Arc spaces and some adjacency problems of plane curves.

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Joint work in progress with Javier Fernández de Bobadilla and Patrick Popescu-Pampu

Arcspace of
$$(\mathbb{C}^2, 0)$$
.

Arc (through the origin) of \mathbb{C}^2 : germ of parametrization through the origin:

$$\begin{array}{rcl} \gamma: & (\mathbb{C},0) & \longrightarrow & (X,O) \subset (\mathbb{C}^2,O) \\ & t & \longmapsto & \left(\sum_i a_i^1 t^i, ..., \sum_i a_i^n t^i\right) \\ & 0 & \longmapsto & O \end{array}$$

Formal arcs are considered: the power series may not converge. It is an infinite affine space.

It is irreducible.

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$$N_i = \{\gamma : \widetilde{\gamma}(0) \in E_i\}$$

The *Nash set* is its closure \overline{N}_i .

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- They are cylindrical: they are determined in order *k*.

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- They are cylindrical: they are determined in order *k*.
- They have finite codimension.
- They are all different: $\overline{N}_i \neq \overline{N}_j$.

Nash sets.

Take a composition of blow ups in point above the origin.

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$$\alpha_{s}(t) = (t^5 + st^3; t^4 + st^4)$$

with $\alpha_s \in N_F$ and $\alpha_0 \in N_E$, then $\alpha_0 \in \overline{N}_F$. Nash sets.

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• GENERALISED NASH PROBLEM: Determine when $\overline{N}_E \subset \overline{N}_F$.

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- Surface singularities (Nash Conjecture, Theorem 2011, J. Fernandez de Bobadilla, M. P. P.): The components of the arcspace are in bijection with exceptional components of the minimal resolution.
- Higher dimensional case (partial result 2014, T. de Fernex, R. Docampo). Components in terminal models give components of the space of arcs. But there are more... still open.

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If some \overline{N}_E isn't a component of the space of arcs, then $\overline{N}_E \subseteq \overline{N}_F$ for some F.

Generalised Nash Problem: describe the inclusions/adjacencies $\overline{N}_E \subset \overline{N}_F$.

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Trivial inclusions: F < E (*E* dominates *F*) implies $N_E \subset \overline{N}_F$



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How to check if $\overline{N}_E \subset \overline{N}_F$?

Theorem (Fernández de Bobadilla, 2009)

Given two divisors above $O \in \mathbb{C}^2$, the following are equivalent:

- ② there exists a convergent family of convergent arcs α realising the inclusion with $\alpha_0 \in \dot{N}_E$ and $\alpha_s \in N_F$.
- for any convergent arc γ ∈ N_E there exists a family of convergent arcs α realising the inclusion with α₀ = γ (and α_s ∈ N_F).

 $\overline{N}_E \subset \overline{N}_F$ doesn't imply F < E.

• $(t^5 + s^3 t^3, (1 + s^4)t^4)$, two different tangents for s = 0 and $s \neq 0$.





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• Multiple examples:



A divisor is determined by the combinatorics + moduli... Minimal model of a divisor by blowing up a finite number of (free or satellite) points.

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Observe: not all the divisors are the final one for the minimal embedded resolution of a branch, but only the blow ups of satellite points.

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Take into account the contact order between E and F, and much more (moduli for free points also count apriori!)...



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Talk about combinatorics of the pair (E, F). We write $(E, F) \equiv (E', F')$... Domination relation F < E = Inclusion of Enriques diagrams.

Valuative criterion in arc spaces (A. Reguera, C. Plenat, S. Ishii...)

- Divisorial valuation ord_E = vanishing order along the divisor E.
- Can be computed intersecting with appropriate arcs γ in N_E :

$$ord_E(f) = I_O(f, \gamma) = ord_t(f \circ \gamma(t)).$$

• Choosing a family of arcs with an approviate $\alpha_0 \in N_E$ and $\alpha_s \in N_F$ we get

 $ord_F(h) \leq ord_t(h \circ \alpha_s(t)) \leq ord_t(h \circ \alpha_0(t)) = ord_E(h)$ forall $h \in \mathbb{C}[[x, y]].$

Valuative criterium in arc spaces (A. Reguera, C. Plènat, S. Ishii...)

- A. Reguera for rational surfaces in Manuscripta math. 1995.C. Plénat in general in Annal Inst. Fourier 2005:
 - $\overline{N}_E \subset \overline{N}_F$ implies $ord_E \leq ord_F$
- S. Ishii in Maximal Divisorial Sets:
 - if F is toric, also the converse is true: $ord_E \leq ord_F$ implies $\overline{N}_E \subset \overline{N}_F$.
 - She found a counterexample for the converse in general (ord_F ≤ ord_E but N_E ⊈ N_F).



$F < E \Rightarrow \overline{N}_E \subset \overline{N}_F \Rightarrow ord_F \leq ord_E$

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Other tools to rule out adjacencies?

• $\overline{N}_E \subsetneq \overline{N}_F$ implies $codim(\overline{N}_F) < codim(\overline{N}_E)$.

In [de Fernex, Ein, Ishii, Lazarsfeld, Mustata'200?]:

$$codim(\overline{N}_E) = 1 + disc(E, \mathbb{C}^2).$$

The *discrepancy* of *E* is the coefficient of *E* in K_{X/\mathbb{C}^2} where $\pi : X \to \mathbb{C}^2$ is any model where *E* appears.

- It is not a sufficient criterium (even with toric examples)
- Neither plus the valuative criterium (counterexample of Ishii with the same discrepancy).

The problem turns very difficult...

Example of topological types.



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- $ord_F \leq ord_E$
- disc(E) + 1 = codim(N_E) = 21 > codim(N_F) = disc(F) + 1 = 17
- $\overline{N}_E \nsubseteq \overline{N}_F$

Recall that $\overline{N}_E \subset \overline{N}_F$ depends on the relative position of E and F, but ... we don't know a priori that it is a combinatorial problem, it also depends on the moduli of the free points!.

Theorem

Assume there exists a wedge α realising the adjacency $\overline{N}_E \subset \overline{N}_F$. If $(E', F') \equiv (E, F)$ then there exists a wedge realising the adjacency $\overline{N}_{E'} \subset \overline{N}_{F'}$.

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Corollary

Assume we have that $\overline{N}_E \subset \overline{N}_F$. Let i_0 be the contact order between E and F. Then, we have that

$$\bigcup_{E'\equiv_{\geq i_0}E}\overline{N}_{E'}\subset\bigcap_{F'\equiv_{\geq i_0}F}\overline{N}_{F'}$$

where $A \equiv_{\geq i_0} B$ means that A has the same combinatorics as B and their contact order is $\geq i_0$.

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We improve the log-discrepancy inequality in many cases. Other conjectures...

$(E,F) \equiv (E',F') \Rightarrow [\overline{N}_E \subsetneq \overline{N}_F \Leftrightarrow \overline{N}_{E'} \subsetneq \overline{N}_{F'}]$

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To change the complex structure we can use:

 Let g : X → Y is a non-ramified covering of differentaible manifolds. A complex structure on Y can be lifted to X so that g is a local biholomorphism.

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- (Grauert-Remmert) Let A be a normal analytic space, let $B \subset A$ be a closed analytic subset such that $A \setminus B$ is dense in A. Let

$$f: U \to A \setminus B$$

be a finite and étale analytic morphism. Then there exists a finite analytic extension

$$\bar{f}: V \to A$$

from a normal analytic space V. Moreover V is unique up to isomorphism.

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We can assume the wedge $\alpha : \mathbb{C}^2 \to \mathbb{C}^2$ is algebraic, that is there exists polynomials $F_1, F_2 \in \mathbb{C}[s, t, x, y]$ such that

$$F_1(s,t,\alpha_1(s,t)) = F_2(s,t,\alpha_2(s,t)) = 0.$$

Coming back to the valuative criterium...

Recall: $\overline{N}_E \subset \overline{N}_F$ implies there exists a family of parametrizations $\alpha(t, s)$ with $\alpha_0(t) \in \dot{N}_E$ and $\alpha_s \in \dot{N}_F$ for all $s \in \Lambda \setminus \{0\}$.

• Deforming a little α , we can assume that

$$\alpha^{-1}(\mathcal{O}) = \{0\} \times \Lambda.$$

• The equation F(x, y, s) of

$$\mathit{Im}[(t,s)\mapsto (\alpha(t,s),s)\in \mathbb{C}^2 imes \Lambda]$$

gives a deformation of plane curves given by $f_s(x, y) := F(x, y, s)$ where $f_0(x, y) = 0$ lifts transversally to E and all $f_s(x, y) = 0$ lift transversally to F for all $s \neq 0$.

 These deformations have a special property: for s ≠ 0 they can be resolved simultaneously by a sequence of blow ups, they fix the free points (for F).

Let f_s be a deformation fixing the free points. If f₀ = 0 has strict transform transverse to some E and f_s = 0 have strict transforms transverse to a fixed F for all s ≠ 0, then

 $ord_F(h) \leq ord_E(h) \ \forall h \in \mathbb{C}[[x, y]].$

We have $I_O(h, f_s) \leq I_O(h, f_0)$ but is not enough...

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Proof

Take embedded resolution $(\tilde{X}, D = \bigcup_i D_i) \to (\mathbb{C}^2, O)$ of $f_s = 0$ and $f_0 = 0$. Look at it in family $\tilde{X} \times \Lambda \to \mathbb{C}^2 \times \Lambda$.

Let Y be the strict transform of F = 0 ($F(x, y, s) := f_s(x, y)$). Observe

$$Y_s = \{\widetilde{f_s = 0}\}$$
 for $s \neq 0$
 $Y_0 = \{\widetilde{f_0 = 0}\} + \sum_k d_k D_k,$ with $d_k \ge 0$

We get $Y_0 \cdot D_i = Y_s \cdot D_i$ for any *i*. Putting $M = (D_i \cdot D_j)$, $(E = D_0, F = D_n)$,

$$egin{aligned} &(1,0,...,0)^t + M(d_1,..,d_n)^t = (0,...,0,1)^t. \ &-M^{-1}(1,0,...,0,-1)^t = (d_1,...,d_n)^t \geq 0. \end{aligned}$$

and the entries of $-M^{-1}$ are exactly $ord_{D_i}(h_{D_i}) = I_O(h_{D_i}, h_{D_i})$.

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Reciprocally, if

$$ord_F(h) \leq ord_E(h) \ \forall h \in \mathbb{C}[[x, y]],$$

then, taking $h_E = 0$ and $h_F = 0$ with strict transform transverse to E and F in a model of E + F, then

$$h_E + s \cdot h_F$$

have strict transform transverse to F for $s \neq 0$ small enough. (Also proved by M. Alberich y J. Roe).

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Proof

Check it works.

• Let f_s be a deformation fixing the free points. If $f_0 = 0$ has strict transform transverse to some E and $f_s = 0$ have strict transforms transverse to a fixed F for all $s \neq 0$, then

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Summarizing:

Proposition

Let E and F be two prime divisors. There exists a deformation f_s of a curve $f_0 = 0$ that lifts transversal to E that fixes the free points for F ($f_s = 0$ has strict transform transverse to F for $s \neq 0$) if and only if $ord_F \leq ord_E$.

Adjacency problems.

- CLASSICAL ONE: Given two topological types f = 0 and g = 0 in (\mathbb{C}^2 , 0), study when there exists a deformation $f_t = 0$ where $f_0 = 0$ has the topological type of f = 0 and $f_t = 0$ the one of g = 0.
- OUR OBSERVATION: deformation fixing the free points of the generic curves are characterized by the valuative criterium.

Proposition

Let E and F be prime divisors over $O \in \mathbb{C}^2$. There exists a deformation f_s of a curve $f_0 = 0$ that lifts transversal to E that fixes the free points for F ($f_s = 0$ has strict transform transverse to F for $s \neq 0$) if and only if $ord_F \leq ord_E$.

Good things about the result and our problem:

- It talks about concrete divisors, not only topological types.
- Takes into account the contact order of *E* and *F*.
- They are very easy to check finite conditions (inequalities for h_D with D in the minimal model of F) => Algorithm!
- Also works for F a non-prime divisor: if $F = \sum_i a_i F_i$ then we the condition is $ord_F := \sum_i a_i ord_{F_i} \le ord_E$.

Bad news:

• Not all the adjacencies are of this type.

We recover many of the adjacencies from Arnol'd's list.





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Only 7 out of the 93 classical adjacencies between simple singularities of $\mu \leq 8$ are not realizable.

We recover many of the adjacencies from Arnol'd's list.

$$A_1 \longleftarrow A_2 \longleftarrow A_3 \longleftarrow A_4 \longleftarrow A_5 \longleftarrow A_6 \longleftarrow A_7 \longleftarrow A_8 \longleftarrow \dots$$

 $D_4 \longleftarrow D_5 \longleftarrow D_6 \longleftarrow D_7 \longleftarrow D_8 \longleftarrow \dots$
 $E_6 \longleftarrow E_7 \longleftarrow E_8$

• For example, $ord_{A_5} \leq ord_{E_6}$ but still there exists a deformation

$$y^3 + x^4 + s^2 y^2 + 2s x^2 y.$$

We recover many of the adjacencies from Arnol'd's list.



Some were not in Arnol'd's list:

 $Z_{11} = S_{2,4,5} \rightarrow E_8, \ Z_{12} = S_{2,4,6} \rightarrow J_{10} = T_{2,3,6}, \ W_{17} \rightarrow Z_{13} = S_{2,4,7}$ Some are not realizable:

$$W_{18} \nrightarrow Z_{17}, Z_{11} \nrightarrow J_{10}, X_9 \nrightarrow E_7.$$

Relation to the study of δ constant stratum.

- Recall Teissier's Theorem: a deformation f_t admits a parametrization in family if and only if it is δ -constant. $(\delta(C, 0) = \dim_{\mathbb{C}}(\mathcal{O}_{\bar{C}, \bar{0}}/\mathcal{O}_{C, 0})).$
- Describe all the $\overline{N}_E \subset \overline{N}_F$ is equivalent to describe which of the deformations fixing the free points are in the δ -constant stratum.
- Our problem is slightly different to the classical study of the δ -constant stratum: may be easier?

Happy birthday and thank you!

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Take combinatorial information (1/0) about n-3 edges (straight/curve) and n-3 vertices (broken between straight/smooth).

Combinatorics induces a partial order: the more straight lines and broken vertices, the bigger.



You get a duality that inverting the partial order just interchanging broken/curve and smooth/straight and reading backwards.



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It is just a combinatorial happening for the moment, will it appear in a deeper context?

 P. Popescu-Pampu, M. Pe Pereira, Fibonacci numbers and self-dual lattice structures for plane branches.Bridging Algebra, Geometry, and Topology, Springer Proceedings in Mathematics Statistics, 96