

LEARNING IN HARMONIC GAMES



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French National Center for Scientific Research (CNRS)

⟨ Workshop on Learning in Games | Toulouse | July 3, 2024 ⟩

About

- 📄 Legacci, M, & Pradeliski, *A geometric decomposition of finite games: Convergence vs. recurrence under exponential weights*, ICML 2024
- 📄 Legacci, M, Papadimitriou, Piliouras, & Pradeliski, *No-regret learning in harmonic games: Extrapolation in the presence of conflicting interests*, preprint, 2024



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B. Pradeliski

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D. Legacci



C. Papadimitriou



G. Piliouras



B. Pradeliski

Outline

- 1 What's a harmonic game?
- 2 No-regret learning
- 3 Learning in continuous time
- 4 Learning in discrete time

Finite games

Finite games

A **finite game** $\Gamma \equiv \Gamma(\mathcal{N}, \mathcal{A}, u)$ consists of

- ▶ A finite set of **players** $\mathcal{N} = \{1, \dots, N\}$
- ▶ A finite set of **actions** (or **pure strategies**) $\mathcal{A}_i = \{1, \dots, A_i\}$ per player $i \in \mathcal{N}$
- ▶ An ensemble of **payoff functions** $u_i: \mathcal{A} \equiv \prod_i \mathcal{A}_i \rightarrow \mathbb{R}, i \in \mathcal{N}$

Notation

- ▶ **Action profile:**

$$\alpha \equiv (\alpha_1, \dots, \alpha_N) \in \mathcal{A} := \prod_{i \in \mathcal{N}} \mathcal{A}_i$$

- ▶ **Pure** - or **realized** - **payoff** of player i

$$u_i(\alpha) \equiv u_i(\alpha_i; \alpha_{-i}) := u_i(\alpha_1, \dots, \alpha_N)$$

- ▶ **Pure** - or **realized** - **payoff vector** of player i

$$v_i(\alpha) \equiv v_i(\alpha_1, \dots, \alpha_N) := (u_i(\beta_i; \alpha_{-i}))_{\beta_i \in \mathcal{A}_i}$$

Mixed extensions

Players can *randomize*

- ▶ **Mixed strategy** of player i

$$x_i = (x_{i\alpha_i})_{\alpha_i \in \mathcal{A}_i} \in \Delta(\mathcal{A}_i) =: \mathcal{X}_i$$

$x_{i\alpha_i}$ = prob. that player i plays $\alpha_i \in \mathcal{A}_i$

- ▶ **Mixed payoff** of player i

$$u_i(x) = \mathbb{E}_{\alpha \sim x} [u_i(\alpha)] = \sum_{\alpha_1 \in \mathcal{A}_1} \dots \sum_{\alpha_N \in \mathcal{A}_N} x_{1,\alpha_1} \dots x_{N,\alpha_N} u_i(\alpha_1, \dots, \alpha_N)$$

- ▶ **Mixed payoff vector** of player i

$$v_i(x) := (u_i(\alpha_i; x_{-i}))_{\alpha_i \in \mathcal{A}_i}$$

vector of *expected* rewards

$v_i(x)$ only depends on x_{-i}

- ▶ **Mixed extension of Γ** : $\tilde{\Gamma} \equiv \Delta(\Gamma)$

Nash equilibrium

Nash equilibrium

[Nash, 1950]

“No player has an incentive to deviate from their chosen strategy if other players don't”

$$u_i(x_i^*; x_{-i}^*) \geq u_i(x_i; x_{-i}^*) \quad \text{for all } x_i \in \mathcal{X}_i, i \in \mathcal{N}$$

Nash equilibrium

Nash equilibrium

[Nash, 1950]

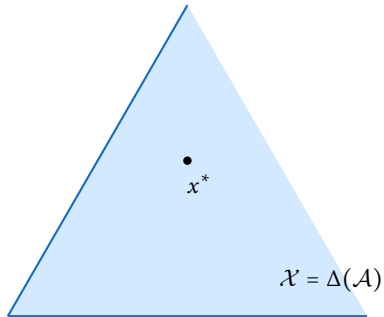
“No player has an incentive to deviate from their chosen strategy if other players don’t”

$$u_i(x_i^*; x_{-i}^*) \geq u_i(x_i; x_{-i}^*) \quad \text{for all } x_i \in \mathcal{X}_i, i \in \mathcal{N}$$

- ▶ **Variational characterization:** $\langle v(x^*), x - x^* \rangle \leq 0$ for all $x \in \mathcal{X}$ # Stampacchia variational inequality
- ▶ **Pure equilibrium:** $\text{supp}(x^*) = \text{singleton}$ # pure strategy profile
- ▶ **Strict equilibrium:** “>” instead of “≥” # unique best response; necessarily pure

Equilibrium configurations

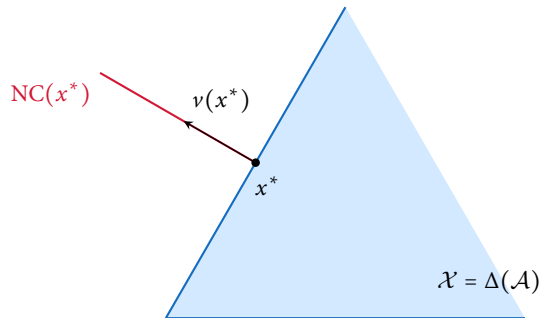
Figure: Different equilibrium configurations: *fully mixed*



#NC(x^*) = normal cone at x^* (outward normal directions)

Equilibrium configurations

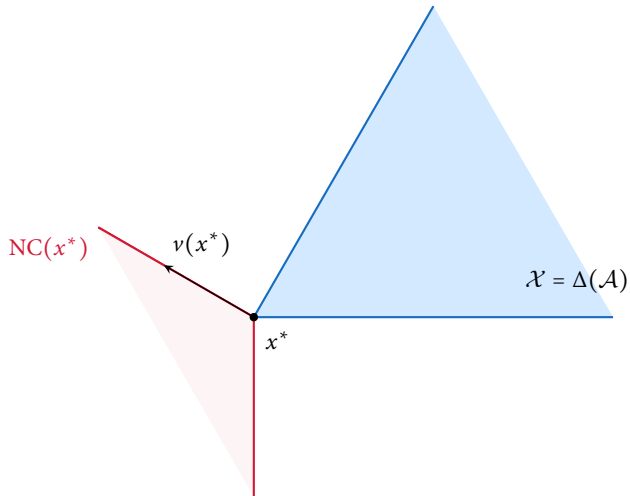
Figure: Different equilibrium configurations: fully mixed vs. *mixed*



$NC(x^*)$ = normal cone at x^* (outward normal directions)

Equilibrium configurations

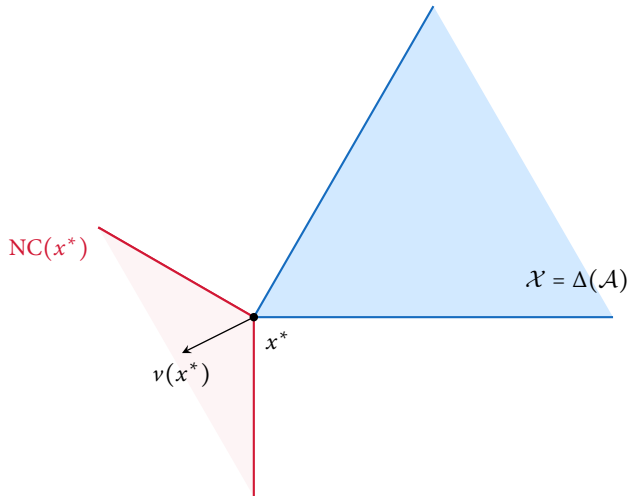
Figure: Different equilibrium configurations: fully mixed vs. mixed vs. *pure*



$NC(x^*)$ = normal cone at x^* (outward normal directions)

Equilibrium configurations

Figure: Different equilibrium configurations: fully mixed vs. mixed vs. pure vs. *strict*



$\# NC(x^*) =$ normal cone at x^* (outward normal directions)

Potential games

Potential games

[Monderer & Shapley, 1996]

The game $\Gamma \equiv \Gamma(\mathcal{N}, \mathcal{A}, u)$ is a **potential game** if it admits a **potential function** $\Phi: \mathcal{A} \rightarrow \mathbb{R}$ such that

$$u_i(\alpha_i; \alpha_{-i}) - u_i(\beta_i; \alpha_{-i}) = \Phi(\alpha_i; \alpha_{-i}) - \Phi(\beta_i; \alpha_{-i}) \quad \text{for all } \alpha, \beta \in \mathcal{A}$$

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Basic properties

- ▶ Player interests **aligned** along a common objective # common interest
- ▶ Improvement paths **always** terminate # no best-response cycles
- ▶ **Always admit pure equilibria** # generically strict

Graphical representation

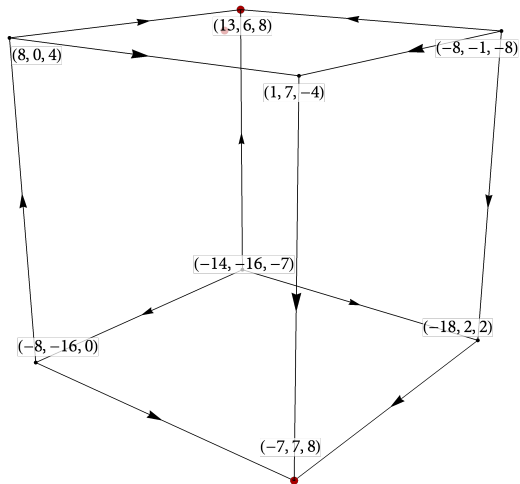


Figure: Response graph of a potential game (Nash in red)

Graphical representation

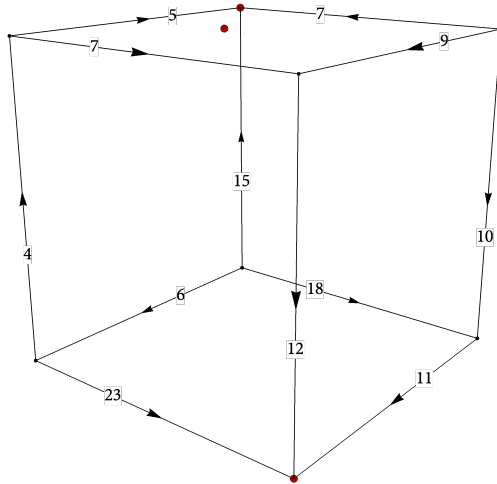


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Harmonic games

Harmonic games

[Candogan et al., 2011; Abdou et al., 2022]

The game $\Gamma \equiv \Gamma(\mathcal{N}, \mathcal{A}, u)$ is a **harmonic game** if it admits a **harmonic measure** $\mu: \prod_i \mathcal{A}_i \rightarrow (0, \infty)$ such that

$$\sum_{i \in \mathcal{N}} \sum_{\beta \in \mathcal{A}_i} \mu_i \beta_i [u_i(\alpha_i; \alpha_{-i}) - u_i(\beta_i; \alpha_{-i})] = 0 \quad \text{for all } \alpha \in \mathcal{A}$$

☞ **uniform harmonic**: $\mu_i = \text{unif}_{\mathcal{A}_i}$

☞ **flow conservation**: no deviation sources or sinks

☞ terminology: harmonic component of Hodge decomposition

Harmonic games

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☞ **uniform harmonic**: $\mu_i = \text{unif}_{\mathcal{A}_i}$

☞ **flow conservation**: no deviation sources or sinks

☞ terminology: harmonic component of Hodge decomposition

Basic properties

- ▶ Player interests **anti-aligned** # conflicts of interest
- ▶ Improvement paths **never** terminate # best-response sinks
- ▶ **No pure equilibria** # at least one fully mixed

Graphical representation

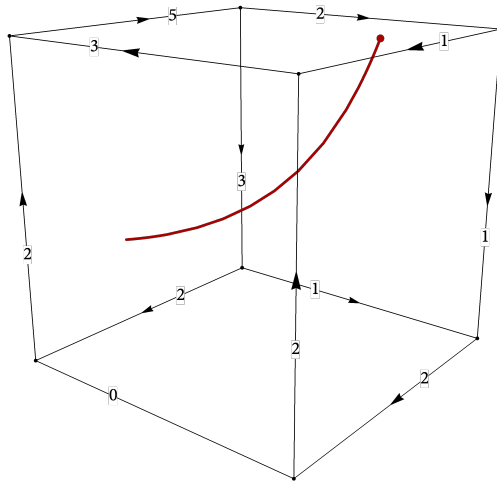




Figure: Response graph of a harmonic game (Nash in red)

Harmonic resonance

Examples of harmonic games

- ▶ Matching Pennies, Rock-Paper-Scissors, Dawkins' battle of the sexes, ...
- ▶ Two-player zero-sum games with a fully mixed equilibrium
- ▶ Cyclic games
- ▶ ...


 Legacci et al. (2024)


 Hofbauer & Schlag (2000)

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Hodge decomposition of games

[Candogan et al., 2011; Abdou et al., 2022]

Any finite game Γ can be decomposed as

$$\Gamma = \Gamma_{\text{pot}} + \Gamma_{\text{harm}}$$

where Γ_{pot} is potential and Γ_{harm} is harmonic

Harmonic resonance

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Remarks:

- ▶ Decomposition **not unique**
- ▶ Harmonic and potential games are **orthogonal**
- ▶ **Harmonic \neq zero-sum** ⚠

must fix measure / gauge

given a measure / metric

zero-sum can be potential

Harmonic resonance

Hodge decomposition of games

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- ▶ **Harmonic \neq zero-sum** ⚠ # zero-sum can be potential

Harmonic games \rightsquigarrow strategic complement of potential games

This talk

Are harmonic and potential games complementary from a dynamic viewpoint?

What is the long-run behavior of no-regret learning in harmonic games?

Outline

- ① What's a harmonic game?
- ② No-regret learning
- ③ Learning in continuous time
- ④ Learning in discrete time

Multi-agent learning

Sequence of events - generic

for each <i>epoch</i> and every <i>player</i> do	# continuous / discrete
Choose <i>action</i>	# continuous / finite
Receive <i>reward</i>	# endogenous / exogenous
Get <i>feedback</i> (maybe)	# full info / oracle / payoff-based
end for	

Defining elements

- ▶ **Time:** continuous or discrete?
- ▶ **Players:** continuous or finite?
- ▶ **Actions:** continuous or finite?
- ▶ **Rewards:** endogenous or exogenous (determined by other players or by “Nature”)?
- ▶ **Feedback:** observe other actions / other rewards / only received?

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end for	

Defining elements

- ▶ **Time:** *continuous* and *discrete*
- ▶ **Players:** ~~continuous~~ or *finite*
- ▶ **Actions:** ~~continuous~~ or *finite*
- ▶ **Rewards:** *endogenous* or *exogenous* (determined by other players or by "Nature")
- ▶ **Feedback:** observe other actions / other rewards / only received?

Regret minimization

Individual regret

$$\text{Reg}_i(T) = \max_{\alpha_i \in \mathcal{A}_i} \sum_{t=1}^T [u_i(\alpha_i; z_{-i,t}) - u_i(z_t)]$$

Regret minimization

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No regret: $\text{Reg}_i(T) = o(T)$

the smaller the better

“The chosen policy is as good as the best fixed strategy in hindsight.”

Worst-case guarantee: at the very least, minimize regret

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Worst-case guarantee: at the very least, minimize regret

Literature:

- ▶ Economics 📖 Hannan (1957); Hart & Mas-Colell (2000); Fudenberg & Levine (1998)
- ▶ Mathematics 📖 Blackwell (1956); Lai & Robbins (1985); Sorin (2024)
- ▶ Computer science 📖 Shalev-Shwartz (2011); Cesa-Bianchi & Lugosi (2006); Lattimore & Szepesvári (2020)

Regret minimization and rationality

Individual regret

$$\text{Reg}_i(T) = \max_{\alpha_i \in \mathcal{A}_i} \sum_{t=1}^T [u_i(\alpha_i; z_{-i,t}) - u_i(z_t)]$$

Does no-regret learning converge to equilibrium?

Regret minimization and rationality

Individual regret

$$\text{Reg}_i(T) = \max_{\alpha_i \in \mathcal{A}_i} \sum_{t=1}^T [u_i(\alpha_i; z_{-i,t}) - u_i(z_t)]$$

Under no-regret learning, **empirical frequencies of play** converge to equilibrium

Regret minimization and rationality

Individual regret

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Under no-regret learning, **empirical frequencies of play** converge to the game's **Hannan set / set of coarse correlated equilibria**

Regret minimization and rationality

Individual regret

$$\text{Reg}_i(T) = \max_{\alpha_i \in \mathcal{A}_i} \sum_{t=1}^T [u_i(\alpha_i; z_{-i,t}) - u_i(z_t)]$$

Under no-regret learning, **empirical frequencies of play** converge to the game's **Hannan set / set of coarse correlated equilibria**

Empirical frequencies of play

$$z_{\alpha,t} = \frac{1}{t} \sum_{\tau=1}^t \mathbb{1}\{\alpha_\tau = \alpha\}$$

Coarse correlated equilibrium

A correlated strategy $z \in \Delta(\mathcal{A})$ is a **coarse correlated equilibrium / Hannan consistent** if

$$u_i(z) \geq u_i(\beta_i; z_{-i}) \quad \text{for all } \beta_i \in \mathcal{A}_i, i \in \mathcal{N}$$

The VZ game

A *coordination game*

	<i>A</i>	<i>B</i>
<i>A</i>	(1, 1)	(0, 0)
<i>B</i>	(0, 0)	(1, 1)

The VZ game

A coordination game with two feeble twins

Viossat & Zapechelnyuk, 2013

	A	A_-	B	B_-
A	$(1, 1)$	$(1, 2/3)$	$(0, 0)$	$(0, -1/3)$
A_-	$(2/3, 1)$	$(2/3, 2/3)$	$(-1/3, 0)$	$(-1/3, -1/3)$
B	$(0, 0)$	$(0, -1/3)$	$(1, 1)$	$(1, 2/3)$
B_-	$(-1/3, 0)$	$(-1/3, -1/3)$	$(2/3, 1)$	$(2/3, 2/3)$

Feeble twins:

- ▶ A_- is **strictly 1/3-dominated** by A
- ▶ B_- is **strictly 1/3-dominated** by B

The VZ game

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A ₋	(2/3, 1)	(2/3, 2/3)	(-1/3, 0)	(-1/3, -1/3)
B	(0, 0)	(0, -1/3)	(1, 1)	(1, 2/3)
B ₋	(-1/3, 0)	(-1/3, -1/3)	(2/3, 1)	(2/3, 2/3)

Feeble twins:

- ▶ A₋ is **strictly 1/3-dominated** by A
- ▶ B₋ is **strictly 1/3-dominated** by B

BUT!

- ▶ Suppose players play (A₋, A₋) and (B₋, B₋) each with prob. 1/2
- ▶ **Distribution of play is a CCE:** $u_i(\alpha_i; z_{-i}) - u_i(z) \leq -1/6$
- ▶ **No regret!**

in fact, **negative** regret

The VZ game

A coordination game with two feeble twins

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	A	A ₋	B	B ₋
A	(1, 1)	(1, 2/3)	(0, 0)	(0, -1/3)
A ₋	(2/3, 1)	(2/3, 2/3)	(-1/3, 0)	(-1/3, -1/3)
B	(0, 0)	(0, -1/3)	(1, 1)	(1, 2/3)
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No-regret play may lead to playing dominated strategies for all time!

Outline

- ① What's a harmonic game?
- ② No-regret learning
- ③ Learning in continuous time
- ④ Learning in discrete time

Game-theoretic learning

Sequence of events – continuous time

Require: finite game $\Gamma \equiv \Gamma(\mathcal{N}, \mathcal{A}, u)$

repeat

At each epoch $t \geq 0$ **do simultaneously** for all players $i \in \mathcal{N}$

continuous time

Choose **mixed strategy** $x_i(t) \in \mathcal{X}_i := \Delta(\mathcal{A}_i)$

mixing

Get **mixed payoff** $u_i(x(t)) = \langle v_i(x(t)), x_i(t) \rangle$

payoff phase

Observe **mixed payoff vector** $v_i(x_t)$

feedback phase

until end

Defining elements

- ▶ **Time:** $t \geq 0$
- ▶ **Players:** many (finite)
- ▶ **Actions:** finite
- ▶ **Payoffs:** endogenous
- ▶ **Feedback:** mixed payoff vectors

multi-agent learning

game-theoretic learning

Learning with exponential weights

Exponential reinforcement mechanism:

- ▶ Score each action based on its cumulative payoff over time:

$$y_{i\alpha_i}(t) = \int_0^t v_{i\alpha_i}(x(\tau)) d\tau$$

- ▶ Play an action with probability exponentially proportional to its score

$$x_{i\alpha_i}(t) \propto \exp(y_{i\alpha_i}(t))$$

Exponential weight dynamics

[Littlestone & Warmuth, 1994; Auer et al., 1995]

$$\dot{y}_i = v_i(x) \quad x_i = \Lambda(y_i) := \frac{\exp(y_i)}{\|\exp(y_i)\|_1} \quad (\text{EWD})$$

Learning with exponential weights

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The replicator dynamics

[Taylor & Jonker, 1978]

$$\dot{x}_{i\alpha_i} = x_{i\alpha_i} [u_i(\alpha_i; x_{-i}) - u_i(x)] \quad (\text{RD})$$

What do the dynamics look like?

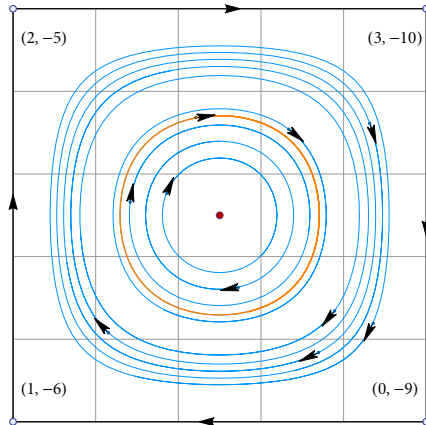


Figure: The replicator dynamics in a 2×2 harmonic game (Nash in red)

What do the dynamics look like?

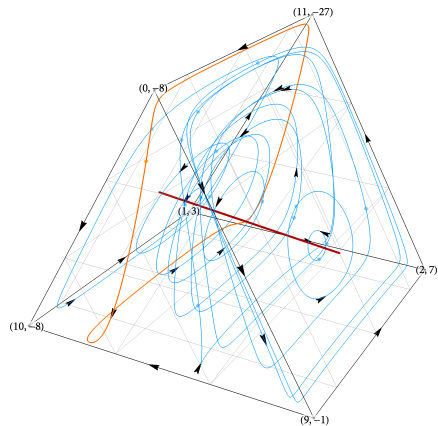


Figure: The replicator dynamics in a 2×3 harmonic game (Nash in red)

What do the dynamics look like?

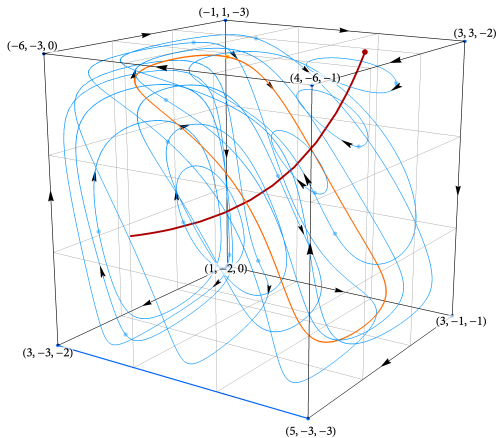


Figure: The replicator dynamics in a $2 \times 2 \times 2$ harmonic game (Nash in red)

What do the dynamics look like?

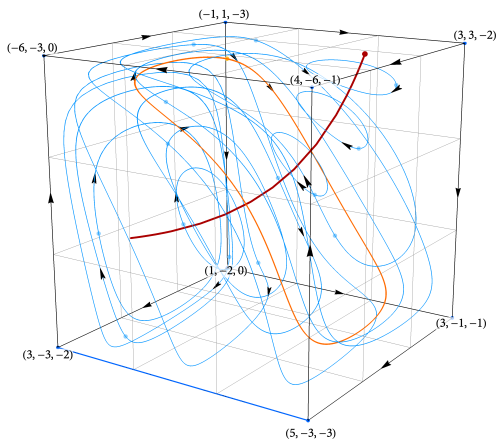


Figure: The replicator dynamics in a $2 \times 2 \times 2$ harmonic game (Nash in red)

Trajectories always periodic!

What's going on? [Geometric version]

- ▶ Dynamic behavior “orthogonal” to potential games
- ▶ Is there a dynamic version of the potential/harmonic decomposition?

convergence to Nash

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convergence to Nash

Geometric decomposition of finite games

[Legacci et al., 2024]

- 1 Any finite game can be decomposed as

$$v(x) = F(x) + B(x)$$

where F is **irrotational** and B is **incompressible** under the Shahshahani metric on \mathcal{X}

- 2 A game is **incompressible** if and only if it is **uniform harmonic**

What's going on? [Geometric version]

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Remarks:

- ▶ **Shahshahani metric** \rightsquigarrow replicator-compatible geometric structure on \mathcal{X}

$$\# g_{\alpha\beta}(x) = \delta_{\alpha\beta}/x_{\alpha}$$

🔥 **Why uniform?**

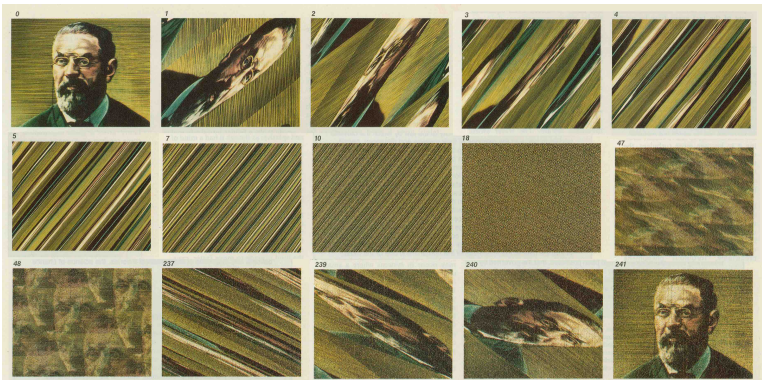
⚠️ highly surprising structural match!

Quasi-periodicity

Poincaré recurrence

A dynamical system is **Poincaré recurrent** if almost all solution trajectories return **arbitrarily close** to their starting point **infinitely many times**

formal definition of "quasi-periodicity"



Incompressibility, recurrence, and harmonic games

Volume preservation

If Γ is incompressible / harmonic, the replicator dynamics are volume-preserving under the Shahshahani metric:

$$\text{vol}_{\text{Shah}}(A) = \text{vol}_{\text{Shah}}(\text{RD}_t(A)) \quad \text{for every measurable set of initial conditions } A \subseteq \mathcal{X}$$

Incompressibility, recurrence, and harmonic games

Volume preservation

If Γ is incompressible / harmonic, the replicator dynamics are volume-preserving under the Shahshahani metric:

$$\text{vol}_{\text{Shah}}(A) = \text{vol}_{\text{Shah}}(\text{RD}_t(A)) \quad \text{for every measurable set of initial conditions } A \subseteq \mathcal{X}$$

Poincaré recurrence [Legacci et al., 2024]

In any uniform harmonic game, the replicator dynamics are Poincaré recurrent.

Exponential weights redux

Exponential weight dynamics

$$\dot{y}_i = v_i(x) \quad x_i = \Lambda(y_i) = \frac{\exp(y_i)}{\|\exp(y_i)\|_1} \quad (\text{EWD})$$

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Softmax interpretation

$$x = \Lambda(y) \quad \iff \quad x = \arg \max_{z \in \mathcal{X}} \left\{ \langle y, z \rangle - \underbrace{\sum_{\alpha \in \mathcal{A}} z_\alpha \log z_\alpha}_{\text{entropic penalty}} \right\}$$

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Is there a general principle in play?

Regularized learning

Replace the “**leader**” / hard arg max

$$x_i(t) = \arg \max_{z_i \in \mathcal{X}} \underbrace{\int_0^t u_i(z_i; x_{-i}(\tau)) d\tau}_{\text{cumulative payoff}}$$

with a “**regularized leader**” / soft arg max:

$$x_i(t) = \arg \max_{z_i \in \mathcal{X}} \left\{ \underbrace{\int_0^t u_i(z_i; x_{-i}(\tau)) d\tau}_{\text{cumulative payoff}} - \underbrace{h_i(z_i)}_{\text{penalty}} \right\}$$

where $h_i: \mathcal{X}_i \rightarrow \mathbb{R}$ is a strongly convex **regularizer** on \mathcal{X}_i

Regularized learning

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where $h_i: \mathcal{X}_i \rightarrow \mathbb{R}$ is a strongly convex **regularizer** on \mathcal{X}_i

Follow-the-regularized-leader (FTRL)

$$\dot{y}_i(t) = v_i(x(t)) \quad x_i(t) = Q_i(y_i(t)) \tag{FTRL-D}$$

where

$$Q_i(y_i) = \arg \max_{z_i \in \mathcal{X}_i} \{ \langle y_i, z_i \rangle - h_i(z_i) \} \tag{\# regularized choice / best response map}$$

Example: Projection dynamics

Euclidean / Ridge regularization

Regularizer:

$$h(x) = \frac{1}{2} \sum_{\alpha} x_{\alpha}^2$$

Choice map:

$$\Pi(y) = \arg \max_{x \in \mathcal{X}} \{ \langle y, x \rangle - (1/2) \|x\|_2^2 \} = \arg \min_{x \in \mathcal{X}} \|y - x\| = \text{proj}_{\mathcal{X}}(y)$$

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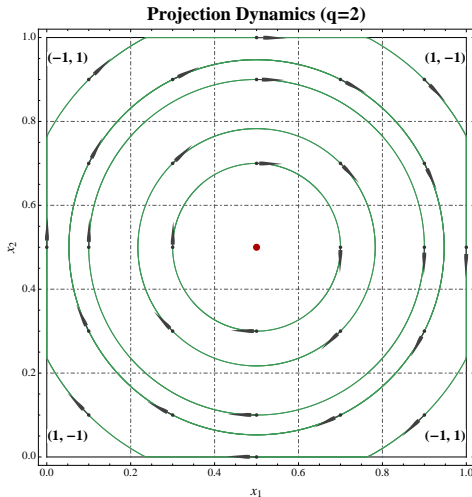
Projection dynamics

[Friedman, 1991; M & Sandholm, 2016]

$$\dot{y}_i(t) = v_i(x(t)) \quad x_i(t) = \Pi(y_i(t)) \quad (\text{PD})$$

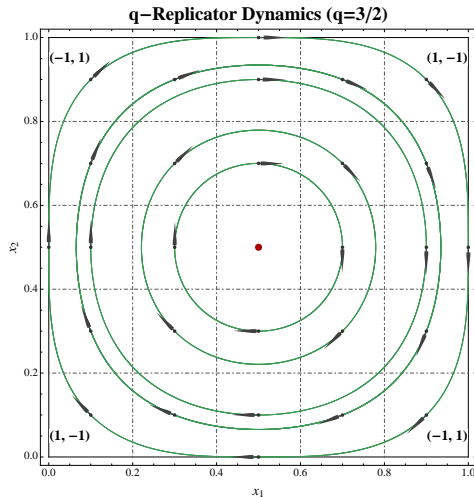
What do the dynamics look like?

The Tsallis kernel: $h(x) = [q(1 - q)]^{-1} \sum_{\alpha} (x_{\alpha} - x_{\alpha}^q)$



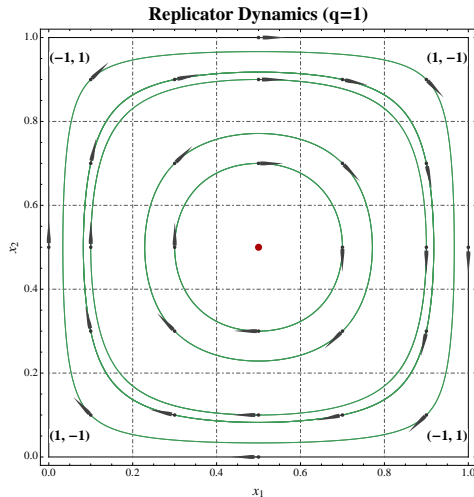
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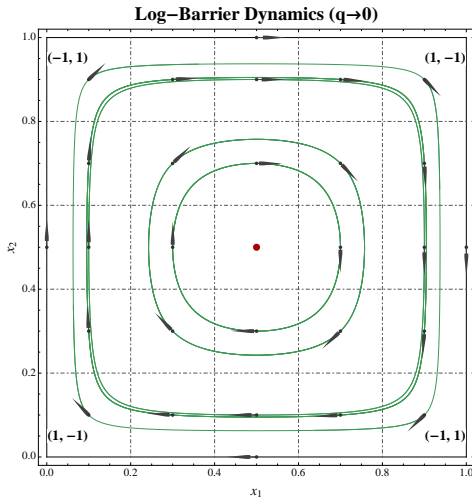
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What's going on? [Dual version]

Poincaré recurrence [Legacci et al., 2024]

The dynamics of FTRL are Poincaré recurrent in any harmonic game

What's going on? [Dual version]

Poincaré recurrence [Legacci et al., 2024]

The dynamics of FTRL are Poincaré recurrent in any harmonic game

Remarks:

- ⚠ **No geometric compatibility** \implies requires completely different proof technique
 - ▶ Leverage tools from convex analysis \rightsquigarrow constant of motion
 - ▶ Simultaneously extend to all harmonic measures and all regularizers

Outline

- ① What's a harmonic game?
- ② No-regret learning
- ③ Learning in continuous time
- ④ Learning in discrete time

Game-theoretic learning

Sequence of events – discrete time

Require: finite game $\Gamma \equiv \Gamma(\mathcal{N}, \mathcal{A}, u)$

repeat

At each epoch $t = 1, 2, \dots$ **do simultaneously** for all players $i \in \mathcal{N}$ # discrete time

Choose **mixed strategy** $x_{i,t} \in \mathcal{X}_i := \Delta(\mathcal{A}_i)$ # mixing

Choose **action** $\alpha_{i,t} \sim x_{i,t}$ and get **realized payoff** $u_i(\alpha_{i,t}; \alpha_{-i,t})$ # payoff phase

Observe **mixed payoff vector** $v_i(x_t) = (u_i(\alpha_i; x_{-i,t}))_{\alpha_i \in \mathcal{A}_i}$ # feedback phase

until end

Defining elements

- ▶ **Time:** $t = 1, 2, \dots$
- ▶ **Players:** many (finite) # multi-agent learning
- ▶ **Actions:** finite
- ▶ **Payoffs:** endogenous # game-theoretic learning
- ▶ **Feedback:** **mixed payoff vectors** # full information, exact

Follow-the-regularized-leader

Follow-the-regularized-leader (FTRL)

$$y_{i,t+1} = y_{i,t} + \gamma_t v_t$$

$$x_{i,t+1} = Q_i(y_{i,t+1}) \equiv \arg \max_{x_i \in \mathcal{X}} \{ \langle y_{i,t+1}, x_i \rangle - h_i(x_i) \} \quad (\text{FTRL})$$

- ▶ **Regularized best responses** instead of logit choice map 📄 M & Sandholm (2016)
- ▶ Every player's **regularizer** $h_i: \mathcal{X}_i \rightarrow \mathbb{R}$ is continuous and strongly convex on \mathcal{X}_i

$$h_i(x'_i) \geq h_i(x_i) + \langle \nabla h_i(x_i), x'_i - x_i \rangle + (K_i/2) \|x'_i - x_i\|^2$$

- ▶ **Template includes:** exponential weights, (lazy) projected gradient ascent, Tsallis-based algorithms, ...

Examples

Example (Ridge regularization)

- ▶ Regularizer:

$$h(x) = \frac{1}{2} \|x\|^2$$

- ▶ Algorithm:

$$y_{t+1} = y_t + \gamma_t v_t \quad x_{t+1} = \Pi_{\mathcal{X}}(y_{t+1})$$

Example (Entropic regularization)

- ▶ Regularizer:

$$h(x) = \sum_{\alpha \in \mathcal{A}} x_{\alpha} \log x_{\alpha}$$

- ▶ Algorithm:

$$y_{t+1} = y_t + \gamma_t v_t \quad x_{t+1} = \Lambda(y_{t+1})$$

Non-convergence of FTRL

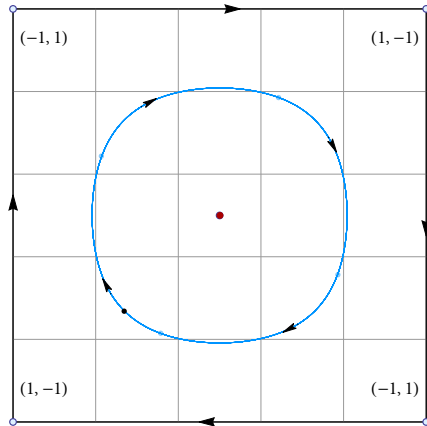


Figure: The replicator dynamics in Matching Pennies

Non-convergence of FTRL

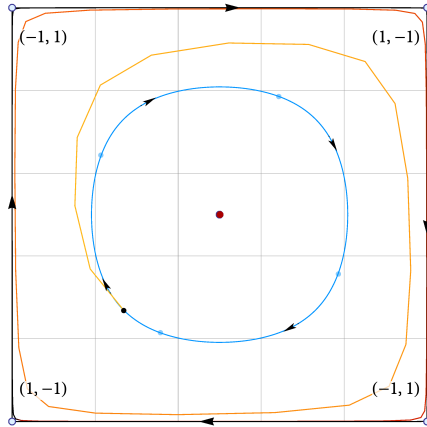


Figure: The FTRL algorithm in Matching Pennies

Non-convergence of FTRL

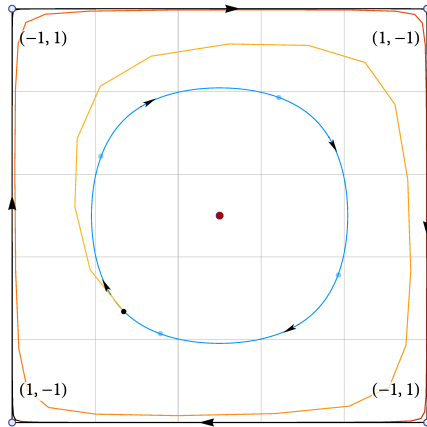


Figure: The FTRL algorithm in Matching Pennies

FTRL does not converge in harmonic games

FTRL with an extrapolation step

Extrapolated FTRL (FTRL+)

- a) **Extrapolation phase:** $y_{i,t+1/2} = y_{i,t} + \eta_i v_{i,t}$ $x_{i,t+1/2} = Q_i(y_{i,t+1/2})$
- b) **Update phase:** $y_{i,t+1} = y_{i,t} + \eta_i v_{i,t+1/2}$ $x_{i,t} = Q_i(y_{i,t+1})$ (FTRL+)

FTRL with an extrapolation step

Extrapolated FTRL (FTRL+)

$$\begin{array}{ll}
 \text{a) Extrapolation phase:} & y_{i,t+1/2} = y_{i,t} + \eta_i v_{i,t} & x_{i,t+1/2} = Q_i(y_{i,t+1/2}) \\
 \text{b) Update phase:} & y_{i,t+1} = y_{i,t} + \eta_i v_{i,t+1/2} & x_{i,t} = Q_i(y_{i,t+1})
 \end{array} \tag{FTRL+}$$

Payoff model

$$\begin{aligned}
 v_{i,t} &= \\
 v_{i,t+1/2} &=
 \end{aligned}$$

FTRL with an extrapolation step

Extrapolated FTRL (FTRL+)

a) <i>Extrapolation phase:</i>	$y_{i,t+1/2} = y_{i,t} + \eta_i v_{i,t}$	$x_{i,t+1/2} = Q_i(y_{i,t+1/2})$	(FTRL+)
b) <i>Update phase:</i>	$y_{i,t+1} = y_{i,t} + \eta_i v_{i,t+1/2}$	$x_{i,t} = Q_i(y_{i,t+1})$	

Payoff model

$$v_{i,t} =$$

$$v_{i,t+1/2} = v_i(x_{t+1/2})$$

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 \end{array}
 \tag{FTRL+}$$

Payoff model

$$\begin{aligned}
 v_{i,t} &= 0 \\
 v_{i,t+1/2} &= v_i(x_{t+1/2})
 \end{aligned}$$

FTRL with an extrapolation step

Extrapolated FTRL (FTRL+)

$$\begin{array}{ll} \text{a) Extrapolation phase:} & y_{i,t+1/2} = y_{i,t} + \eta_i v_{i,t} & x_{i,t+1/2} = Q_i(y_{i,t+1/2}) \\ \text{b) Update phase:} & y_{i,t+1} = y_{i,t} + \eta_i v_{i,t+1/2} & x_{i,t} = Q_i(y_{i,t+1}) \end{array} \quad (\text{FTRL+})$$

Payoff model

$$\begin{aligned} v_{i,t} &= v_i(x_t) \\ v_{i,t+1/2} &= v_i(x_{t+1/2}) \end{aligned}$$

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Payoff model

$$\begin{aligned}
 v_{i,t} &= v_i(x_{t-1/2}) \\
 v_{i,t+1/2} &= v_i(x_{t+1/2})
 \end{aligned}$$

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 \end{array} \tag{FTRL+}$$

Payoff model

$$\begin{aligned}
 v_{i,t} &= \lambda_i v_i(x_t) + (1 - \lambda_i) v_i(x_{t-1/2}) \\
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FTRL with an extrapolation step

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Payoff model

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 v_{i,t+1/2} &= v_i(x_{t+1/2})
 \end{aligned}$$

Examples

- ▶ $\lambda_i = 1$: Mirror-prox / Extra-gradient 📖 Korpelevich, 1976; Nemirovski, 2004
- ▶ $\lambda_i = 0$: Optimistic FTRL / Popov updates 📖 Popov, 1980; ?

Long-run behavior of FTRL+

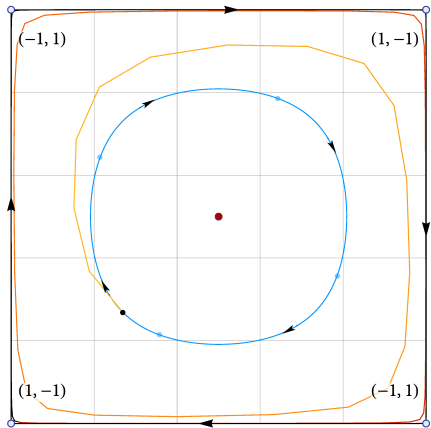


Figure: FTRL in Matching Pennies ✗

Long-run behavior of FTRL+

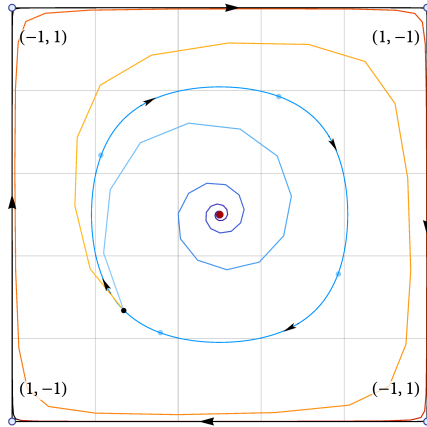


Figure: Mirror-Prox in Matching Pennies ✓

Long-run behavior of FTRL+

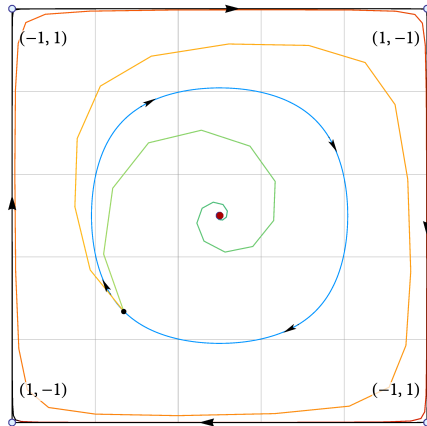


Figure: Optimistic FTRL in Matching Pennies ✓

Long-run behavior of FTRL+

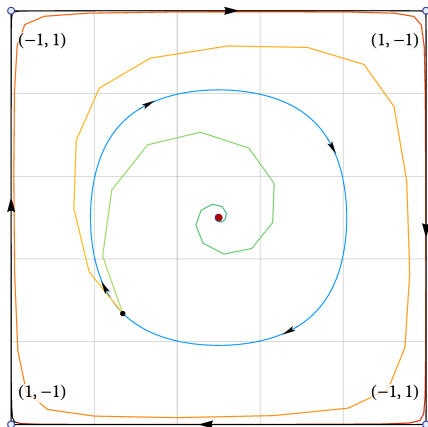


Figure: Optimistic FTRL in Matching Pennies ✓

Does (FTRL+) converge in harmonic games?

Long-run behavior of FTRL+

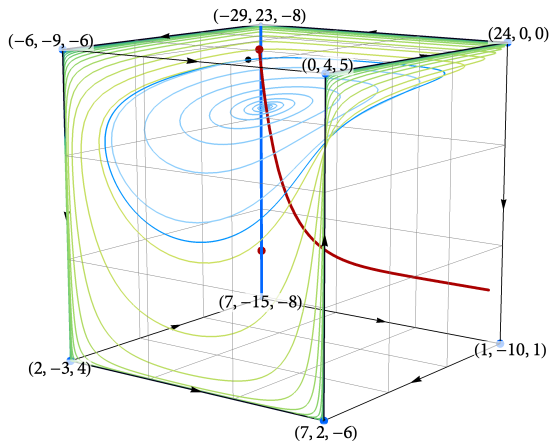


Figure: (FTRL+) in a $2 \times 2 \times 2$ harmonic game ✓

Does (FTRL+) converge in harmonic games?

The long-run behavior of FTRL+

Guarantee 1: Constant regret

[Legacci et al., 2024]

Assume:

- ▶ Γ is μ -harmonic
- ▶ Each player follows (FTRL+) with $\eta_i \leq m_i K_i [2(N+2) \max_j m_j L_j]^{-1}$

Then: (FTRL+) enjoys the bound

$$\text{Reg}_i(T) \leq \frac{H_i}{\eta_i} + \frac{2L_i}{N+2} \sum_{j \in \mathcal{N}} \frac{H_j}{\eta_j L_j} = \mathcal{O}(1)$$

where $H_i = \max h_i - \min h_i$, and L_i is the Lipschitz modulus of v_i

The long-run behavior of FTRL+

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Guarantee 2: Convergence

[Legacci et al., 2024]

Assume:

- ▶ Γ is μ -harmonic
- ▶ Each player follows (FTRL+) with $\eta_i \leq m_i K_i [2(N+2) \max_j m_j L_j]^{-1}$

Then: the sequence x_t generated by (FTRL+) converges to a Nash equilibrium

Take-aways and conclusions

Main take-aways:

- ▶ Harmonic games behave “orthogonally” to potential games in terms of learning
- ▶ No-regret learning in continuous time is recurrent
- ▶ No-regret learning in discrete time may be divergent...

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- ▶ Harmonic games behave “orthogonally” to potential games in terms of learning
- ▶ No-regret learning in continuous time is recurrent
- ▶ No-regret learning in discrete time may be divergent...
- ▶ ...but an extrapolation step recovers convergence to Nash equilibrium
- ▶ ...and guarantees constant regret

This is just a first peek:

- ▶ Rate of convergence?
- ▶ Inexact / Payoff-based information
- ▶ Adaptive / Agnostic step-size policies

difficult, but not hopeless

two-step policies?

AdaGrad-like?

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