

# 13ème Atelier en Évaluation des Performances

Toulouse, Dec 2-4, 2024

Decision-Epoch Matters: Unveiling its Impact on the Stability of  
Scheduling with Randomly Varying Connectivity

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Application example:  
Supply chains

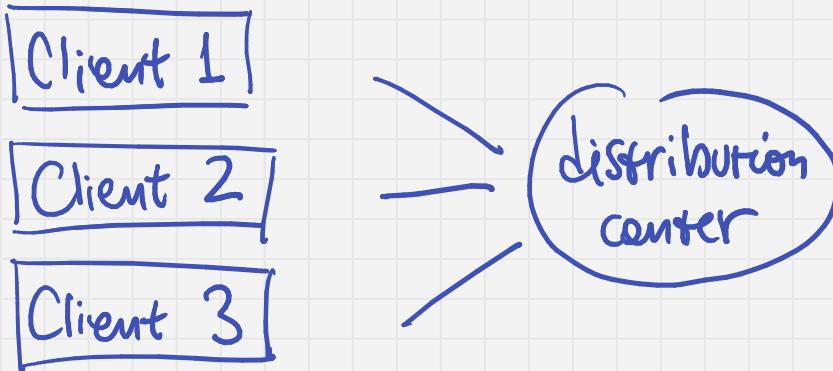
Queuing ; random  
systems ; environment ;  
timing restrictions over  
decision making

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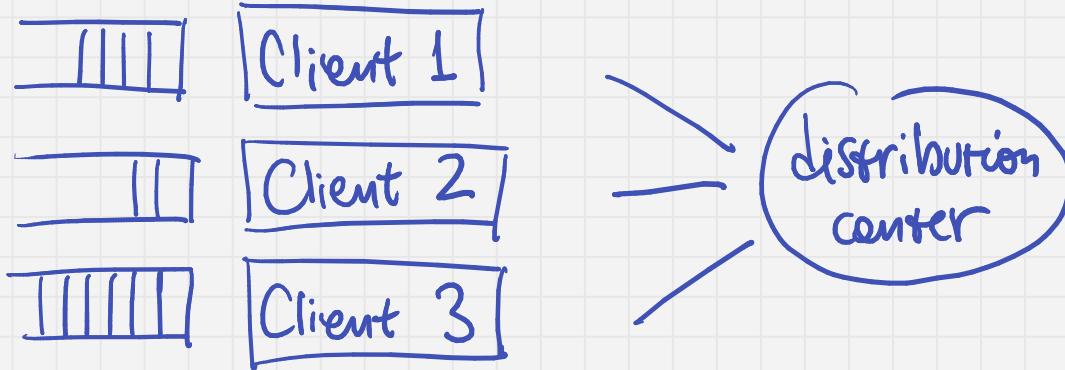
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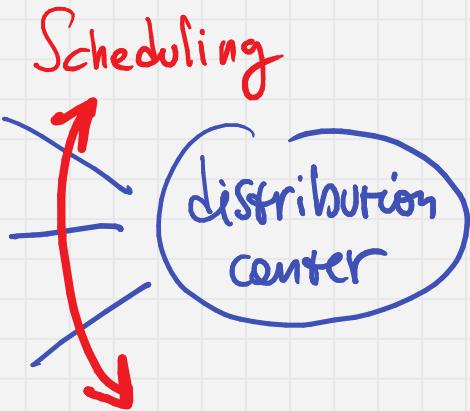
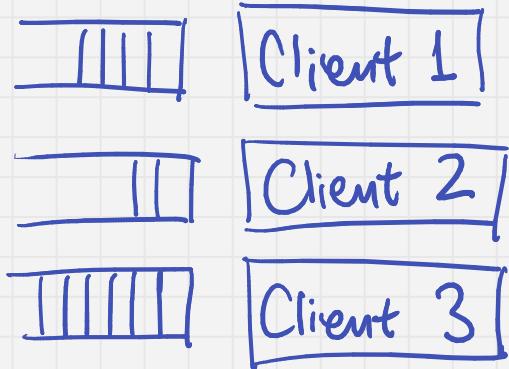


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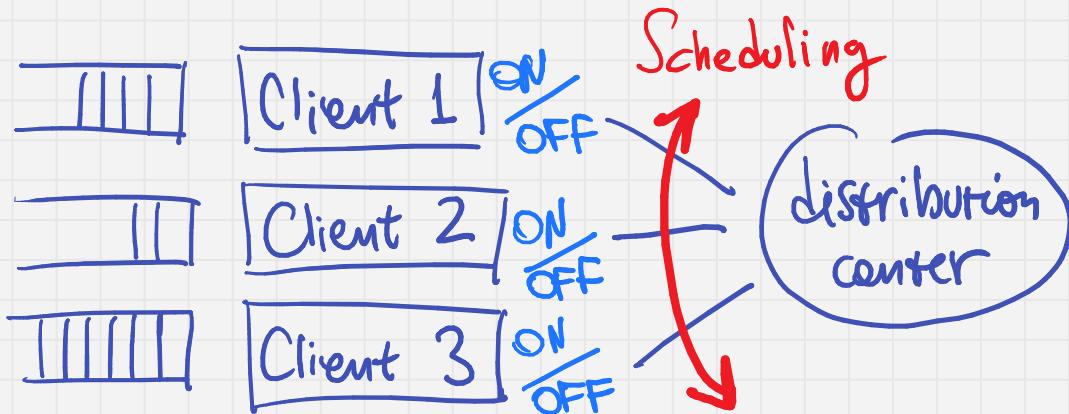
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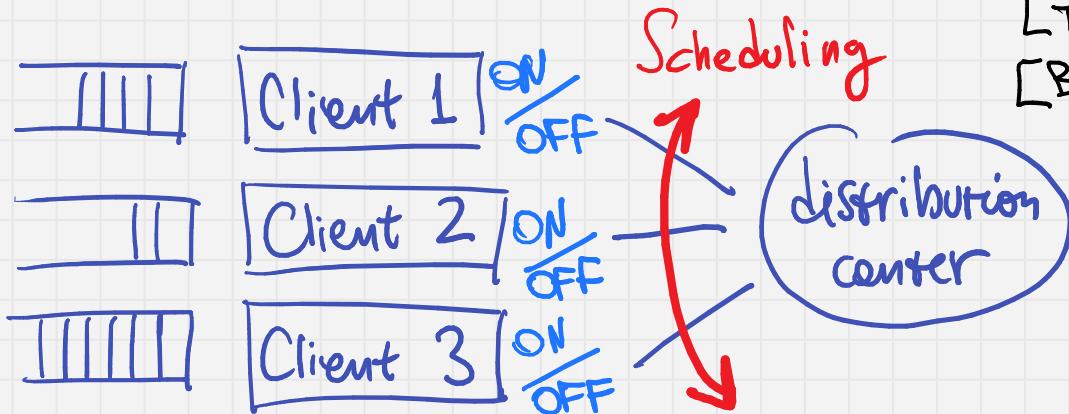
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queueing systems ; random environment ; timing restrictions over decision making



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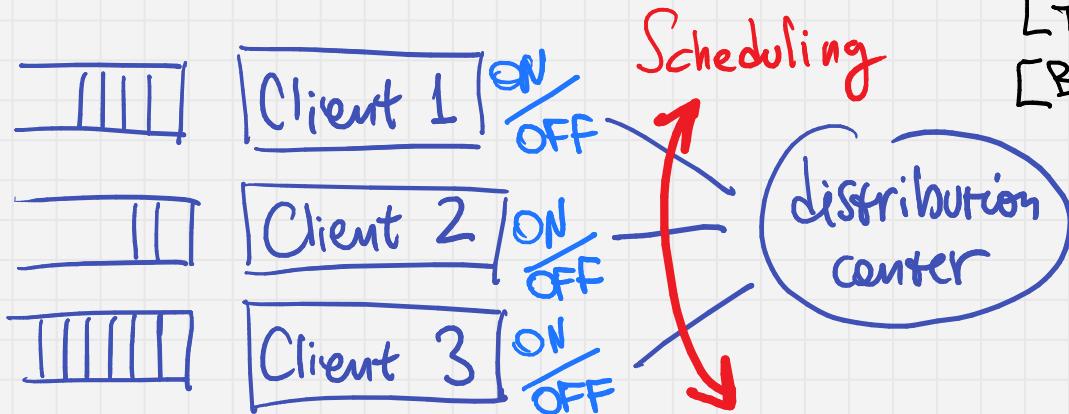
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[Tassiulas, Ephremides; 1993]

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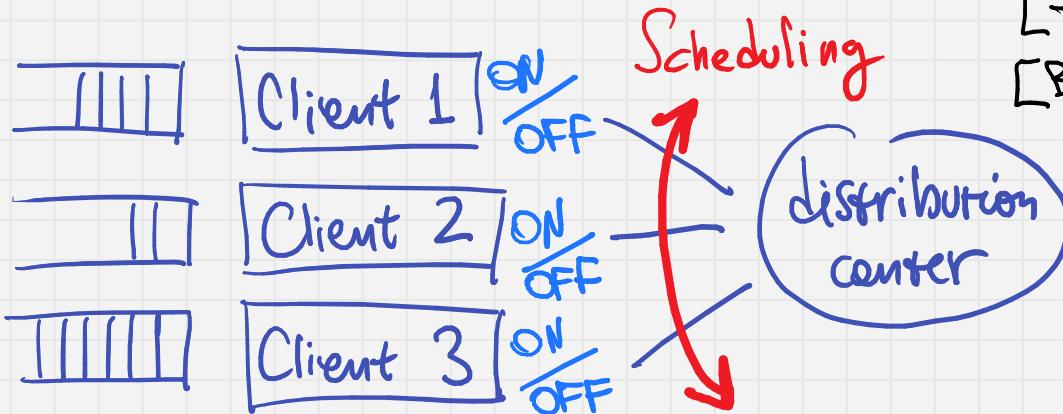
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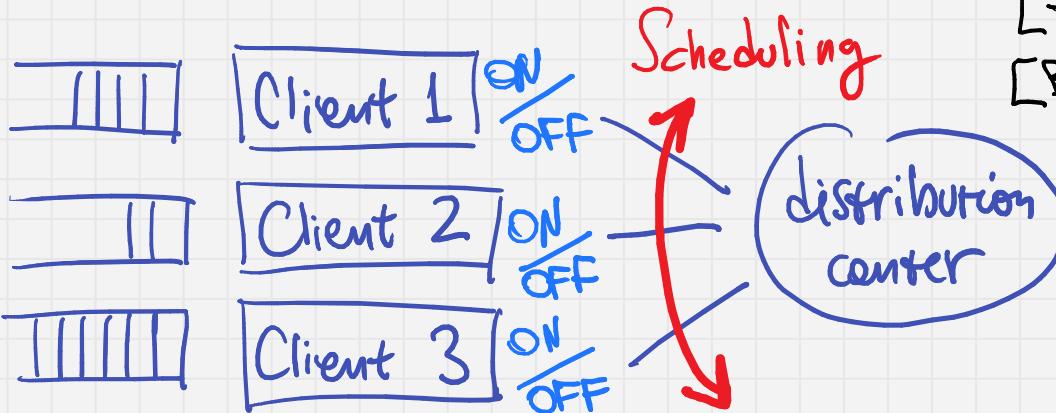
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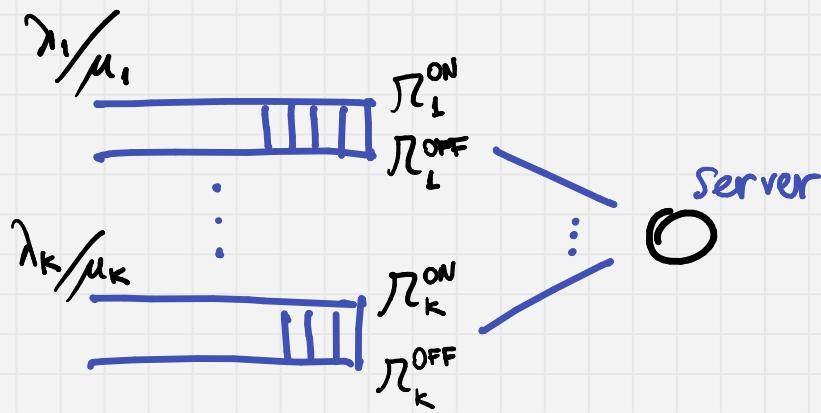
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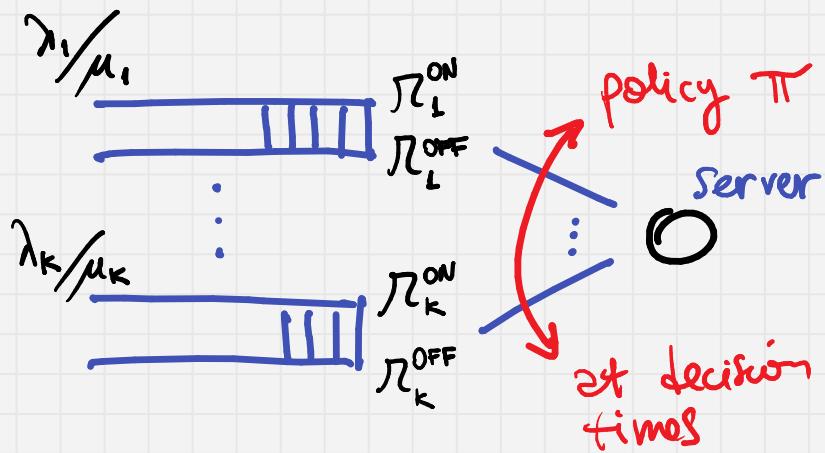
Performance guarantees?

Optimal policies?

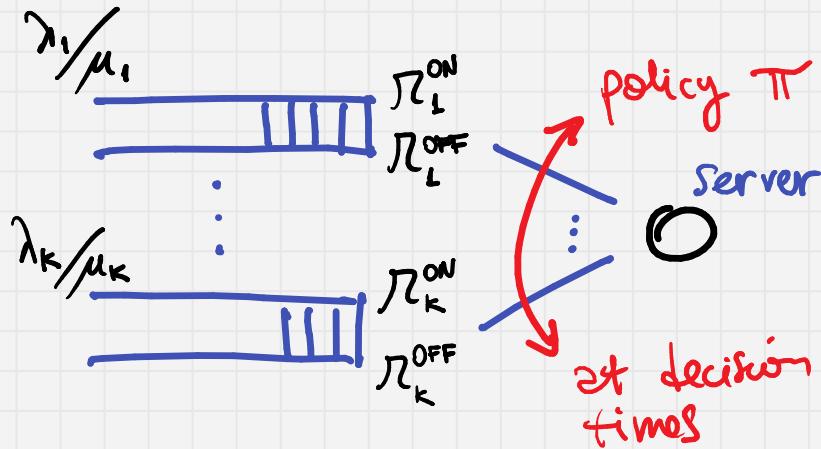
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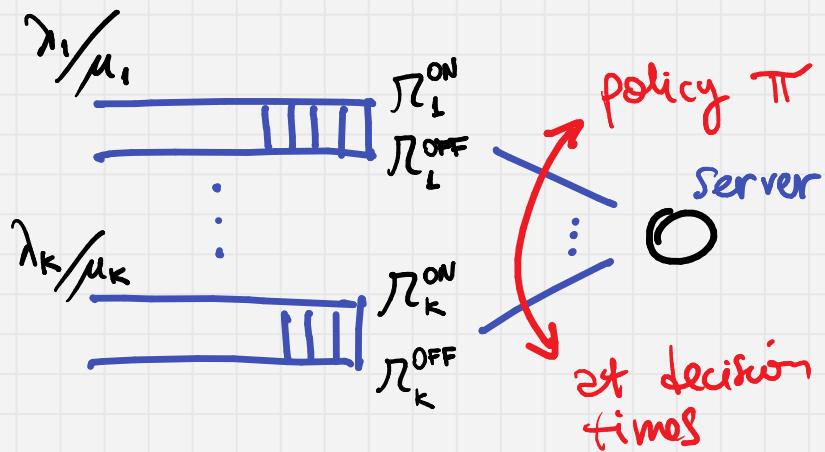
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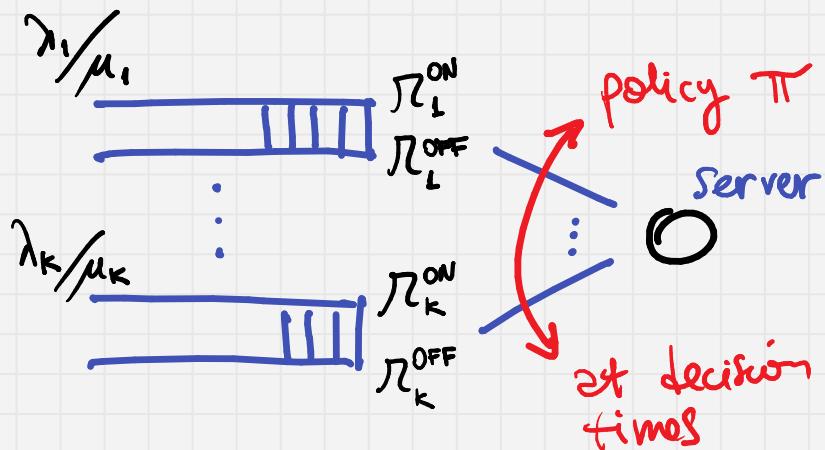
Timing restrictions:

Setting 1: unconstrained

Setting 2: non-preemptive

Setting 3:  $\gamma$ -rate exponentially

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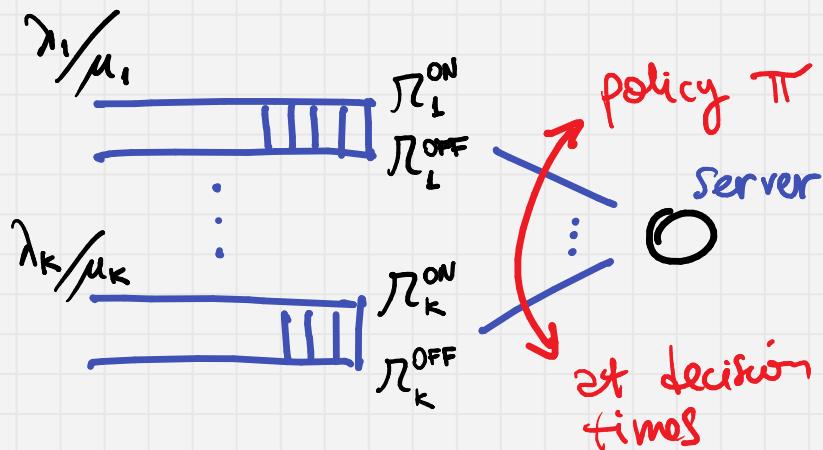
positive recurrent (stable)

Setting  $j \}$  → Markov process

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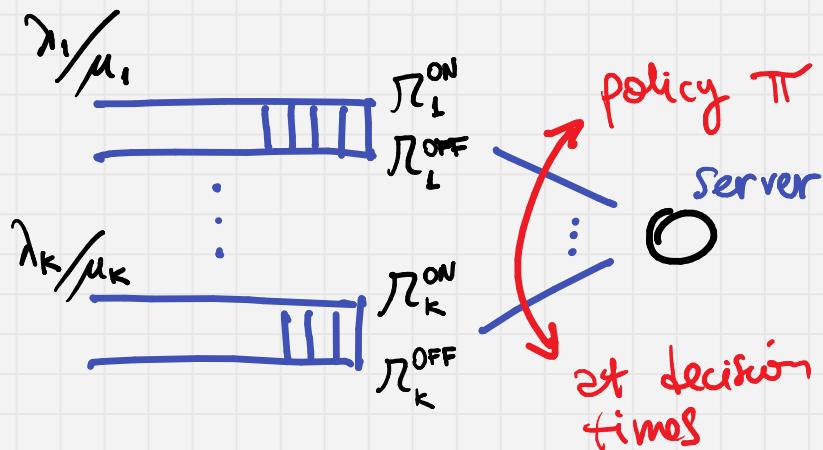
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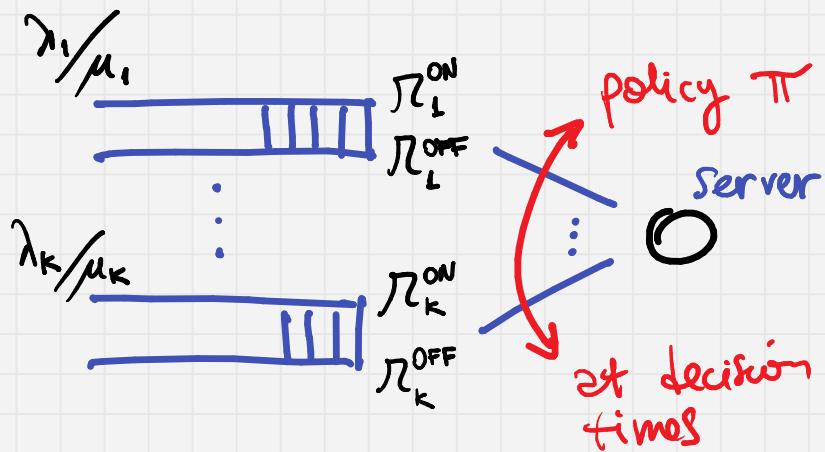
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$\pi^*$  is throughput optimal if

$$SR_j(\pi^*) = MSR_j$$

## Main theorem

$$f_i = \frac{\lambda_i}{\mu_i}, \quad p_i^{\text{ON}} = \% \text{ time ON queue } i = \frac{R_i^{\text{ON}}}{R_i^{\text{ON}} + R_i^{\text{OFF}}} = 1 - p_i^{\text{OFF}}$$

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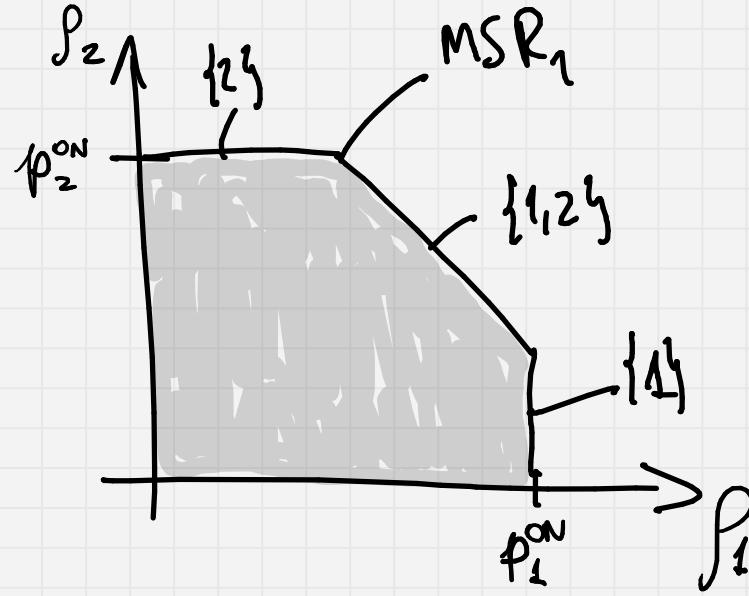
MSR<sub>1</sub>:  $\forall L \subseteq \{1, \dots, k\}, \underbrace{\sum_{i \in L} p_i}_{\text{load at } L} < \underbrace{1 - \prod_{i \in L} p_i^{\text{OFF}}}_{\text{maximum capacity}}$

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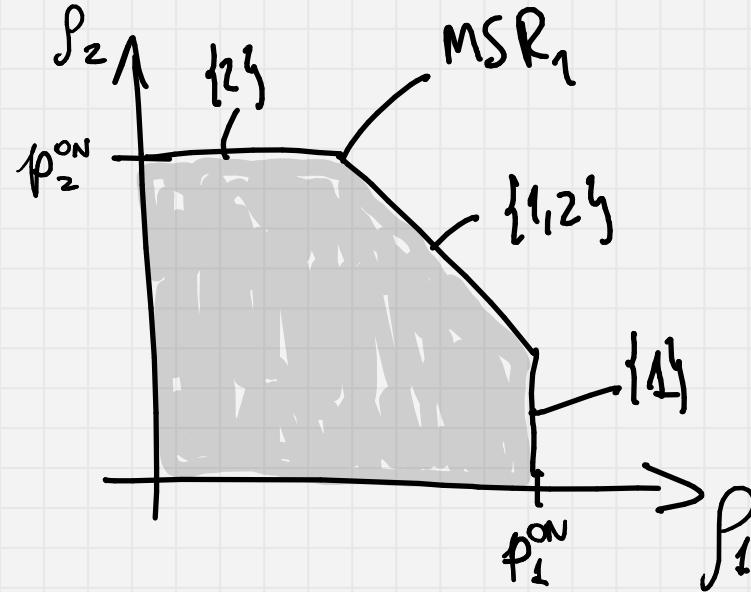
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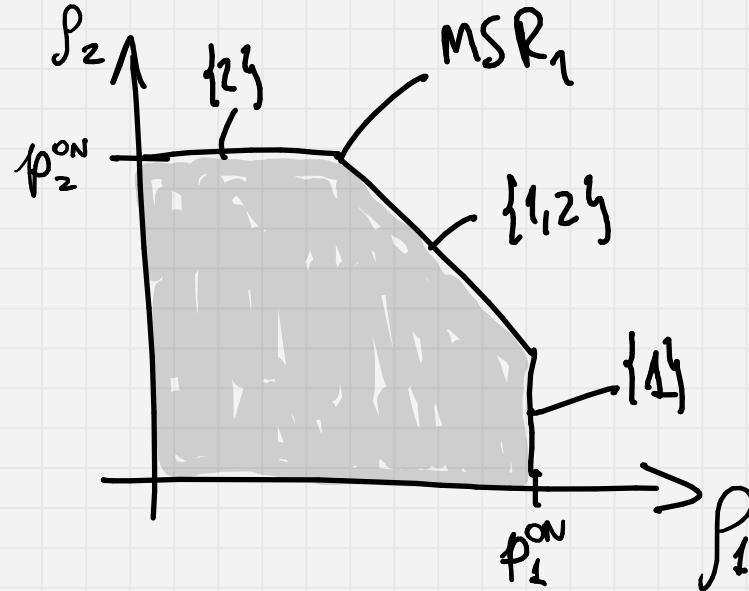
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(unconstrained)

↑ not new

$\text{MSR}_2$ :  $\sum_{i=1}^k \frac{f_i}{p_i^{\text{ON}}} \leq 1$

(non-preemptive)



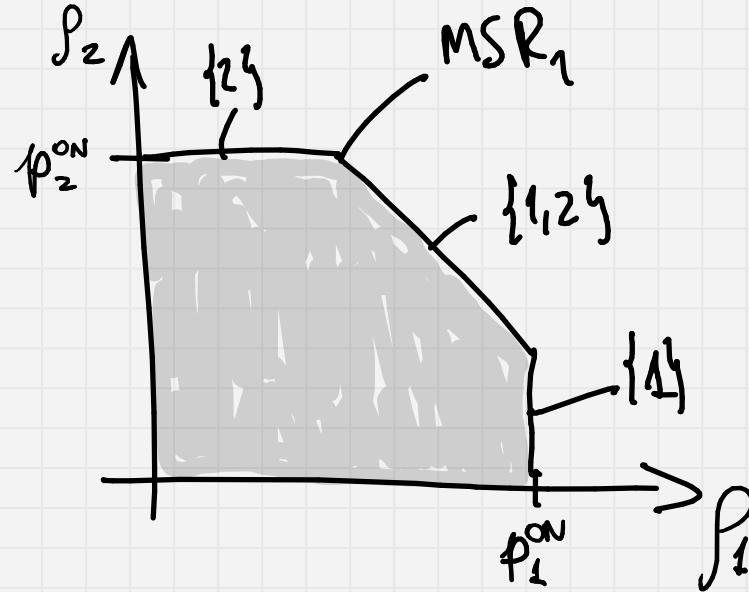
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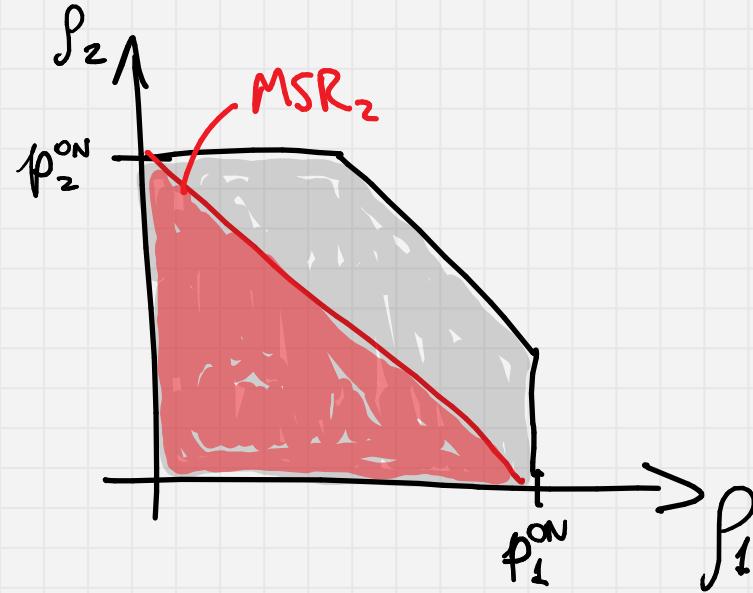
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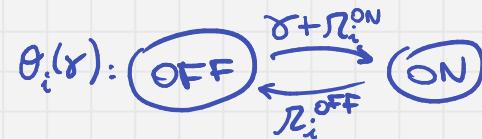
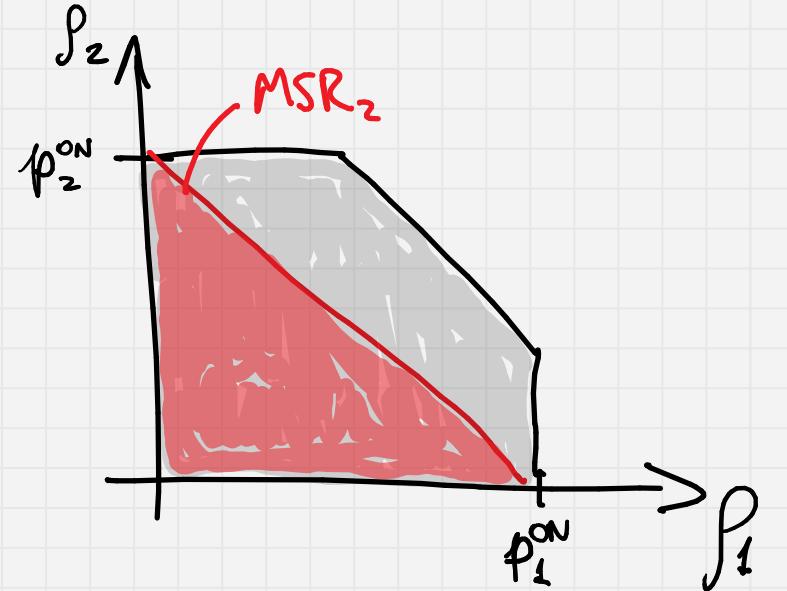
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(unconstrained)

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$\text{MSR}_2$ :  $\sum_{i=1}^k \frac{p_i}{p_i^{\text{ON}}} < 1 \leftarrow$  multi-class M/G/1  
(non-preemptive)

$\text{MSR}_3(r)$  - CONNECTED POLICIES:  $\forall L \subseteq \{1, \dots, k\}, \sum_{i \in L} \frac{p_i}{\theta_i(r)} < 1 - \prod_{i \in L} p_i^{\text{OFF}}$   
(r-exponentials)



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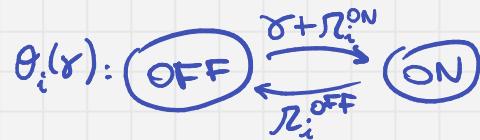
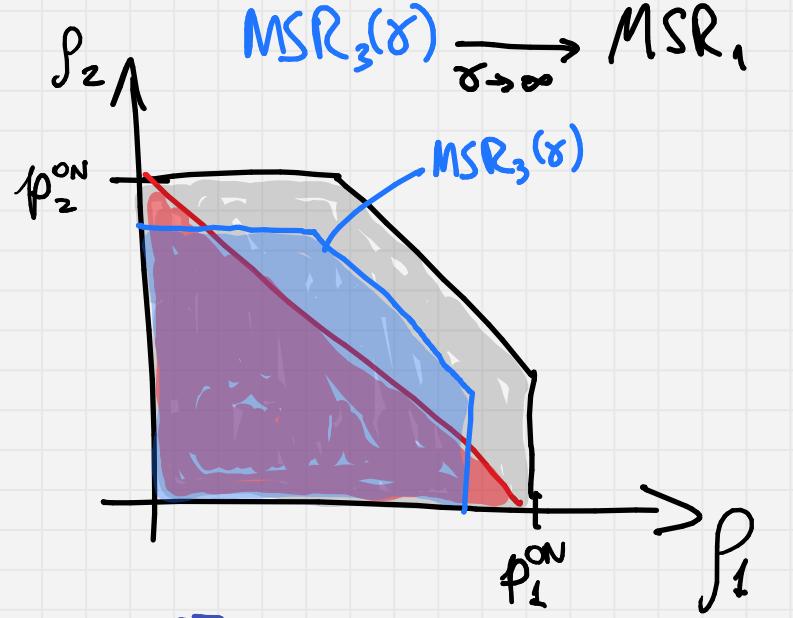
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$\text{MSR}_3(\gamma)$  - CONNECTED POLICIES:  $\forall L \subseteq \{1, \dots, k\}, \sum_{i \in L} \frac{p_i}{\theta_i(\gamma)} < 1 - \prod_{i \in L} p_i^{\text{OFF}}$

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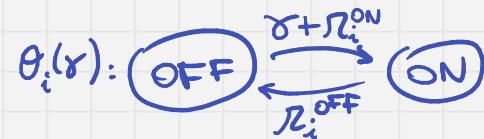
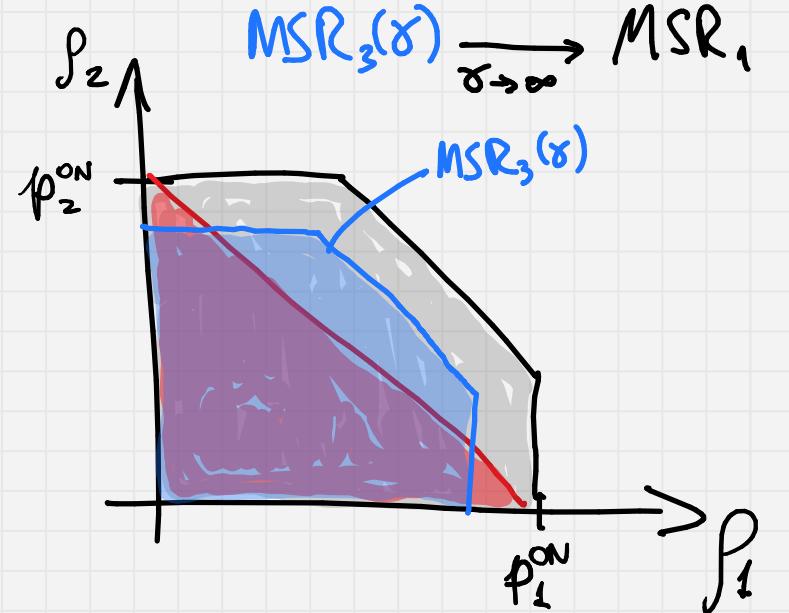
$\underbrace{\text{load at } L}_{\text{maximum capacity}}$

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$\forall c_1, \dots, c_K > 0, \arg \max_{i \in \{1, \dots, K\}} c_i Q_i \mathbb{1}\{E_i = \text{ON}\}$  is optimal.



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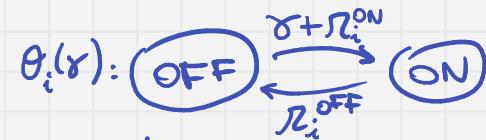
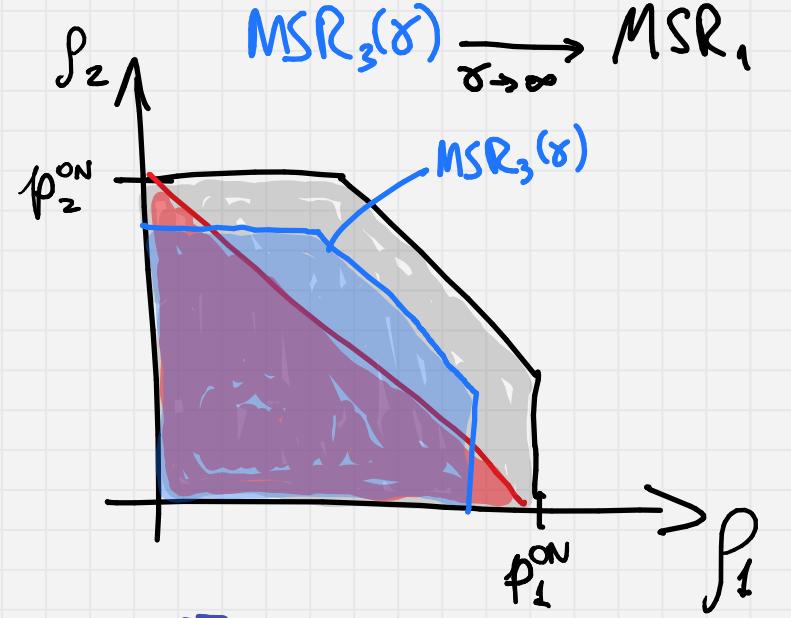
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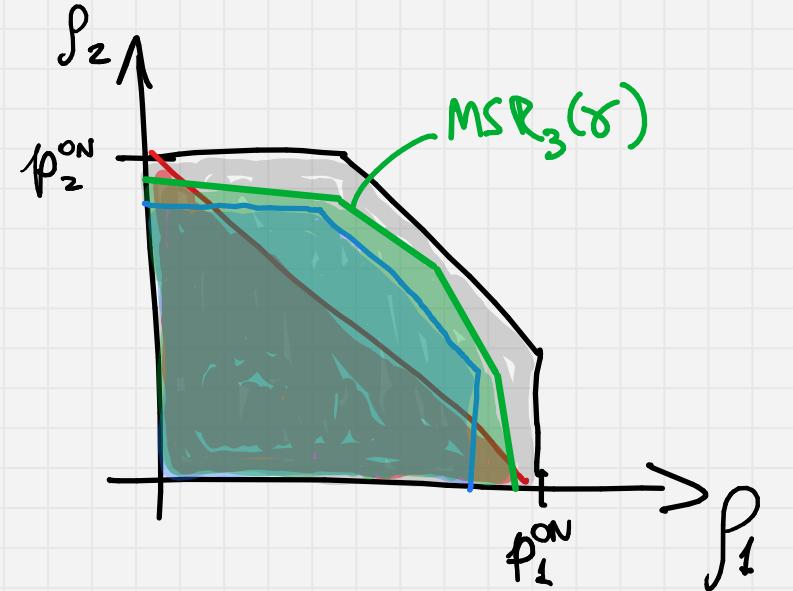
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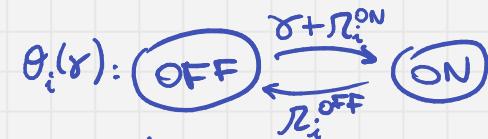
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$\text{MSR}_3(\gamma) - \text{NON-CONNECTED POLICIES}:$  Minkowski sum optimal:  $\arg \max_{\substack{i \in \{1, \dots, K\} \\ E \in \{0, 1\}}} c_i Q_i (\mathbb{1}\{E_i = \text{ON}\} + k_i \mathbb{1}\{E_i = \text{OFF}\})$

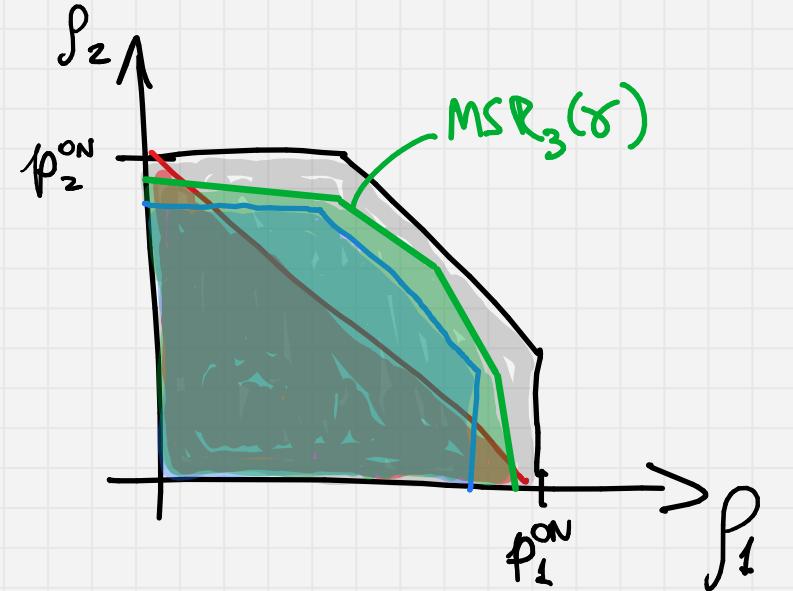
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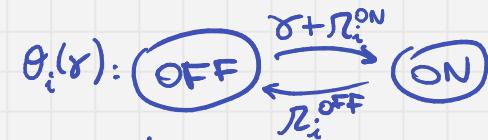
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└ or at L      maximum capacity

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$\text{MSR}_3(\gamma)$  — NON-CONNECTED POLICIES: Minkowski sum of triangles — optimal:  $\arg \max_{i \in \{1, \dots, K\}} c_i Q_i (\mathbb{1}\{E_i = \text{ON}\} + \gamma_i \mathbb{1}\{E_i = \text{OFF}\})$

Obs: they're all agnostic policies.

[Andrews, Kumaran, Ramanan, Stolyar, Vijayakumar, Whiting; 2004]

## Proof: test for fluid limits (TFL)

Classical result  
(Dai, Rybko-Stolyar) : Stable Fluid model  $\implies$  Stable stochastic model

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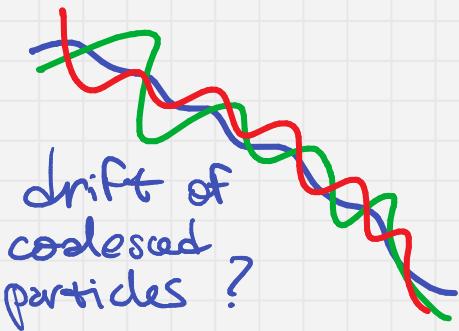
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↑  
Maximum  
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fluid Lyapunov



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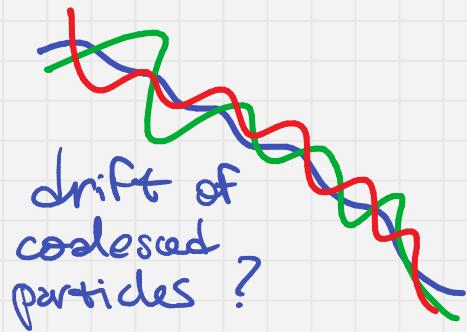
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Stable  
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wave length  
fluid Lyapunov

$$\frac{M(t_2) - M(t_1)}{t_2 - t_1} \leq -\delta$$

$t_1 < t_2$



# Proof: test for fluid limits (TFL)

Classical result  
(Dai, Rybko-Stolyar)

Stable  
Fluid  
model



Maximum  
wave length  
fluid Lyapunov

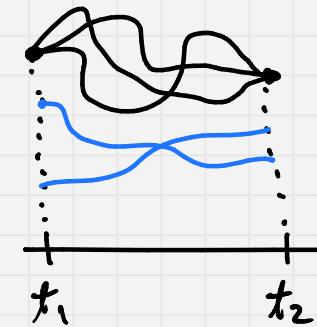
Stable  
stochastic  
model

$$\frac{M(t_2) - M(t_1)}{t_2 - t_1} \leq -\delta$$

~~$t_1 < t_2$~~

TFL

for  $t_1 < t_2$  s.t.



# Proof: test for fluid limits (TFL)

Classical result  
(Dai, Rybko-Stolyar)

: Stable Fluid model  $\Rightarrow$  Stable stochastic model

↑  
Maximum wave length  
fluid Lyapunov

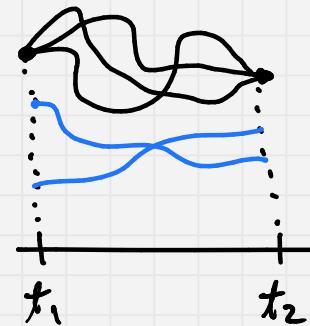
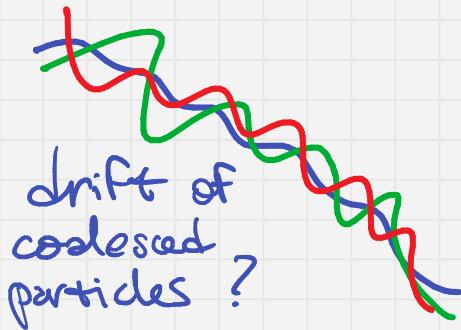
$$\frac{M(t_2) - M(t_1)}{t_2 - t_1} \leq -\delta$$

~~$t_1 < t_2$~~

TFL.

for  $t_1 < t_2$  s.t.

Under the formal assumption,  
the control of the slope follows easily!



## Conclusions

- In a queuing network with a random environment, we studied how timing restrictions on decision making affect the stability of the system.  
More precisely, for different timing restriction settings,
  - we explicitly described the MSRs,
  - and we provided a family of MW optimal policies.
- We developed a formal test for fluid limits (TFL) which, if passed, it guarantees the stability of the model.

Thanks !