

# Asymmetric data-driven interdiction problems with cost uncertainty: a distributionally robust optimization approach (*working paper*)

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# General interdiction problem

- A deterministic interdiction problem is a zero-sum game between a *leader* (attacker) and a *follower* (defender).
- The leader allocates limited interdiction resources to target components that are crucial to the follower's operations.
- In response, the follower observes the leader's actions and aims to maximize its own profit within the impacted environment.
- The leader is informed about the follower's objective and chooses attacks that minimize the follower's profit.

# General interdiction problem

## Deterministic interdiction (DI) problem

$$[\text{DI}]: \min_{x \in X} \max_{y \in Y(x)} c^T y, \quad (1)$$

where

$$X = \{x \in \{0, 1\}^m : Hx \leq h\} \text{ and } Y(x) = \{y \in \mathbb{R}_+^n : Fy + Lx \leq f\} \quad (2)$$

are, respectively, the leader's and the follower's feasible sets, and  $c$  is a deterministic profit vector.

# Example: min-cost flow interdiction [1]

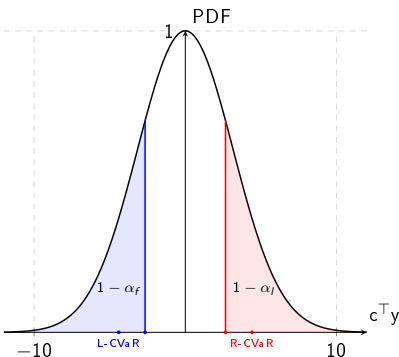
- Let  $G = (N, A)$  be a directed graph, where  $N$  and  $A$  are, respectively, the sets of nodes and directed arcs.
- Assume that the leader has a single budget constraint, i.e.,

$$X = \{x \in \{0, 1\}^{|A|} : \sum_{(i,j) \in A} r_{ij} x_{ij} \leq R\}.$$

- The follower aims to maximize  $-\sum_{(i,j) \in A} c_{ij} y_{ij}$  s.t.

$$y \in Y(x) = \left\{ y \in \mathbb{R}_+^{|A|} : \sum_{j: (i,j) \in A} y_{ij} - \sum_{j: (j,i) \in A} y_{sj} = d_i \quad \forall i \in N, \right. \\ \left. 0 \leq y_{ij} \leq u_{ij}(1 - x_{ij}) \right\}, \text{ where } \sum_{i \in N} d_i = 0.$$

# Stochastic interdiction problem



**Figure:** Illustration of the right-tail CVaR (red) and the left-tail CVaR (blue) for normal distribution with  $\alpha_l = \alpha_f = 0.8$  [2].

- We assume that  $c$  is a random vector governed by an unknown *true* distribution  $Q^*$ .
- If  $Q^*$  is given, then the decision-makers solve a stochastic interdiction (SI) problem of the form:

$$[\text{SI}]: \quad z^* := \min_{x, y} \rho_l(y, Q^*)$$

$$\text{s.t. } x \in X$$

$$y \in \operatorname{argmax}_{\tilde{y} \in Y(x)} \rho_f(\tilde{y}, Q^*),$$

where

$$\rho_l(y, Q^*) := \min_{t_l \in \mathbb{R}} \left\{ t_l + \frac{1}{1 - \alpha_l} \mathbb{E}_{Q^*} \{ (c^\top y - t_l)^+ \} \right\} \text{ and}$$

$$\rho_f(y, Q^*) := \max_{t_f \in \mathbb{R}} \left\{ t_f + \frac{1}{1 - \alpha_f} \mathbb{E}_{Q^*} \{ (c^\top y - t_f)^- \} \right\}.$$

# Distributionally robust interdiction model

- A1.** The feasible sets  $X$  and  $Y(x)$  for each  $x \in X$  are non-empty and bounded.
- A2.** The *support* of  $\mathbb{Q}^*$  is known and given by a non-empty compact polyhedral set

$$S = \{c \in \mathbb{R}^n : Bc \leq b\}.$$

- A3.** The leader and the follower have access to two i.i.d. training data sets, generated according to  $\mathbb{Q}^*$ ,

$$\hat{C}_I = \{\hat{c}_I^{(k)} \in S, k \in \{1, \dots, k_I\}\} \text{ and } \hat{C}_f = \{\hat{c}_f^{(k)} \in S, k \in \{1, \dots, k_f\}\}.$$

Furthermore,  $\hat{C}_f \subseteq \hat{C}_I$  or, equivalently,  $k_f \leq k_I$  and  $\hat{c}_f^{(k)} = \hat{c}_I^{(k)}$  for each  $k \in \{1, \dots, k_f\}$ .

# Distributionally robust interdiction model

## Basic DRI model

$$\begin{aligned} \text{[DRI]:} \quad \hat{z}_b^* &:= \min_{x,y} \left\{ \max_{Q_I \in \mathcal{Q}_I} \rho_I(y, Q_I) \right\} \\ \text{s.t. } x &\in X \\ y &\in \operatorname{argmax}_{\tilde{y} \in Y(x)} \left\{ \min_{Q_f \in \mathcal{Q}_f} \rho_f(\tilde{y}, Q_f) \right\}, \end{aligned}$$

where

$$\mathcal{Q}_i := \left\{ Q \in \mathcal{Q}_0(S) : W^P(\hat{Q}_i, Q) \leq \varepsilon_i \right\} \quad i \in \{I, f\}$$

are the Wasserstein ambiguity sets,  $\hat{Q}_i(\hat{C}_i)$  is an empirical distribution of the data and  $W^P(\hat{Q}_i, Q)$  is the type-1 Wasserstein distance w.r.t.  $\ell_p$ -norm.

# Basic DRI model

- Notably, **[DRI]** is strongly *NP*-hard, as even its deterministic version, **[DI]**, is known to be strongly *NP*-hard.
- If  $\mathcal{Q}_I$  and  $\mathcal{Q}_f$  are defined in terms of  $\ell_1$  or  $\ell_\infty$  norm, then **[DRI]** admits a single-level MILP reformulation of polynomial size. This reformulation builds on LP reformulations of the worst-case CVaR problems for both decision-makers [3] and a strong duality-based reformulation of the follower's problem:

$$\begin{aligned} \hat{z}_b^* = & \min_{x, y, \nu, s, \mu, \beta, \gamma, \lambda, t} \left\{ t_I + \frac{1}{1 - \alpha_I} (\lambda_I \varepsilon_I + \frac{1}{k_I} \sum_{k \in K_I} s_I^{(k)}) \right\} \\ \text{s.t. } & \text{primal and dual feasibility (follower),} \\ & x \in X \\ & \left. \begin{aligned} \hat{c}_I^{(k)\top} y - t_I + \Delta_I^{(k)\top} \nu_I^{(k)} &\leq s_I^{(k)} \\ \|B^\top \nu_I^{(k)} - y\|_* &\leq \lambda_I \\ \nu_I^{(k)} \geq 0, s_I^{(k)} &\geq 0 \end{aligned} \right\} \forall k \in K_I \\ & (-Lx + f)^\top \beta_f = t_f - \frac{1}{1 - \alpha_f} (\varepsilon_f \lambda_f + \frac{1}{k_f} \sum_{k \in K_f} s_f^{(k)}). \end{aligned}$$



# Basic DRI model

- As for the related one-stage DRO model [3], we demonstrate that, under mild assumptions, the basic DRI model is *asymptotically consistent*.
- That is, when the leader and the follower acquire more data, their optimal solutions and objective function values converge, in a sense, to those of the underlying stochastic programming problem [SI].
- Given  $x \in X$ , the follower's convergence, as  $k_f \rightarrow \infty$ , is due to Theorem 3.6 in [3].
- The leader's convergence, as  $k_l \rightarrow \infty$  and  $k_f \rightarrow \infty$ , is based on the following assumption:
  - For every fixed  $x \in X$ , the full information follower's problem  $\max_{y \in Y(x)} \rho_f(y, \mathbb{Q}^*)$  has a **unique** optimal solution.

The proof also exploits discreteness of  $X$ , closedness of the follower's optimal solution set and the bounded convergence theorem.

# Example: packing interdiction

- $n = 4$  items, leader can block  $\leq 2$  items, follower selects  $\leq 1$  item.
- Let

$$S = \{c \in \mathbb{R}^4 : c_1 \in [0, 10], c_2 \in [1, 11], c_3 \in [6, 12], c_4 \in [7, 13]\},$$

with the true distribution  $\mathbb{Q}^*$  being a *discrete uniform distribution* with expected costs  $\bar{c}^* = (5, 6, 9, 10)^\top$ .

- Full-information **[SI]** model:

$$z^* = \min_{x \in X} \max_{y \in Y(x)} \bar{c}^{*\top} y,$$

where

$$X = \left\{x \in \{0, 1\}^4 : \sum_{i=1}^4 x_i \leq 2\right\} \text{ and } Y(x) = \left\{y \in [0, 1]^4 : y \leq 1 - x, \sum_{i=1}^4 y_i \leq 1\right\}$$

$$\bar{c}_1^* = 5$$

Item 1

$$\bar{c}_2^* = 6$$

Item 2

$$\bar{c}_3^* = 9$$

Item 3

$$\bar{c}_4^* = 10$$

Item 4

**Figure:** Optimal solution:  $x^* = (0, 0, 1, 1)^\top$ ,  $y^* = (0, 1, 0, 0)^\top$ ,  $z^* = 6$ .

# Example: packing interdiction

- We assume that both decision-makers implement a myopic *sample average approximation* based on  $k_I = k_f = 5$  samples given by:

$$\hat{C}_I = \begin{pmatrix} 10 & 11 & 8 & 7 \\ 6 & 7 & 7 & 9 \\ 10 & 10 & 7 & 9 \\ 9 & 11 & 6 & 7 \\ 10 & 11 & 7 & 8 \end{pmatrix} \quad \text{and} \quad \hat{C}_f = \begin{pmatrix} 3 & 3 & 8 & 13 \\ 1 & 2 & 9 & 10 \\ 3 & 11 & 11 & 13 \\ 10 & 3 & 10 & 12 \\ 4 & 1 & 12 & 7 \end{pmatrix},$$

with average profits  $\bar{c}_I := (9, 10, 7, 8)^\top$  and  $\bar{c}_f := (4, 5, 10, 11)^\top$ .

- Notably, Assumption **A3** is violated.

# Example: packing interdiction

- **True basic model:** Hypothetical benchmark assuming the leader knows the follower's data and its optimal policy.
- **Pessimistic approximation:** Assumes no knowledge of the follower's data; leader optimizes against the worst-case feasible policy.
- **Semi-pessimistic approximation:** Leader has partial information about the follower's data, e.g., it is aware of the first two columns in  $\hat{C}_f$ . Then, the leader constructs an uncertainty set to estimate the follower's average profits and makes a robust decision.
- **Augmented basic model:** The leader replaces the missing follower's data with its own, effectively assuming both players use the same data.

Model	Leader's Decision	Leader Expected Obj.	Leader's True Obj.
Full information	$(0, 0, 1, 1)^T$	6	6
True basic	$(0, 0, 0, 1)^T$	7	9
Pessimistic	$(1, 1, 0, 0)^T$	8	10
Semi-pessimistic	$(0, 0, 0, 1)^T$	7	9
Augmented basic	$(1, 1, 0, 0)^T$	8	10

# True basic model

## True basic model

$$\begin{aligned} [\text{DRI}^*]: \quad z_b^* &:= \min_{x,y} \left\{ \max_{Q_I \in \mathcal{Q}_I} \rho_I(y, Q_I) \right\} \\ \text{s.t. } x &\in X \\ y &\in \operatorname{argmax}_{\tilde{y} \in Y(x)} \left\{ \min_{Q_f \in \mathcal{Q}_f(\hat{\mathcal{C}}_f^*)} \rho_f(\tilde{y}, Q_f) \right\}, \end{aligned}$$

where  $\hat{\mathcal{C}}_f^*$  denotes the true follower's data set, and the second part of Assumption **A3** does not necessarily hold.

## Pessimistic approximation

$$[\mathbf{DRI-P}]: \hat{z}_p^* := \min_{x \in X} \max_{y \in Y(x)} \max_{Q_I \in \mathcal{Q}_I} \rho_I(y, Q_I).$$

- The leader in **[DRI-P]** disregards any available information about the true follower's data set  $\hat{C}_f^*$ .
- It selects the worst-case *feasible* follower's policy in terms of the leader's objective function value.

# Pessimistic approximation

- We show that **[DRI-P]** is  $\Sigma_2^P$ -hard by a reduction from the dominating set interdiction problem [4].
- To solve **[DRI-P]**, we design a Benders decomposition algorithm tailored to two special cases of the problem, where the leader is either *risk-neutral* or *ambiguity-free*.
- The algorithm for each case is based on the standard decomposition techniques for bilevel optimization and leverages a disjoint bilinear structure of the inner optimization problem.

# Semi-pessimistic approximation

## Semi-pessimistic approximation

$$\begin{aligned} \text{[DRI-SP]: } \hat{z}_{sp}^* &:= \min_{x \in X} \max_{\hat{C}_f \in \hat{S}_I} \min_y \left\{ \max_{Q_I \in \mathcal{Q}_I} \rho_I(y, Q_I) \right\} \\ \text{s.t. } y &\in \operatorname{argmax}_{\tilde{y} \in Y(x)} \left\{ \min_{Q_f \in \mathcal{Q}_f(\hat{C}_f)} \rho_f(\tilde{y}, Q_f) \right\}. \end{aligned}$$

- We relax Assumption **A3**, i.e., for each  $k \in \{1, \dots, k_f\}$ , the leader either knows that

$$\hat{c}_f^{(k)} \subseteq \hat{S}_I^{(k)} := \left\{ c \in \mathbb{R}^n : \underline{b}^{(k)} \leq c \leq \bar{b}^{(k)} \right\} \subseteq S \text{ or } \hat{c}_f^{(k)} \in \hat{C}_I.$$

- The leader in [DRI-SP] assumes the worst-case possible realization of  $\hat{C}_f^*$  set w.r.t. the uncertainty set  $\hat{S}_I = \hat{S}_I^{(1)} \times \dots \times \hat{S}_I^{(k_f)}$ .



# Semi-pessimistic approximation

- We show that **[DRI-SP]** is  $\Sigma_2^P$ -hard from the robust optimistic bilevel problem with interval uncertainty [5].
- We propose to use a discretization of the uncertainty set  $\hat{S}_I$ , based on a finite number of scenarios for the follower's data set  $\hat{C}_f$ . The discretization is shown to admit a single-level MILP formulation, whose size however is proportional to the number of scenarios.
- To justify our approach, we show that the discretized semi-pessimistic approximation is almost surely *robust* with respect to the follower's data in an asymptotic sense, i.e., when the number of scenarios tends to infinity.

# Computational settings

- We consider a class of general interdiction problems defined as:

$$\min_{x \in X} \max_{y \in Y(x)} c^\top y,$$

where

$$X = \{x \in \{0, 1\}^n : Hx \leq h\} \quad \text{and} \quad Y(x) = \{y \in \mathbb{R}_+^n : \tilde{F}y \leq \tilde{f}, y \leq U(1 - x)\}.$$

- $d_I = \dim(h) = 1$ ,  $\tilde{d}_f = \dim(\tilde{f}) = 10$  and  $U = I$ .
- All elements of  $H$  and  $\tilde{F}$  are generated uniformly at random from the interval  $[0.01, 1]$ , whereas

$$h_j = 0.4 \sum_{i=1}^n H_{ji} \quad \forall j \in \{1, \dots, d_I\}, \quad f_j = 0.4 \sum_{i=1}^n F_{ji} \quad \forall j \in \{1, \dots, \tilde{d}_f\}.$$

# Computational settings

- The support set is given by:

$$S := \{c \in \mathbb{R}^n : c_i \in [0.01, 1] \quad \forall i \in N := \{1, \dots, n\}\}.$$

- The true distribution  $\mathbb{Q}^*$  is defined as a product of truncated normal distributions for each component  $c_i$ ,  $i \in N$ , with increasing mean and variance, similar to [3].
- We set  $\varepsilon_l := \frac{\delta_l}{\sqrt{k_l}}$  and  $\varepsilon_f := \frac{\delta_f}{\sqrt{k_f}}$  for some  $\delta_l, \delta_f \in \mathbb{R}_+$ , and  $\alpha_l = \alpha_f = 0.95$ .
- To define  $\hat{S}_l$ , it is assumed that the first  $k_{lf} \leq k_f$  samples in  $\hat{C}_f^*$  are known to the leader; for each  $i \in N$  and  $k \in \{k_{lf} + 1, \dots, k_f\}$ , let

$$(\hat{c}_f^{(k)})_i \in [c_{ki}^* - \kappa \Delta_{ki}, c_{ki}^* + \kappa(1 - \Delta_{ki})] \cap [0.01, 1].$$

Here,  $\kappa$  is a fixed noise level, and  $\Delta_{ki} \in [0, 1]$  is a shift parameter ( $\Delta_{ki}$  is 0 with probability 0.5, and with probability 0.5, it is selected uniformly at random from the interval  $[0, 1]$ ).

# Performance metrics

- Let  $(\hat{x}^*, \hat{y}^*)$  be our obtained solution and  $\hat{z}^*$  the respective optimal objective function value.
- The follower's relative out-of-sample loss is defined as:

$$RL_f^{(\text{out})} = \frac{\rho_f(\hat{y}^*, Q^*)}{\max_{y \in Y(\hat{x}^*)} \rho_f(y, Q^*)} \leq 1$$

- The leader's relative out-of-sample loss is defined as:

$$RL_l^{(\text{out})} = \frac{\rho_l(\hat{y}^*, Q^*)}{\rho_l(\tilde{y}^*, Q^*)} \geq 1,$$

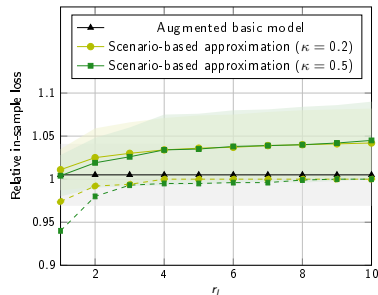
where  $\tilde{y}^*$  solves the full information leader's problem.

- The relative leader's in-sample loss is defined as:

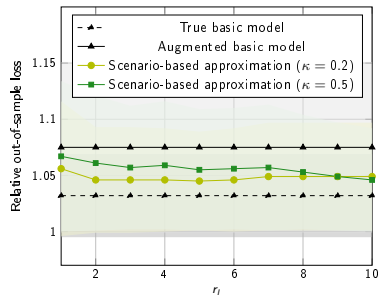
$$RL_l^{(\text{in})} = \frac{\hat{z}^*}{z_b^*}.$$

# Selected computational results

- $n = 10$ , 10 random test instances, and 10 data sets for each instance.



(a)

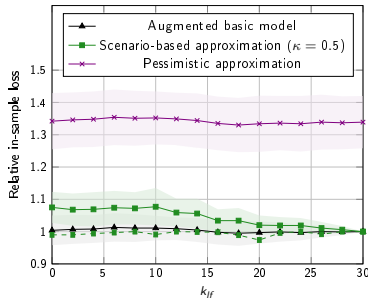


(b)

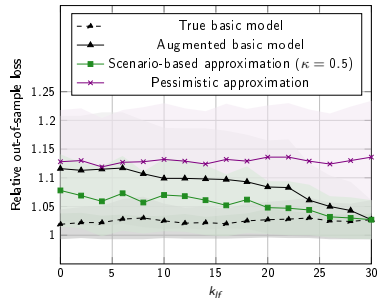
**Figure:** The average relative in-sample and out-of-sample loss (with MADs) as a function of the number of scenarios,  $r_I$ , for  $k_I = k_f = 30$ ,  $\delta_I = \delta_f = 0.1$  and  $k_{If} = 20$ . The dashed lines in (a) correspond to the empirical 5% percentile of the relative in-sample loss.

# Selected computational results

- $n = 10$ , 10 random test instances, and 10 data sets for each instance.



(a)

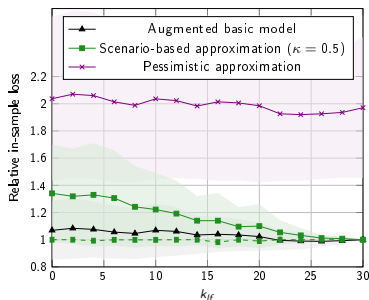


(b)

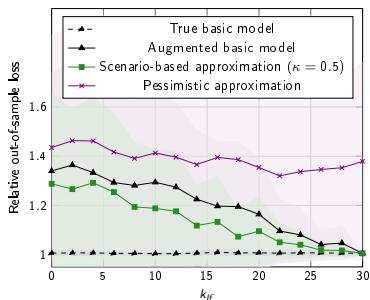
**Figure:** The average relative in-sample and out-of-sample loss of the *ambiguity-free* leader (with MADs) as a function of  $k_{lf}$ , for  $k_l = k_f = 30$ ,  $\delta_l = \delta_f = 0.1$  and  $\alpha_l = 0.9$ . The follower is assumed to be *risk-averse*. The dashed line corresponds to the empirical 5% percentile of the relative in-sample loss for the scenario-based semi-pessimistic approximation.

# Selected computational results

- $n = 10$ , 10 random test instances, and 10 data sets for each instance.



(a)

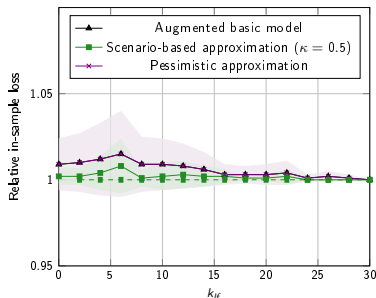


(b)

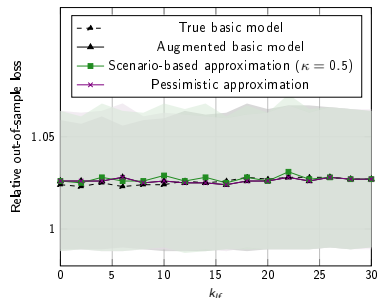
**Figure:** The average relative in-sample and out-of-sample loss of the *risk-neutral* leader (with MADs) as a function of  $k_{lf}$ , for  $k_l = k_f = 30$  and  $\delta_l = \delta_f = 0.1$ . The follower is assumed to be *risk-averse*. The dashed line corresponds to the empirical 5% percentile of the relative in-sample loss for the scenario-based semi-pessimistic approximation.

# Selected computational results

- $n = 10$ , 10 random test instances, and 10 data sets for each instance.



(a)

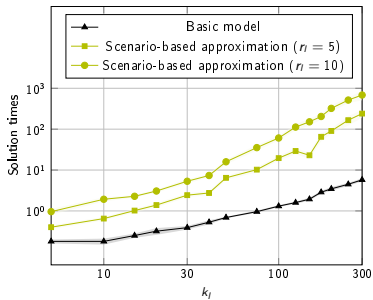


(b)

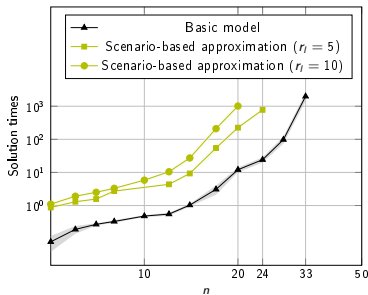
**Figure:** The average relative in-sample and out-of-sample loss of the *risk-neutral* leader (with MADs) as a function of  $k_{lf}$ , for  $k_l = k_f = 30$  and  $\delta_l = \delta_f = 0.1$ . The follower is assumed to be *risk-neutral*. The dashed line corresponds to the empirical 5% percentile of the relative in-sample loss for the scenario-based semi-pessimistic approximation.



# Selected computational results



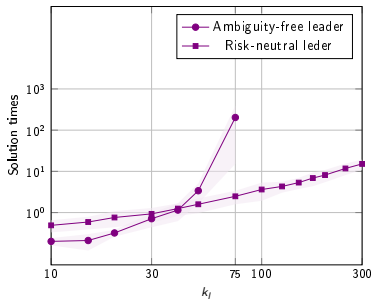
(a)  $k_f = k_l$ ,  $n = 10$ .



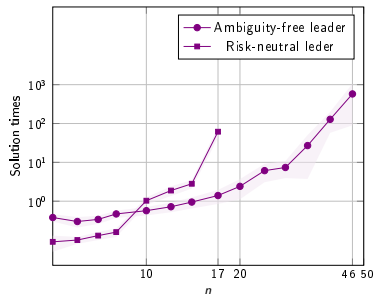
(b)  $k_l = k_f = 30$ .

**Figure:** Average solution times in seconds (with MADs) of the basic and semi-pessimistic formulations as a function of  $k_l = k_f$  (a) and  $n$  (b), for  $\delta_l = \delta_f = 0.1$ ,  $k_{lf} = \lfloor \frac{2}{3} k_f \rfloor$  and  $\kappa = 0.2$ , evaluated over 10 random test instances. The time limit is set to 60 minutes.

# Selected computational results



(a)  $n = 10$ .



(b)  $k_I = 30$ .

**Figure:** Average solution times in seconds (with MADs) of the pessimistic formulation as a function of  $k_I$  (a) and  $n$  (b), for  $\delta_I = 0.1$  and  $\alpha_I = 0.9$ , evaluated over 10 random test instances. The time limit is set to 60 minutes.

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