# New Perspectives for Multi-Armed Bandits and Their Applications

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CMLA, ENS Paris-Saclay

## Motivations & Objectives

- Size of data: n patients with some proba of getting cured
- Choose one of two treatments to prescribe





Patients cured or dead

1) Inference: Find the best treatment between the red and blue
2) Cumul: Save as many patients as possible

2) **Cumul:** Save as many patients as possible

- Size of data: *n* banners with some proba of click
- Choose one of two ads to display





- Banner clicked or ignored

1) Inference: Find the best ad between the red and blue

2) Cumul: Get as many clicks as possible

- Size of data: *n* auctions with some expected revenue
- Choose one of two strategies(bid/opt out) to follow



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Auction won or lost

- 1) Inference: Find the best strategy between the red and blue
- 2) Cumul: Win as many profitable auctions as possible

- Size of data: *n* mails with some proba of spam
- Choose one of two actions: spam or ham





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Mail correctly or incorrectly classified

- 1) Inference: Find the best strategy between the red and blue
- 2) Cumul: Minimize number of errors as possible

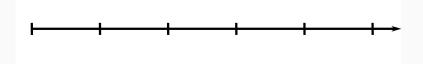
- Size of data: n patients with some proba of getting cured
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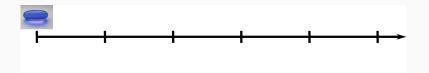


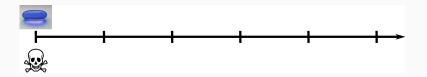
Patients cured ♥ or dead ♀

1) Inference: Find the best treatment between the red and blue

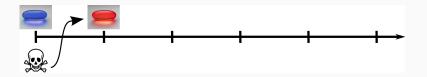
2) **Cumul:** Save as many patients as possible



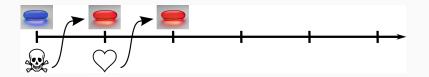


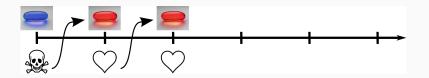


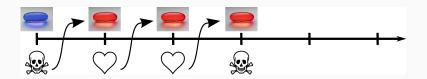


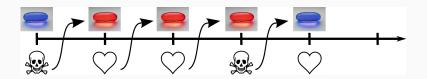


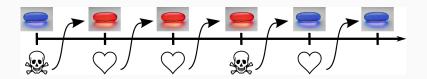


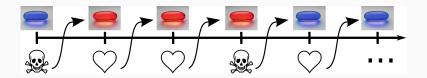


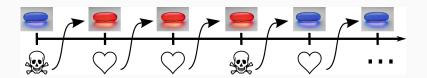












- Patients arrive and are treated sequentially.
- Save as many as possible.



Stochastic Multi-Armed Bandit

#### K-Armed Stochastic Bandit Problems

– K actions  $i \in \{1, ..., K\}$ , outcome  $X_t^i \in \mathbb{R}$  (sub-)Gaussian, bounded

$$X_1^i, X_2^i, \ldots, \sim \mathcal{N}(\mu^i, 1)$$
 i.i.d.

- Non-Anticipative Policy:  $\pi_t(X_1^{\pi_1}, X_2^{\pi_2}, \dots, X_{t-1}^{\pi_{t-1}}) \in \{1, \dots, K\}$
- Goal: Maximize expected reward  $\sum_{t=1}^{T} \mathbb{E} X_t^{\pi_t} = \sum_{t=1}^{T} \mu^{\pi_t}$
- Performance: Cumulative Regret

$$R_T = \max_{i \in \{1,2\}} \sum_{t=1}^T \mu^i - \sum_{t=1}^T \mu^{\pi_t} = \Delta_i \sum_{t=1}^T \mathbb{1} \{ \pi_t = i \neq \star \}$$

with  $\Delta_i = \mu^* - \mu^i$ , the "gap" or cost of error i.

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#### Most Famous algorithm [Auer, Cesa-Bianchi, Fisher, '02]

• UCB - "Upper Confidence Bound"

$$\pi_{t+1} = \arg\max_{i} \left\{ \overline{X}_{t}^{i} + \sqrt{\frac{2\log(t)}{T^{i}(t)}} \right\},$$

where 
$$T^i(t) = \sum_{t=1}^t \mathbb{1}\{\pi_t = i\}$$
 and  $\overline{X}_t^i = \frac{1}{T_t^i} \sum_{s:i_s=i} X_s^i$ .

#### Regret:

$$\mathbb{E} R_T \lesssim \sum_k \frac{\log(T)}{\Delta_k}$$

#### Worst-Case:

$$\mathbb{E} R_T \lesssim \sup_{\Delta} \kappa \frac{\log(T)}{\Delta} \wedge T\Delta$$
$$\approx \sqrt{\kappa T \log(T)}$$

## Ideas of proof $\pi_{t+1} = rg \max_i \left\{ \overline{X}_t^i + \sqrt{\frac{2 \log(t)}{T^i(t)}} \right\}$

· 2-lines proof:

$$\pi_{t+1} = i \neq \star \iff \overline{X}_t^{\star} + \sqrt{\frac{2\log(t)}{T^{\star}(t)}} \leq \overline{X}_t^i + \sqrt{\frac{2\log(t)}{T^i(t)}}$$

$$"\implies \Delta_i \leq \sqrt{\frac{2\log(t)}{T^i(t)}} \implies T^i(t) \lesssim \frac{\log(t)}{\Delta_i^2}$$

• Number of mistakes grows as  $\frac{\log(t)}{\Delta_i^2}$ ; each mistake costs  $\Delta_i$ .

Regret at stage T 
$$\lesssim \sum_i \frac{\log(7)}{\Delta_i^2} \times \Delta_i \approx \sum_i \frac{\log(7)}{\Delta_i}$$

- " $\Longrightarrow$ " actually happens with overwhelming proba
- "optimal": no algo can always have a regret smaller than  $\sum_i \frac{\log(T)}{\Delta_i}$

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#### Other Algos

 Other algo, ETC [Perchet, Rigollet], pulls in round robin then eliminates

$$R_T \lesssim \sum_k \frac{\log(T\Delta^k)}{\Delta^k}$$
, worst case  $R_T \leq \sqrt{T\log(K)K}$ 

Other algo, MOSS [Audibert, Bubeck], variants of UCB

$$R_T \lesssim K \frac{\log(T\Delta^{\min}/K)}{\Delta^{\min}}$$
, worst case  $R_T \leq \sqrt{TK}$ 

• Infinite number of actions  $x \in [0,1]^d$  with  $\Delta(x)$  1 Lipschitz. Discretize + UCB gives

$$R_T \lesssim T\varepsilon + \sqrt{\frac{T}{\varepsilon}} \leq T^{2/3}$$

Very interesting....

useful?

no...

Here is a list of reasons

#### On the basic assumptions

- Stochastic: Data are not iid, patients are different ill-posedness, feature selection/model selection
- 2. Different Timing: several actions for one reward pomdp, learn trade bias/variance
- Delays: Rewards not received instantaneously grouping, evaluations
- Combinatorial: Several decisions at each stage combinatorial optimization, cascading
- 5. Non-linearity: concave gain, diminishing returns, etc

Investigating (past/present/futur) them

#### Patients are different

- · We assumed (implicitly?) that all patients/users are identical
- Treatments efficiency 9proba of clicks) depend on age, gender...
- Those covariates or contexts are observed/known before taking the decision of blue/red pill

The decision (and regret...) should ultimately depend on it

#### General Model of Contextual Bandits

- Covariates:  $\omega_t \in \Omega = [0, 1]^d$ , i.i.d., law  $\mu$  (equivalent to)  $\lambda$ The cookies of a user, the medical history, etc.
- Decisions:  $\pi_t \in \{1,..,K\}$ The decision can (should) depend on the context  $\omega_t$
- Reward:  $X_t^k \in [0,1] \sim \nu^k(\omega_t)$ ,  $\mathbb{E}[X^k|\omega] = \mu^k(\omega)$ The expected reward of action k depend on the context  $\omega$
- Objectives: Find the best decision given the request Minimize regret  $R_T := \sum_{t=1}^T \mu^{\pi^*(\omega_t)}(\omega_t) \mu^{\pi_t}(\omega_t)$

#### Regularity assumptions

1. Smoothness of the pb: Every  $\mu^k$  is  $\beta$ -hölder, with  $\beta \in (0,1]$ :

$$\exists L > 0, \ \forall \omega, \omega' \in \mathcal{X}, \ \|\mu(\omega) - \mu(\omega')\| \le L\|\omega - \omega'\|^{\beta}$$

2. Complexity of the pb: ( $\alpha$ -margin condition)  $\exists C_0 > 0$ ,

$$\mathbb{P}_{X}\left[0<\left|\mu^{1}(\omega)-\mu^{2}(\omega)\right|<\delta\right] \,\leq\, C_{0}\delta^{\alpha}$$

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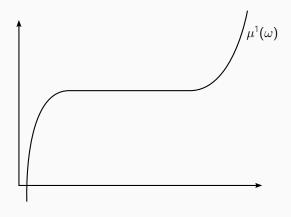
2. Complexity of the pb: ( $\alpha$ -margin condition)  $\exists C_0 > 0$ ,

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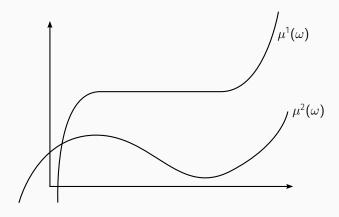
where  $\mu^*(\omega) = \max_k \mu^k(\omega)$  is the maximal  $\mu^k$  and  $\mu^{\sharp}(\omega) = \max \left\{ \mu^k(\omega) \text{ s.t. } \mu^k(\omega) < \mu^*(\omega) \right\}$  is the second max.

With K > 2:  $\mu^*$  is  $\beta$ -Hölder but  $\mu^{\sharp}$  is not continuous.

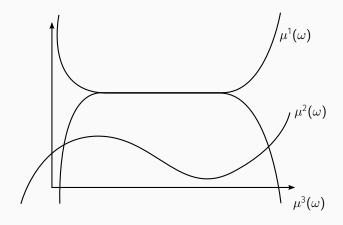
### Regularity: an easy example ( $\alpha$ big)



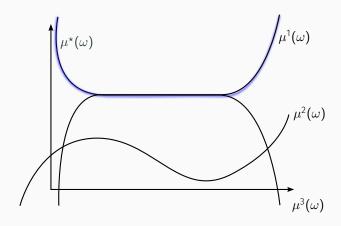
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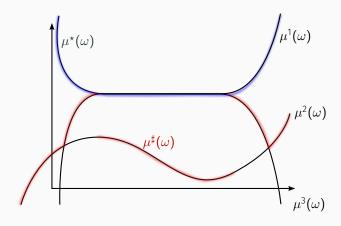
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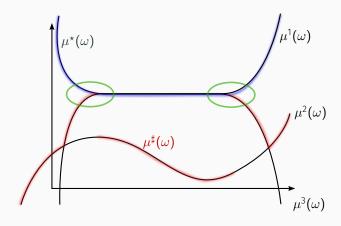
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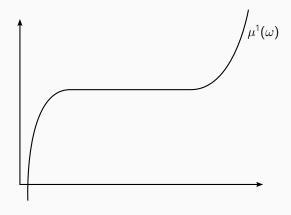


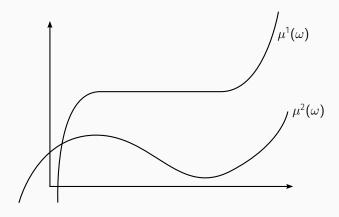
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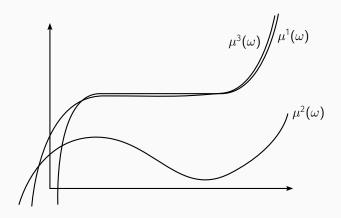


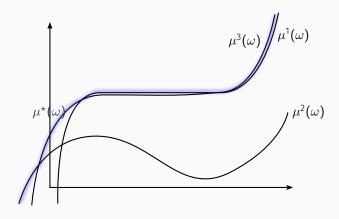
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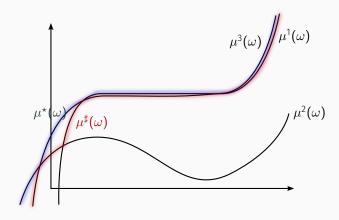


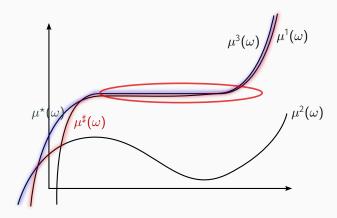




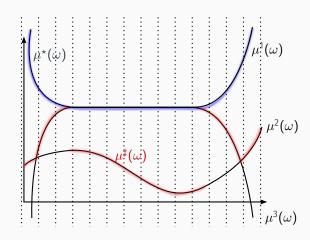




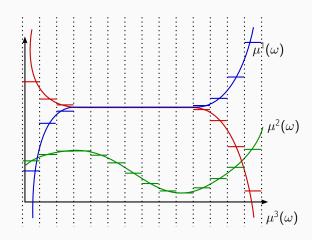




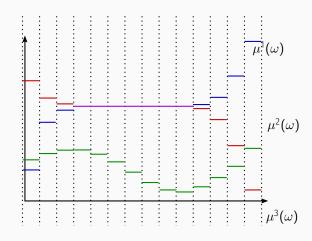
## Binned policy



# Binned policy



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#### Binned Successive Elimination (BSE)

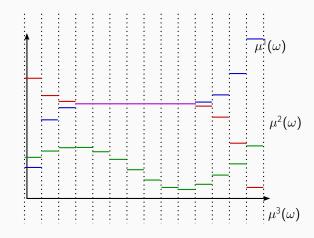
#### Theorem [P. and Rigollet ('13)]

If 
$$\alpha < 1$$
,  $\mathbb{E}[R_T(\mathrm{BSE})] \lesssim T\left(\frac{K\log(K)}{T}\right)^{\frac{\beta(1+\alpha)}{2\beta+d}}$ , bin side  $\left(\frac{K\log(K)}{T}\right)^{\frac{1}{2\beta+d}}$ .

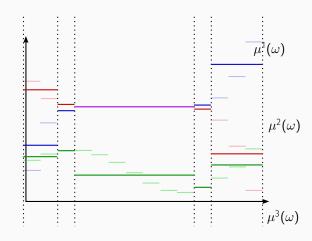
For K = 2, matches lower bound: minimax optimal w.r.t. T.

- Same bound with full monit [Audibert and Tsybakov, '07]
- No log(*T*): difficulty of nonparametric estimation washes away the effects of exploration/exploitation.
- $\alpha$  < 1: cannot attain fast rates for easy problems.
- Adaptive partitioning!

# Suboptimality of (BSE) for $\alpha \geq 1$



# Suboptimality of (BSE) for $\alpha \geq 1$



#### Adaptive BSE (ABSE)

#### Theorem [P. and Rigollet ('13)]

For all 
$$\alpha$$
,  $\mathbb{E}[R_T(\mathrm{ABSE})] \lesssim T\left(\frac{K\log(K)}{T}\right)^{\frac{\beta(1+\alpha)}{2\beta+d}}$ .

For K = 2, matches lower bound: minimax optimal w.r.t. T.

• Same bound than (BSE) even for easy problems  $\alpha \geq$  1.

#### This is not the solution

1. **dimensions** dependent bound:  $T^{1-\frac{\beta}{2\beta+d}}$ 

 $d=+\infty$  and  $\beta=0$ , lots of contexts, no regularity Online selection of models ? Ill-posed pb  $\mu(\cdot)$  not  $\beta$ -holder

#### Estimation/Approx errors

Performance = Approx Error + Regret( $\beta$ , d, T)

2. Non-stationarity of arms: Value are not i.i.d., evolve with time. Ex. ads for movies.

Cumulative objectives clearly not the solution.

Discount? How, why, at which speeds?

3. Non-stationarity of sets of arms:

Arms arrive and disappears

How incorporate a new arm? which index?

#### This was really not the solution

1. Non-stationarity of **sets** of arms:

Arms arrive and disappears

How incorporate a new arm? which index?

2. Contexts (covariates) are not in  $\mathbb{R}^d$ 

Rather descriptions, texts, id, images...How to embed? training set is influenced by algorithms...



#### **Example of Repeated Auctions**



#### Ad slot sold by lemonde.fr. 2nd-price auctions

- · Several (marketing) companies places bids
- · Highest bid wins (...), say criteo, pays to lemonde 2nd bid (...)
- criteo chooses ad of a client, fnac or singapore airlines
- criteo paid by the client if the user clicks on the ad

Main Problem: Repeated auctions with unknown private valuation Learn valuations, find which ad to display & good strategies

#### Repeated auctions

- 1. Can be modeled as a bandit pb with Extra Structure
- Actually, Criteo (Google, Facebook) paid if the user buys something after the click

Needs several "costly" auctions to seal a deal

Auctions lost can also help to seal deal (competitor displays ad for free)

Optimal strategy in repeated auctions, learn it? (POMDP?)

Reward timing per user, decision timing by opportunities

#### Other examples - repeated A/B tests

 Companies test new technologies (algo, hardware, etc.) before putting in productions. Sequences of AB tests

Timing of Decisions: each day, continue, stop or validate the current AB test

Timing of Rewards: Total improvements of implemented techno.

 The longer AB test are, the more confident (reduces variance) but less and less implementation

Online tradeoff risks/performances



#### Rewards are not observed immediately

- Clinical trials: have to wait 6 months to see results.
  - A trial length is 3 year : 6 phases Regret is still  $\sqrt{T}$
- Marketing (ad displays), only see if users buy
  - No feedback is either no sale (forever) or no sale **yet** 
    - Build estimators with censured/missing data
      - Feasible with iid data... but they are not!

# **Combinatorial Structure**

#### Large Decision spaces



- · Choose not to display 1 ad, but 4, 6, 10...
- Paid if sales after click (even if unrelated)
  - Lots of correlations (between products, positions, colors/style of banner, **time**, etc.)

Some products are seen, other are not (carrousels...)

Too many possibilities of (almost) equal performances

Compete with the best 
$$R_T \leq \sqrt{KT}$$
  
but at least top 5%,  $R_T \leq \sqrt{\log(K)\frac{1}{5\%}T}$ ??

## Bandit theory is quite neat

To be "applied", or relevant, need LOTS of work

Anybody is welcome to join & collaborate!

Model selection, Feature extractions, Missing Data, Censured Data,

Combinatorial Optimization, New techniques estimators..