

Control theory and splitting methods

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Goal of this talk

- ▶ **highlight the deep links between numerical splitting methods and control theory**
- ▶ provide (control-inspired) proofs of conjectures of the order theory of splitting methods

Control theory: the infinitesimal question

Let $m \in \mathbb{N}^*$ and $f_0, \dots, f_m : \mathbb{R}^d \rightarrow \mathbb{R}^d$ smooth such that $f_0(0) = 0$.

$$\dot{x}(t) = f_0(x(t)) + u_1(t)f_1(x(t)) + \dots + u_m(t)f_m(x(t))$$

(Small-state) Small-time local controllability (STLC):

$\forall T, \delta > 0, \exists r > 0, \forall x^* \in B_{\mathbb{R}^d}(0, r), \exists u \in L^1((0, T); \mathbb{R}^m)$
such that $x(T; u, 0) = x^*$ and $x([0, T]) \subset B(0, \delta)$.

- For linear systems $\dot{x} = Ax + Bu$: NSC = Kalman rank condition. [1960]
- If $f_0 = 0$, NSC = LARC: $\text{Lie}(f_1, \dots, f_m)(0) = \mathbb{R}^d$. [Chow, Rashevski 1930']
- If $f_0 \neq 0$, the iterated Lie brackets at 0 still have the answer, [Krener1973]
 $\text{Lie}(f_0, \dots, f_m)(0) = \mathbb{R}^d$ is necessary, [Herman 1963, Nagano 1966]
but not sufficient: bad Lie brackets $\text{ad}_{f_1}^{2k}(f_0)(0)$. [Sussmann 1983, Stefani 1986]



Rudolf Kalman
1930-2016



Sophus Lie
1842-1899

Splitting methods

We want to compute the flow of $f_0 + f_1$ from the flows of f_0 and f_1 .

Splitting method of order N : (α, β) such that $\forall f_0, f_1, \forall x_0$

$$e^{T(f_0+f_1)}x_0 = e^{\alpha_1 T f_0} e^{\beta_1 T f_1} \dots e^{\alpha_k T f_0} e^{\beta_k T f_1} x_0 + O_{T \rightarrow 0}(T^{N+1}). \quad (1)$$

Ex: Lie - Trotter : $e^{T(f_0+f_1)}x_0 = e^{T f_0} e^{T f_1} x_0 + O(T^2)$

Strang : $e^{T(f_0+f_1)}x_0 = e^{\frac{1}{2}T f_0} e^{T f_1} e^{\frac{1}{2}T f_0} x_0 + O(T^3)$

- Without constraint i.e. $(\alpha, \beta) \in \mathbb{R}^k \times \mathbb{R}^k$: existence $\forall N$.
- With only forward flows of f_0 i.e. $(\alpha, \beta) \in \mathbb{R}_+^k \times \mathbb{R}^k$: order ≤ 2 .
[Goldman Kaper 1996, Blanes Casas 2005]
- If $(\alpha, \beta) \in \mathbb{R}_+^k \times \mathbb{C}_+^k$ order 6 OK, if $(\alpha, \beta) \in \mathbb{C}_+^k \times \mathbb{C}_+^k$ order 44 OK
[Castella Chartier Descombes Vilmart 2009, Hansen Ostermann 2009, Blanes Casas Chartier Murua 2016]
- With additional commutator flows: with $[f_1, [f_1, f_0]]$ order 4 OK
[Takahashi Imada 1984, Auzinger Hofstatter Koch 2019]



Sophus Lie
1842-1899



Hale Trotter
1931-2022



Gilbert Strang
1934-

Splitting methods with $+f_0$ as control systems

- An $(\mathbb{R}_+, \mathbb{K})$ -splitting method $e^{\alpha_k T f_0} e^{\beta_k T f_1} \dots e^{\alpha_1 T f_0} e^{\beta_1 T f_1}$
(i.e. with $(\alpha, \beta) \in \mathbb{R}_+^k \times \mathbb{K}^k$) is a trajectory of the control system

$$\dot{x}(t) = f_0(x(t)) + u(t)f_1(x(t))$$

where $u = \sum_j \beta_j T \delta_{t=\tau_j}$, $\beta_j \in \mathbb{K}$, $\tau_j = \sum_{j' \leq j} \alpha_{j'} T$.

Ex: Lie - Trotter $e^{T f_0} e^{T f_1} \longrightarrow u(t) = \delta_0(t)$
Strang $e^{\frac{1}{2} T f_0} e^{T f_1} e^{\frac{1}{2} T f_0} \longrightarrow u(t) = \delta_{T/2}(t)$

- A splitting method involving $f_2 := [f_1, [f_1, f_0]]$

$$e^{\alpha_k T f_0} e^{\beta_k T f_1} e^{\alpha'_k T f_0} e^{\gamma_k T f_2} \dots e^{\alpha_1 T f_0} e^{\beta_1 T f_1} e^{\alpha'_1 T f_0} e^{\gamma_1 T f_2}$$

(where $(\alpha, \beta, \gamma) \in \mathbb{R}_+^{2k} \times \mathbb{R}^k \times \mathbb{R}^k$) is a trajectory of the control system

$$\dot{x}(t) = f_0(x(t)) + u_1(t)f_1(x(t)) + u_2(t)f_2(x(t))$$

where u_1, u_2 are finite sums of Dirac masses with disjoint supports.



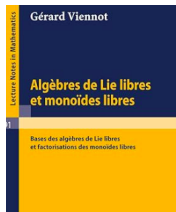
Paul Dirac
1902-1984

Common tools: (1) Free algebras

- ▶ Let $X := \{X_0, X_1\}$ be non-commutative **unknowns**.
- ▶ Let $\mathcal{A}(X)$ be the **free algebra** over X , i.e. the vector space of non-commutative polynomials. e.g. $7X_0^2 + 3X_1X_0 + 2X_0X_1$
- ▶ Let $\mathcal{L}(X)$ the **free Lie algebra** over X , i.e. the smallest vector subspace of $\mathcal{A}(X)$ containing X_0, X_1 , and stable by the Lie bracket (commutator) operation $[a, b] := ab - ba$.
- ▶ Let $\widehat{\mathcal{A}}(X)$ be the algebra of **formal power series**. e.g. $\sum_{k=0}^{\infty} k! X_0^k X_1$
$$\mathcal{L}(X) \subset \mathcal{A}(X) \subset \widehat{\mathcal{A}}(X)$$
- ▶ One can “**evaluate**” (although not injective)

$$b \in \mathcal{L}(X) \hookrightarrow f_b \in C^\omega(\mathbb{R}^d; \mathbb{R}^d) \hookrightarrow f_b(0) \in \mathbb{R}^d$$

$$\begin{aligned} \text{e.g. } [X_1, X_0] &= X_1X_0 - X_0X_1 \rightarrow [f_1, f_0] = (Df_0)f_1 - (Df_1)f_0 \\ &\rightarrow [f_1, f_0](0) =: f_{(X_1, X_0)}(0) \end{aligned}$$



C. Reutenauer

Common tools: (2) Formal linear differential equation

Let $u \in L^1(0, T)$.

$$\begin{cases} \dot{S}(t) = S(t)(X_0 + u(t)X_1 + \sum_{b \in A} u_b(t)b) \\ S(0) = 1 \end{cases}$$

Definition

The solution is the formal-series valued function $S : [0, T] \rightarrow \hat{\mathcal{A}}(X)$ whose components S_n of degree n satisfy $S_0(t) = 1$

$$S_{n+1}(t) = \int_0^t \left(S_n(\tau)(X_0 + u(\tau)X_1) + \sum_{b \in A} u_b(\tau)S_{n+1-|b|}(\tau)b \right) d\tau.$$

Chen-Fliess expansion: $S(t) = \sum_{n \in \mathbb{N}} \sum_{\sigma \in \{0,1\}^n} c_\sigma(t, u) X_{\sigma_1} \cdots X_{\sigma_n}$.
If $u = \beta \delta_\tau$, take the limit $\epsilon \rightarrow 0$ on the solutions for $u = \frac{\beta}{\epsilon} 1_{(\tau-\epsilon, \tau+\epsilon)}$.



Kuo Tsai Chen
1923-1987



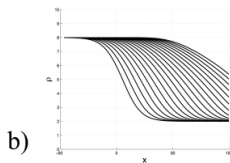
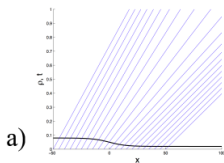
Michel Fliess
1945-

Linearization principle (converse characteristics method)

One transforms the **nonlinear ODE in \mathbb{R}^d** : $\dot{x} = f_0(x) + u(t)f_1(x)$, into a **linear equation in $\text{Op}(C^\infty(\mathbb{R}^d, \mathbb{R}))$** : $\frac{d}{dt}L(t) = L(t)\underbrace{(f_0 \cdot \nabla)}_{X_0} + u(t)\underbrace{f_1 \cdot \nabla}_{X_1}$ on which we use the algebra developed for the formal differential equation $\dot{S}(t) = S(t)(X_0 + u(t)X_1)$.

The zero-order operator $\left| \begin{array}{l} L(t) : C^\infty(\mathbb{R}^d, \mathbb{R}) \rightarrow C^\infty(\mathbb{R}^d, \mathbb{R}) \\ \varphi \mapsto (p \mapsto \varphi(x(t; u, p))) \end{array} \right.$
satisfies $\forall \varphi \in C^\infty(\mathbb{R}^d, \mathbb{R}), \forall p \in \mathbb{R}^d$,

$$\begin{aligned} \frac{d}{dt}(L(t)\varphi)(p) &= D\varphi(x(t; f, p)) \cdot (f_0 + u(t)f_1)(x(t; f, p)) \\ &= \left(L(t)(f_0 \cdot \nabla + u(t)f_1 \cdot \nabla)\varphi \right)(p) \end{aligned}$$



Truncation and approximation

$$\begin{cases} \dot{S}(t) = S(t)(X_0 + u(t)X_1 + \sum_{b \in A} u_b(t)b) & \text{in } \widehat{\mathcal{A}}(X) \\ S(0) = 1 \end{cases}$$

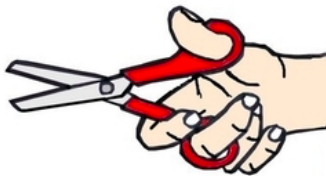
If $S_N(t; X, u)$ is the component of $S(t)$ on polynomials of degree $\leq N$ then, $\forall f_0, f_1, u, x_0$, the solution of

$$\begin{cases} \dot{x}(t) = f_0(x(t)) + u(t)f_1(x(t)) + \sum_{b \in A} u_b(t)f_b(x(t)) \\ x(0) = x_0 \end{cases}$$

satisfies

$$x(1; Tf, u, x_0) = S_N(1; Tf, u)(Id)(x_0) + \mathcal{O}_{T \rightarrow 0}(T^{N+1})$$

Splitting has to do with the reachability of $e^{X_0+X_1}$ for $S_N(t; X, u)$.



Splitting method = free trajectory reaching $e^{X_0+X_1}$

$\mathcal{A}^N(X)$: Free nilpotent algebra of index N {polynom. of deg. $\leq N$ }

$\mathcal{L}^N(X)$: Free nilpotent Lie algebra of index N $\mathcal{L}(X) \cap \mathcal{A}^N(x)$

$G^N(X) = \exp(\mathcal{L}^N(X))$: Lie group analytic submanifold

$$\begin{cases} \dot{S}_N(t) = S_N(t)(X_0 + u(t)X_1 + \sum_{b \in A} u_b(t)b) & \text{in } G^N(X) \\ S_N(0) = 1 \end{cases} \quad (2)$$

Splitting method of order N involving the elements of A

\Leftrightarrow **Trajectory of (2) with controls $u, u_b =$ finite sums of Dirac masses** with disjoint supports **such that** $S_N(1; X, u) = e^{X_0+X_1}$

\Leftarrow **Controllability of the system (2)**

$$u(t) = a\underline{u}\left(\frac{t}{T}\right) \Rightarrow S_N(t; u) = \lambda S_N\left(\frac{t}{T}, \underline{u}\right) \text{ where } \left| \begin{array}{l} \lambda: \mathcal{A}^N(X) \rightarrow \mathcal{A}^N(X) \\ X_0 \mapsto TX_0 \\ X_1 \mapsto aX_1 \end{array} \right.$$

$$e^A e^B = e^Z \text{ where } Z = A + B + \frac{1}{2}[A, B] + \frac{1}{12}([A, [A, B]] + [B, [B, A]]) + \dots$$



Henry F. Baker
1866-1956



John E. Campbell
1862-1924



Felix Hausdorff
1868-1942



Wilhelm Magnus
1907-1990

The simplest case

Theorem

For every $N \in \mathbb{N}^*$, there exists an (\mathbb{R}, \mathbb{R}) -splitting method of order N .

Theorem (Chow 1939, Rashevski 1938)

If $f_0, f_1 : \mathbb{R}^d \rightarrow \mathbb{R}^d$ are smooth vector fields s.t. $\text{Lie}(f_0, f_1)(x_0) = \mathbb{R}^d$ then the system $\dot{x}(t) = u_0(t)f_0(x(t)) + u_1(t)f_1(x(t))$ is STLC at x_0 .

$$\begin{cases} \dot{S}_N(t) = S_N(t)(u_0(t)X_0 + u_1(t)X_1) & \text{in } G^N(X) \\ S_N(0) = 1 \end{cases}$$

For $f_0(S) = SX_0$ and $f_1(S) = SX_1$ we have $f_b(S) = Sb$ thus $\text{Lie}(f_0, f_1)(1) = \mathcal{L}^N(X)$.



Wei Liang Chow
1911-1995



Petr Rashevskii
1907-1983

Complex systems with a complex control $u : (0, T) \rightarrow \mathbb{C}$

Theorem (B, Laurent, Marbach 2024)

For every $N \in \mathbb{N}^*$, there exists an $(\mathbb{R}_+, \mathbb{C})$ -splitting method of order N .

Theorem (B, Laurent, Marbach 2024, (Sussmann 1987))

If $f_0, f_1 : \mathbb{C}^d \rightarrow \mathbb{C}^d$ are holomorphic, $f_0(0) = 0$ and $\text{Lie}(f_0, f_1)(0) = \mathbb{C}^d$ then the system $\dot{x}(t) = f_0(x(t)) + u(t)f_1(x(t))$ is STLC.

$$\begin{cases} \dot{S}_N(t) = S_N(t)(X_0 + u(t)X_1) & \text{on } G^N(X) \\ S_N(0) = 1 \end{cases}$$

We build a trajectory from 1 to e^{TX_0} with a controllable linearized system
[\[Coron's return method\]](#)

- LARC provides u with a controllable linearized system.
- Then replacing u by $u \diamond e^{i\frac{2\pi}{N}} u \diamond \dots \diamond e^{i\frac{2(N-1)\pi}{N}} u : (0, NT) \rightarrow \mathbb{C}$ increases the valuation of the polynomial $e^{-TX_0} S_N(T; X, u) - 1$. And iterate.

Why does the valuation of \mathcal{Z}_N increase?

$$\begin{cases} \dot{S}_N(t) = S_N(t)(X_0 + u(t)X_1) & \text{in } G^N(X) \\ S_N(0) = 1 \end{cases}$$

$$S_N(t, u) = e^{tX_0} e^{\mathcal{Z}_N(t, u)} \text{ where } \mathcal{Z}_N(t, u) \in \mathcal{L}_0^N(X)$$

- $\mathcal{Z}_N(T_1 + T_2, u_1 \diamond u_2) = \mathcal{Z}_N(T_1, u_1) + \mathcal{Z}_N(T_2, u_2) + \text{higher valuation}$

$$e^{(T_1+T_2)X_0} e^{\mathcal{Z}(T_1+T_2, u_1 \diamond u_2)} = e^{T_2 X_0} e^{\mathcal{Z}(T_2, u_2)} e^{T_1 X_0} e^{\mathcal{Z}(T_1, u_1)}$$

$$\text{thus } e^{\mathcal{Z}(T_1+T_2, u_1 \diamond u_2)} = e^{-T_1 X_0} e^{\mathcal{Z}(T_2, u_2)} e^{T_1 X_0} e^{\mathcal{Z}(T_1, u_1)} + \text{CBH}$$

- $\mathcal{Z}_N(T, u) = \sum_{j=1}^{N-1} \underbrace{Z_j}_{\deg_{X_1}=j} \Rightarrow \mathcal{Z}_N(T, e^{i\theta} u) = \sum_{j=1}^{N-1} e^{ij\theta} Z_j.$

Finally with $v := u \diamond e^{i\frac{2\pi}{N}} u \diamond \dots \diamond e^{i\frac{2(N-1)\pi}{N}} u$ we obtain

$$\mathcal{Z}_N(NT, v) = \sum_{k=0}^{N-1} \sum_{j=1}^{N-1} e^{i\frac{2k\pi}{N} j} Z_j + \text{higher valuation}$$

Obstructions to splitting methods of order N

- The maximal order of an $(\mathbb{R}_+, \mathbb{R})$ -splitting method is 2.
[Goldman-Kaper 1996, Blanes, Casas 2005]
- The max order of an $(\mathbb{R}_+, \mathbb{R})$ splitting method involving $\text{ad}_{X_1}^2(X_0)$ is 4
[Takahashi Imada 1984, Auzinger Hofstatter Koch 2019]

What happens next?

$\forall k \in \mathbb{N}^*$, $W_k := \text{ad}_{\text{ad}_{X_0}^{k-1}(X_1)}^2(X_0)$ is an obstruction to STLC [B. Marbach 2022]

Theorem (B, Laurent, Marbach 2024)

The maximal order of an $(\mathbb{R}_+, \mathbb{R})$ -splitting method involving W_1, \dots, W_{N-1} is $2N$. (W_{2N} is an obstruction to order $2N + 1$)

Sussmann's product, coordinates of the 2nd kind [1986]:

if \mathcal{B} is a Hall basis of $\mathcal{L}(X)$ then $S(t; X, u) = \overleftarrow{\prod}_{b \in \mathcal{B}} \exp(\xi_b(t, u)b)$



Marshall Hall Jr
1910-1990



Hector Sussmann
1946-

Easy computation of the $\xi_b(t, u)$: Lazard elimination

$$\begin{cases} \dot{S}(t) = S(t)(X_0 + u(t)X_1) \\ S(0) = 1 \end{cases} \quad \text{Goal : } S(t) = \prod_{b \in \mathcal{B}}^{\leftarrow} e^{\xi_b(t; u)b}$$

Step 1: Eliminate X_1 by looking for $S(t)$ of the form $S_1(t)e^{\xi(t)X_1}$.

$$\begin{aligned} \dot{S}_1(t) &= S(t)(X_0 + u(t)X_1)e^{-\xi_{X_1}(t)X_1} + S(t)(-\dot{\xi}(t)X_1)e^{-\xi(t)X_1} \\ &= S_1(t)e^{\xi(t)X_1}X_0e^{-\xi(t)X_1} \quad \text{because } \dot{\xi}(t) = u(t) \\ &= S_1(t) \underbrace{\sum_{m \in \mathbb{N}} \frac{\xi(t)^m}{m!} \cdot \dot{\xi}_{X_0}(t) \text{ad}_{X_1}^m(X_0)}_{=\dot{\xi}_{\text{ad}_{X_1}^m(X_0)}(t)} \rightarrow \text{does not involve } X_1 \end{aligned}$$

Step $(j+1)$: $\dot{S}_j(t) = \sum_{b \in H_j} \dot{\xi}_b(t)b$

Choose a bracket $a \in H_j$ and look for $S_j(t)$ of the form $S_{j+1}(t)e^{\xi_a(t)a}$.

$$\dot{S}_{j+1}(t) = S_{j+1}(t) \sum_{m \in \mathbb{N}} \sum_{b \in H_j \setminus \{a\}} \underbrace{\frac{(\xi_a(t))^m}{m!} \dot{\xi}_b(t) \text{ad}_a^m(b)}_{\dot{\xi}_{\text{ad}_a^m(b)}(t)}$$

The concept of "Hall basis" of $\mathcal{L}(X)$ is adapted to this algorithm.

[Viennot 1978, Sussmann 1986]

And reciprocally...

Theorem (B, Laurent, Marbach 2024)

There exists an $(\mathbb{R}_+, \mathbb{R})$ -splitting method of order $2N$ involving W_1, \dots, W_{N-1} .

- ▶ If $\text{ad}_{f_1}^2(f_0) = f_j$ then the controllability of

$$\dot{x} = f_0(x) + u_1(t)f_1(x) + \dots + u_m(t)f_m(x)$$

is equivalent to the controllability of

$$\dot{x} = f_0(x) + u_1(t)f_1(x) + \dots + u_m(t)f_m(x) + u_{m+1}(t)[f_1, f_0](x)$$

- ▶ Let $M_0 := X_1$ and $M_{j+1} = [M_j, X_0]$. Then $W_j = \text{ad}_{M_{j-1}}^2(X_0)$.

$$\begin{cases} \dot{S}_N(t) = S_N(t)(X_0 + \sum_{b \in A} u_b(t)b) & \text{in } G^N(X) \\ S_N(0) = 1 \end{cases}$$

$$A_0 := \{X_1, W_1, \dots, W_{N-1}\}$$

$$\Leftrightarrow A_1 = A_0 \cup \{M_1\} \quad \text{because } \text{ad}_{X_1}^2(X_0) = W_1 \in A_0$$

$$\Leftrightarrow A_2 = A_1 \cup \{M_2\} \quad \text{because } \text{ad}_{M_1}^2(X_0) = W_2 \in A_1, \quad \text{etc...}$$

$$A_N \supset \{X_1, M_1, \dots, M_{2N}\} \text{ and } \text{Lie}(M_1, \dots, M_{2N}) = \mathcal{L}_{\emptyset}^N(X)$$

Conclusion

- ▶ Splitting methods \sim trajectories of (free) control systems.
- ▶ Control theory has tools to prove obstructions and to prove abstract existence results.
- ▶ They formalise the order theory of splitting methods with sign constraints on the coefficients.
- ▶ Splitting methods with $(\mathbb{R}^+, \mathbb{C})$ coefficients exist for any order.
- ▶ There exists an $(\mathbb{R}_+, \mathbb{R})$ splitting method of order $2N$ involving W_1, \dots, W_{N-1} . (a 100% quadratic theory)

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Quiz

Who are they?

