Online Stochastic Matching Under the First-Come-First-Matched Policy

Céline Comte



"Online Stochastic Matching" Workshop Toulouse – September 27, 2024

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- Set \mathcal{V}_i of neighbors of node $i \rightarrow$ possible matches.



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The evolution of the sequence of unmatched item classes defines a Markov process whose transition rates depend on the graph G and the arrival rates μ_i , $i \in \mathcal{V}$.

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- We let $\mathbb{I}^{\scriptscriptstyle +} = \mathbb{I} \setminus \{ \emptyset \}$ denote the family of nonempty independent sets.



Outline

Literature review

Basic results

- Stability condition
- Definition of the discrete-time Markov chain
- Product-form stationary distribution and partial balance

3 Performance analysis

- Normalization constant
- Performance metrics
- Heavy-traffic analysis

4 Concluding remarks

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1 Literature review

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Early works

- René Caldentey, Edward H. Kaplan, and Gideon Weiss. "FCFS Infinite Bipartite Matching of Servers and Customers". Advances in Applied Probability 41.3 (Sept. 2009), pp. 695–730. DOI: 10.1239/aap/1253281061
- Ivo Adan and Gideon Weiss. "Exact FCFS Matching Rates for Two Infinite Multitype Sequences". *Operations Research* 60.2 (Apr. 1, 2012), pp. 475–489. DOI: 10.1287/opre.1110.1027

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Stability (i.e., positive recurrence) condition

 Ana Bušić, Varun Gupta, and Jean Mairesse. "Stability of the Bipartite Matching Model". *Advances in Applied Probability* 45.2 (June 2013), pp. 351–378. DOI: 10.1239/aap/1370870122

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Performance evaluation

- Ivo Adan et al. "Reversibility and Further Properties of FCFS Infinite Bipartite Matching". Mathematics of Operations Research 43.2 (Dec. 12, 2017), pp. 598–621. DOI: 10.1287/moor.2017.0874
- Céline Comte and Jan-Pieter Dorsman. "Performance Evaluation of Stochastic Bipartite Matching Models". Performance Engineering and Stochastic Modeling. Ed. by Paolo Ballarini et al. Lecture Notes in Computer Science. Cham: Springer International Publishing, 2021, pp. 425–440. DOI: 10.1007/978-3-030-91825-5_26

Stability

• Jean Mairesse and Pascal Moyal. "Stability of the Stochastic Matching Model". Journal of Applied Probability 53.4 (Dec. 2016), pp. 1064–1077. DOI: 10.1017/jpr.2016.65

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Performance evaluation

- Pascal Moyal, Ana Bušić, and Jean Mairesse. "A Product Form for the General Stochastic Matching Model". *Journal of Applied Probability* 58.2 (June 2021), pp. 449–468. DOI: 10.1017/jpr.2020.100
- Céline Comte. "Stochastic Non-Bipartite Matching Models and Order-Independent Loss Queues". *Stochastic Models* 38.1 (Jan. 2, 2022), pp. 1–36. DOI: 10.1080/15326349.2021.1962352

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- The Markov process is positive recurrent if and only if

$$\rho(\mathcal{I}) \triangleq \frac{\sum_{i \in \mathcal{I}} \mu_i}{\sum_{i \in \mathcal{V}(\mathcal{I})} \mu_i} < 1, \quad \mathcal{I} \in \mathbb{I}^+.$$



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In our toy example:

$$\begin{cases} \rho(\{1\}) = \frac{\mu_1}{\mu_2}, & \rho(\{2\}) = \frac{\mu_2}{\mu_1 + \mu_3 + \mu_4}, & \rho(\{3\}) = \frac{\mu_3}{\mu_2 + \mu_4}, \\ \rho(\{4\}) = \frac{\mu_4}{\mu_2 + \mu_3}, & \rho(\{1,3\}) = \frac{\mu_1 + \mu_3}{\mu_2 + \mu_4}, & \rho(\{1,4\}) = \frac{\mu_1 + \mu_4}{\mu_2 + \mu_3}. \end{cases}$$

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• This condition can be satisfied only if the graph is non-bipartite.

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- The set of unmatched items is an independent set, meaning that the state space is

$$\mathcal{C} = \bigcup_{\mathcal{I} \in \mathbb{I}} \mathcal{C}_{\mathcal{I}},$$

where $C_{\mathcal{I}}$ is the set of sequences $c \in \mathcal{V}^*$ where the set of classes is \mathcal{I} , for each $\mathcal{I} \in \mathbb{I}$



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• Stationary distribution [MM16]: for $(c_1, c_2, \ldots, c_\ell) \in \mathcal{C} \setminus \{\emptyset\}$,

$$\pi(c_1, c_2, \dots, c_{\ell}) = \pi(\emptyset) \prod_{p=1}^{\ell} \frac{\mu_{c_p}}{\sum_{i \in \mathcal{V}(\{c_1, c_2, \dots, c_p\})} \mu_i}.$$



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• In our toy example:

$$\pi((1, 1), (3), (1), (3)) = \pi(\emptyset) \frac{\mu_1}{\mu_2} \frac{\mu_1}{\mu_2} \frac{\mu_3}{\mu_2 + \mu_4} \frac{\mu_1}{\mu_2 + \mu_4} \frac{\mu_3}{\mu_2 + \mu_4}.$$

• In other words, for each $\mathcal{I}\in\mathbb{I}^{\scriptscriptstyle +}$ and $(c_1,c_2,\ldots,c_{\ell-1},c_\ell)\in\mathcal{C}_{\mathcal{I}}$,

$$\pi(c_1, c_2, \dots, c_{\ell-1}, c_\ell) = \pi(c_1, c_2, \dots, c_{\ell-1}) \frac{\mu_{c_\ell}}{\sum_{i \in \mathcal{V}(\mathcal{I})} \mu_i}.$$

Transitions out of and into a state



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"Partial" balance equations

• Balance between state c and states with -1 item:

$$\pi(c_1,\ldots,c_\ell)\left(\sum_{i\in\mathcal{V}(\mathcal{I})}\mu_i\right)=\pi(c_1,\ldots,c_{\ell-1})\mu_{c_\ell},$$

• Balance between state c and states with +1 class-i item:

$$\pi(c_1,\ldots,c_\ell)\mu_i = \sum_{p=1}^{\ell+1} \pi(c_1,\ldots,c_{p-1},i,c_p,\ldots,c_\ell) \left(\sum_{j\in\mathcal{V}_i\setminus\mathcal{V}(\{c_1,\ldots,c_{p-1}\})} \mu_j\right), \quad i\notin\mathcal{V}(\mathcal{I}).$$

What are product-form stationary distributions useful for?

- $\bullet\,$ Compute performance metrics $\rightarrow\,$ in the rest of this talk
- \bullet Analyze scaling regimes \rightarrow in the rest of this talk
- \bullet Optimization and learning \rightarrow preprint + ongoing work



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Variants of the model

- Bipartite vs. non-bipartite graph
- Abandonment (a.k.a. renegging)
- Admission control



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Normalization constant

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$$\pi(\mathcal{I}) = \sum_{c \in \mathcal{C}_{\mathcal{I}}} \pi(c), \quad \mathcal{I} \in \mathbb{I}.$$



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$$\pi(\mathcal{I}) = rac{
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• Key step:
$$\pi(c_1, \ldots, c_\ell, i) = \pi(c_1, \ldots, c_\ell) \frac{\mu_i}{\sum_{j \in \mathcal{V}(\mathcal{I})} \mu_j}, \ \mathcal{I} \in \mathbb{I}^+, \ i \in \mathcal{I}, \ (c_1, \ldots, c_\ell) \in \mathcal{C}_{\mathcal{I}} \cup \mathcal{C}_{\mathcal{I} \setminus \{i\}}.$$

Performance metrics

• Waiting probability of class *i*:

$$\omega_i = \sum_{\mathcal{I} \in \mathbb{I}: i \notin \mathcal{V}(\mathcal{I})} \pi(\mathcal{I}),$$



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• Mean queue length:

$$L = \sum_{\mathcal{I} \in \mathbb{I}^+} \ell(\mathcal{I}), \quad \text{where} \quad \ell(\mathcal{I}) = \frac{\pi(\mathcal{I})}{1 - \rho(\mathcal{I})} + \frac{\rho(\mathcal{I})}{1 - \rho(\mathcal{I})} \left(\sum_{i \in \mathcal{I}} \frac{\mu_i}{\sum_{j \in \mathcal{I}} \mu_j} \ell(\mathcal{I} \setminus \{i\}) \right).$$

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Similar formulas for the per-class performance.

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• The mean waiting time follows by applying Little's law.

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- The mean waiting time follows by applying Little's law.
- Similar results for the bipartite variant of the model [CD21].

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 - the mean number of unmatched items is $\sim \frac{\rho(\mathcal{I})}{1-\rho(\mathcal{I})}$.
- Take-away: **minimizing the maximum load** is likely a good heuristic to optimize performance.



$$\begin{array}{c} \mu_1 \longrightarrow \\ \mu_3 \longrightarrow \end{array} \begin{array}{c} 3 \\ 1 \\ 1 \\ \end{array} \begin{array}{c} \mu_2 + \mu_4 \end{array}$$

 $\simeq {\rm M}/{\rm M}/{\rm 1}$ multi-class queue
















Numerical results: Cycle with a chord



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• The complexity is "hidden" in the description of the state space and normalization constant.

Outline

Literature review

2 Basic results

- Stability condition
- Definition of the discrete-time Markov chain
- Product-form stationary distribution and partial balance

3 Performance analysis

- Normalization constant
- Performance metrics
- Heavy-traffic analysis

4 Concluding remarks

Take-away

- Product-form stationary distribution
- Closed-form expressions for performance metrics
- Heavy-traffic analysis



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- Product-form stationary distribution
- Closed-form expressions for performance metrics
- Heavy-traffic analysis

Future works

- Extensions to state-dependent arrival rates? hypergraphs? other policies?
- More fundamental relationship between balance, reversibility, and insensitivity?

