# International Conference on Networks, Games, Control and Optimization (NetGCoop)

The next edition will be held in Bilbao from 08/10/2025 to 10/10/2025



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# Stability and performance of multi-class queueing systems with unknown service rates: A scheduling-based approach

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- *C* classes of trafic flows.
- *N* mutually exclusive service options/modes.
- Input-queued system.
- Real-world examples : channel/frequency selection in wireless communications (e.g. WiFi networks or cognitive radio systems), ...





- Time slotted operation.
- The number of jobs that arrive per class is independent across time slots.





Each time slot,

- a single service mode/option can be selected,
- when service mode/option s is selected in time slot t, (up to)  $R_{c,s}(t)$  class-c jobs will be served
- neither realizations nor statistics of  $R_{c,s}(t)$  are known to scheduling agent





Scheduling agent :

- observes the global state of the queues,
- can infer service rates  $R_{c,s}(t)$  from evolution of queue lengths, but does not have any advance knowledge of realizations or underlying statistics
  - in stark contrast to conventional assumptions.





### **Objective:**

Design scheduling algorithm that

- achieves maximum stability (throughput optimality), and
- provides (near-)optimal response times.





For analysis purposes, we assume that

- the number of jobs that arrive per time slot and class ~ Geometric with mean  $\lambda_c$  for class c.
- the number of served jobs of class c at service option  $s \sim$  Geometric with mean  $\mu_{s,c}$ .



**Stability region:** Given the set of mean arrival rates  $(\lambda_1, \lambda_2, ..., \lambda_C)$  and mean service rates  $\mu_s = (\mu_{s,1}, \mu_{s,2}, ..., \mu_{s,C})$  for service option s. There exists a vector  $(\overline{\lambda_1}, \overline{\lambda_2}, ..., \overline{\lambda_C}) \in convex hull (\mu_1, \mu_2, ..., \mu_N)$  such that  $(\lambda_1, \lambda_2, ..., \lambda_C) < (\overline{\lambda_1}, \overline{\lambda_2}, ..., \overline{\lambda_C})$  component-wise.



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### The above stability condition ...

- is **necessary** for all algorithms that do not have advance knowledge of the realizations of the service rates  $R_{c,s}(t)$  and
- **sufficient** for the algorithm that we will propose.

maximum stability for our algorithm.

#### The above stability condition ...

 is not necessary in case of scheduling algorithms that do have advance knowledge of the realizations of service rates (channel-aware, `opportunistic').



**Proposition :** Consider the system with a single traffic class and N service options with service rates  $\mu_s$  for service option s. The system is stable if the following holds:

 $\lambda_1 < \max\{\mu_1, \mu_2, \dots, \mu_N\}.$ 



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### Algorithm :

• Q(t) = the number of jobs in queue at time t with

 $Q(t) = Q(t-1) + A(t) - R_{S(t)}(t),$ 

and A(t) and  $R_{S(t)}(t)$ , number of arrivals and departures when the service option is S(t).

- Fix Z(0) = Q(0) the threshold value : for every t > 0:
  - If  $Z(t-1) \ge Q(t-1)$ : Z(t) = Z(t-1) and S(t) = S(t-1).
  - If Z(t-1) < Q(t-1) or Q(t) = 0:  $Z(t) = \frac{Z(t-1)}{2}$  and  $S(t) \sim Unif(serv.opt)$

### **3. STABILITY REGION SINGLE CLASS**





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**Proposition :** Consider the system with a single traffic class with arrival rates  $(\lambda_1, \lambda_2)$ , and N service options with service rates  $\mu_s = (\mu_{s,1}, \mu_{s,2})$  for service option s. There exists a vector

 $(\overline{\lambda_1}, \overline{\lambda_2}) \in convex hull (\mu_1, \mu_2, ..., \mu_N)$ such that  $(\lambda_1, \lambda_2) < (\overline{\lambda_1}, \overline{\lambda_2})$  component-wise.



### **Queue dynamics :**

•  $Q_c(t)$  = the number of class c jobs in queue at time t with  $Q_c(t) = Q_c(t-1) + A_c(t) - R_{c,S(t)}(t)$ ,

and  $A_c(t)$  and  $R_{c,S(t)}(t)$ , number of arrivals and departures of class c when the service option is S(t).

- $L(t) = \sum_{c=1}^{C} Q_c^2(t).$
- Z(t) = the threshold value at time t, with fix  $Z(0) = L(0) = \sum_{c=1}^{C} Q_c^2(0)$ .
- $(X_1(t), X_2(t)) =$  the queue length per class of the threshold value a time t, with fix  $(X_1(0), X_2(0)) = (Q_1(0), Q_2(0))$ .



### **4. THRESHOLD BASED SCHEDULING ALGORITHM**

### Algorithm :

- At each time slot t > 1:
  - If  $Z(t-1) \ge L(t)$ :

keep going : Z(t) = Z(t - 1),  $(X_1(t), X_2(t)) = (X_1(t - 1), X_2(t - 1)),$ S(t) = S(t - 1).

• If 
$$Z(t-1) < L(t)$$
:  
Update:  $Z(t) = L(t)$   
 $(X_1(t), X_2(t)) = (Q_1(t), Q_2(t))$   
 $S(t) \sim Unif(service option)$ 

- fix  $\sigma$  small lower bound for Z(t).

• If  $min_{c \in C} \{Q_c(t)\} = 0$  and  $max_{c \in C} \{Q_c(t)\} \ge X_{argmax_{c \in C}} \{Q_c(t)\}(t) :$ Update: Z(t) = L(t) $(X_1(t), X_2(t)) = (Q_1(t), Q_2(t))$ 

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Example : C = 2 and N = 4  $S_1 \qquad \mu_{s_1,1} = 4, \\ \mu_{s_1,2} = 1$   $\lambda_1 = 3$   $\lambda_2 = 2$  Sch. agent  $S_2 \qquad \mu_{s_2,1} = 1, \\ \mu_{s_2,2} = 1$   $S_3 \qquad \mu_{s_3,1} = 2, \\ \mu_{s_3,2} = 4$  $S_4 \qquad \mu_{s_4,1} = 2, \\ \mu_{s_4,2} = 1$ 

	Se	Sec. of service options :													
time															
	0						6 - 12		3	14	15	16	17		
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18			19	20	21 2		22223- 27								
	2		2	1	2		2 1								



# **4. STABILITY REGION TWO CLASSES**



#### Challenges in the stochastic process:

 The number of times to sample a feasible service option with 'good' trajectory is unbounded, as well as the amount that the threshold value increases.

### Improvement in the fluid limit:

Vanish in the limit.



### **4. STABILITY REGION TWO CLASSES**



### **Conclusions :**

- We provide an algorithm that is maximally stable, for system that do not have any advance knowledge of the service rates.
- However, `lazy' scheduling by randomly resampling without really acting as long as things seem to move in the right direction.

### Future work :

• Improve (response time) performance, by using an active learning algorithm that learns to sample the best combination of service options.



### **5. CONCLUSION AND FUTURE RESEARCH**

# **THANK YOU!**



X-model\*: 2 traffic classes with 2 service options where either:

- server 1 (server 2) serves both classes simultaneously or,
- server 1 and server 2 both serve the same class



\* Yuan Zhong, Instability and stability of parameter agnostic policies in parallel server systems, Performance 2023.

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