)_1$

semi-axis: $\|\epsilon_2\| = \sqrt{\frac{N}{N-1}} \mu_2$

volume:

$$V_{\mathcal{E}_S} = \frac{\pi^{N/2}}{\Gamma(N/2+1)} \frac{N^{N/2}}{(N-1)^{N/2}} \sqrt{\prod_{k=1}^N \mu_k}$$

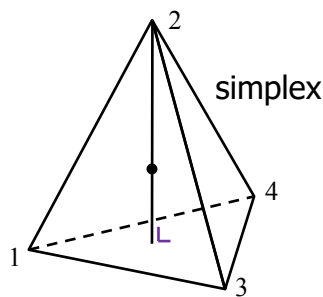
Hence,

$$V_{\mathcal{E}_S}^2 = \frac{(N\pi)^N}{(\Gamma(N/2+1))^2 (N-1)^N} \prod_{k=1}^N \mu_k$$

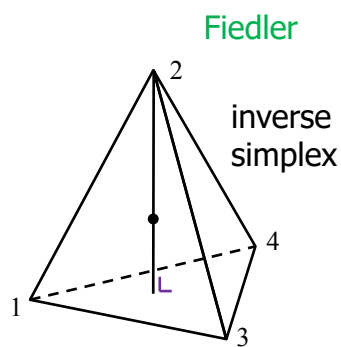
$$V_{N\text{-ellipsoid}} = \frac{\pi^{N/2}}{\Gamma(N/2+1)} \prod_{k=1}^N \epsilon_k$$



altitude(s) in a simplex



$$\|a_{\{2\}}\|^2 = \frac{1}{Q_{22}^+}$$



$$\|a_2^+\|^2 = \frac{1}{d_2}$$

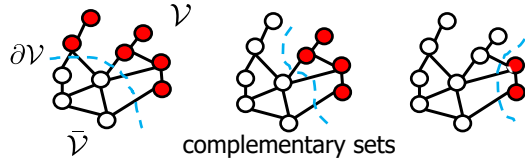
The altitude from a vertex s_i^+ to the complementary face $F_{\{i\}}^+$ in the inverse simplex (dual graph representation) has a length equal to the inverse degree of node i

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recall that $Q_{ii}^+ = v_i$ (nodal potential, best spreader)



Cut size of graph and altitude of simplex



$$\text{Cut size: } |\partial V| = u_V^T Q u_V$$

$$Q = S^T S$$

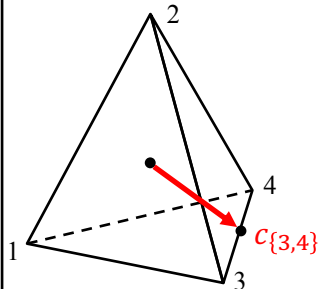
$$\begin{aligned} |\partial V| &= u_V^T S^T S u_V \\ &= \|S u_V\|_2^2 \\ &= \|c_V\|_2^2 V^2 \end{aligned}$$

altitude:

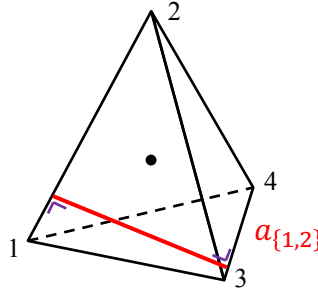
vector from face F_V
to face F_V and
orthogonal to both
faces

$$|\partial^+ V| = u_V^T Q^+ u_V$$

$$\|a_V^+\|^2 = \frac{1}{|\partial V|}$$



$$\|c_{\{3,4\}}\|^2 = \frac{\partial\{3,4\}}{4}$$



$$\|a_{\{1,2\}}\|^2 = \frac{1}{\partial^+\{1,2\}}$$

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$c_V = S \frac{u_V}{|V|}$ is the centroid of face F_V with $(u_V)_i = 1_{i \in V}$



Metrics $\sqrt{\omega_{ij}}$ and ω_{ij}

$$\|s_i^\dagger - s_j^\dagger\|_2 = \sqrt{\omega_{ij}}$$

the Euclidean distance between vertices of inverse simplex S^\dagger



vertices of S^\dagger are an embedding of nodes of the graph G
according to the metric $\sqrt{\omega_{ij}}$ (a.o. obeying the triangle inequality)

$$\text{Also } \|s_i^\dagger - s_j^\dagger\|_2^2 = \omega_{ij} \text{ is a metric}$$



Inverse simplex S^\dagger of the graph G with positive link weights is
hyperacute

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Unpublished recent work with K. Devriendt



Generalization metrics

Q^\dagger is the Gram matrix of a hyperacute simplex S^+ \rightarrow determines a metric



$m_{ij}^{(f)} = (e_i - e_j)^T (f(Q))^\dagger (e_i - e_j)$ is a metric when $f(Q)$ is a Laplace matrix

"Statistical physics" metrics on a graph

$$m_{ij}^{(stat(g))} = (e_i - e_j)^T (e^{(Q+pQ_K)t} - I + gQ_K)^\dagger (e_i - e_j) = \sum_{k=1}^{N-1} \frac{((z_k)_i - (z_k)_j)^2}{e^{(\mu_k + pN)t + g}}$$

with $t = \frac{1}{k_B T}$ and $pN = -E_f$ (chemical potential or Fermi energy)

$g = -1 \rightarrow$ Bose-Einstein

$g = 0 \rightarrow$ Maxwell-Boltzmann

$g = 1 \rightarrow$ Fermi-Dirac

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Unpublished recent work with K. Devriendt



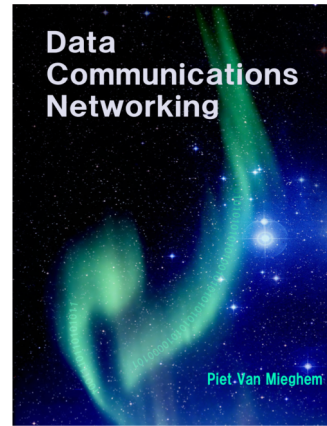
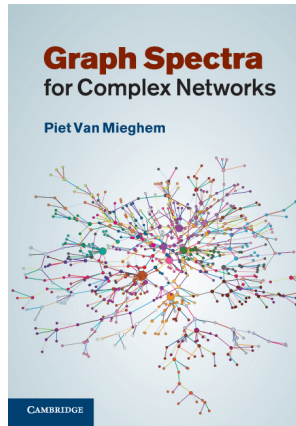
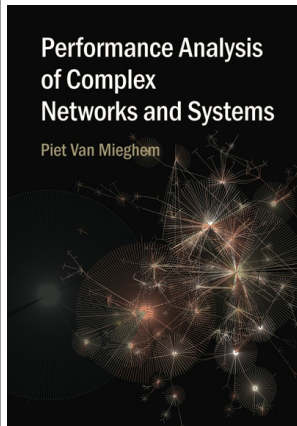
Summary

- Linearity between process and graph naturally leads to the weighted Laplacian Q and its pseudoinverse Q^\dagger
- Spectral decomposition of the weighted Laplacian Q and its pseudoinverse Q^\dagger provides an $N-1$ dimensional simplex representation of each graph,
 - allowing computations in the $N-1$ dim. Euclidean space (in which a distance/norm is defined)
 - geometry for (undirected) graphs
- **Open:** "Which network problems are best solved in the simplex representation?"

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Books



Articles: <http://www.nas.ewi.tudelft.nl>

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Thank You

Piet Van Mieghem
NAS, TU Delft
P.F.A.VanMieghem@tudelft.nl

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