Game theory for cumulative-prospect-theoretic agents

Venkat Anantharam based on the thesis work of Soham Phade University of California, Berkeley

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Soham Phade



Network economics today



Broadband network



Cloud Computing



Smart Grid



Ad Auctions



Telemedicine



Labor Markets

How to incorporate human preferences?

Decision Making under Uncertainty

Decision Maker = Choose between different Lotteries

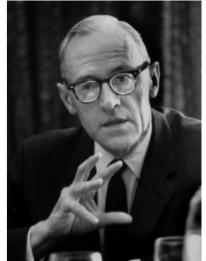
$L_2 =$	probability	0.15	0.05	0.6	0.2
11 2	outcome	8	5	-1	-2



Expected Utility Theory (EUT)

(Von Neumann-Morgenstern 1947)





Lottery

$$U(L) = 0.1u(10) + 0.2u(5) + 0.15u(2) + 0.1u(0) + 0.25u(-1) + 0.2u(-3)$$

Expected utility of lottery L Outcome χ +0.25u(-1) + 0.2u(-3)

Lottery with higher Expected Utility is preferred.

Utility function

utility

u(x)



Allais

Experiment 1			
Gamble 1A Gamble 1B			
Winnings Chance		Winnings	Chance
\$1 million	100%	\$1 million	89%
		Nothing	1%
		\$5 million	10%



Allais



Allais

Experiment 1			
Gamble 1A Gamble 1B			
Winnings Chance		Winnings	Chance
\$1 million	100%	\$1 million	89%
		Nothing	1%
		\$5 million	10%

Experiment 2				
Gamble 2A Gamble 2B				
Winnings Chance		Winnings	Chance	
Nothing	89%	Nothing	90%	
\$1 million	11%			
		\$5 million	10%	



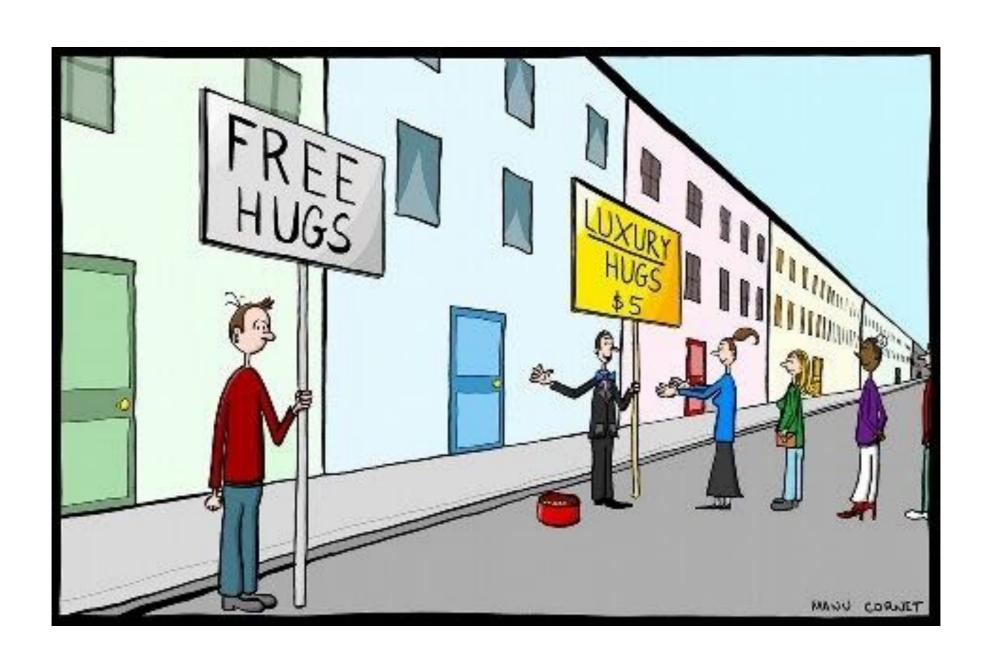
Allais

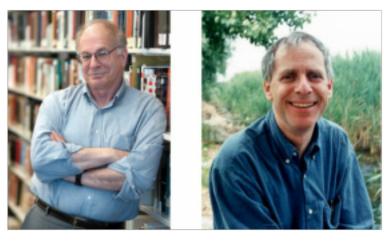
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People often do NOT follow EUT!

Behavioral Aspects





Kahneman

Tversky

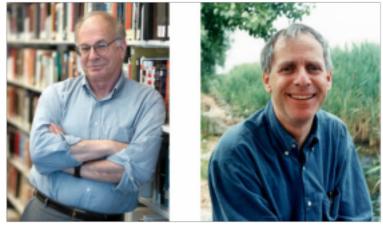
Kahneman - Tversky Cumulative Prospect Theory (1992)

Reference Point

$$r \in \mathbb{R}$$

Gains : Outcomes $\geq r$

Losses: Outcomes < r



Kahneman

Tversky

Kahneman - Tversky Cumulative Prospect Theory (1992)

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$$r \in \mathbb{R}$$

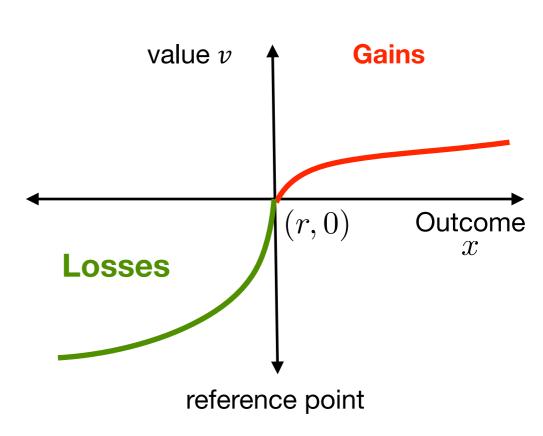
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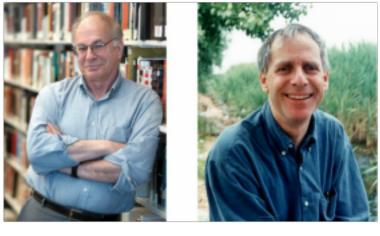
Losses: Outcomes < r

Value Function

$$v: \mathbb{R} \to \mathbb{R}$$

- 1. v(x) is continuous in x
- 2. v(r) = 0
- 3. it is strictly increasing in x





Kahneman

Tversky

Kahneman - Tversky Cumulative Prospect Theory (1992)

Probability Weighting Functions

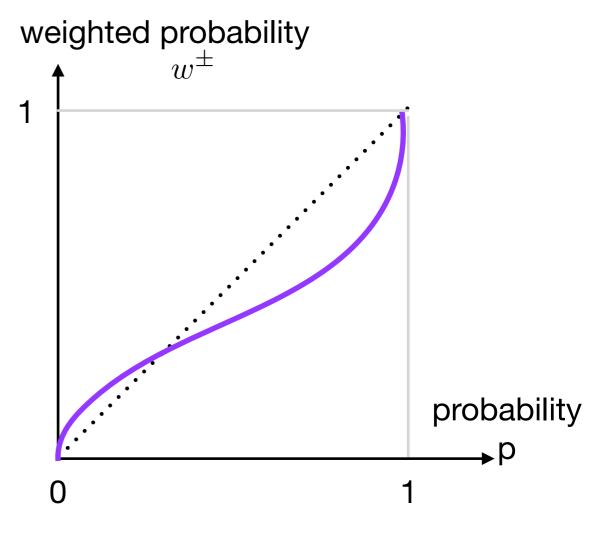
$$w^+:[0,1]\to[0,1]$$

$$w^+:[0,1]\to[0,1]$$
 $w^-:[0,1]\to[0,1]$

Gains

Losses

- 1. they are continuous
- 2. they are strictly increasing
- **3.** $w^{\pm}(0) = 0$ and $w^{\pm}(1) = 1$



Suppose
$$r=0$$
 Gains Losses $L= \begin{bmatrix} \text{probability} & 0.1 & 0.2 & 0.15 & 0.1 & 0.25 & 0.2 \\ \text{outcome} & 10 & 5 & 2 & 0 & -1 & -3 \end{bmatrix}$

CPT Value of Lottery L

$$V(L) = V^{gain}(L) + V^{loss}(L)$$

$$V^{gain}(L) = v(10)[w^{+}(0.1)] + v(5)[w^{+}(0.1 + 0.2) - w^{+}(0.1)]$$
$$+ v(2)[w^{+}(0.1 + 0.2 + 0.15) - w^{+}(0.1 + 0.2)]$$

$$V^{loss}(L) = v(-3)[w^{-}(0.2)] + v(-1)[w^{-}(0.2 + 0.25) - w^{-}(0.2)]$$

Allais Paradox Resolved

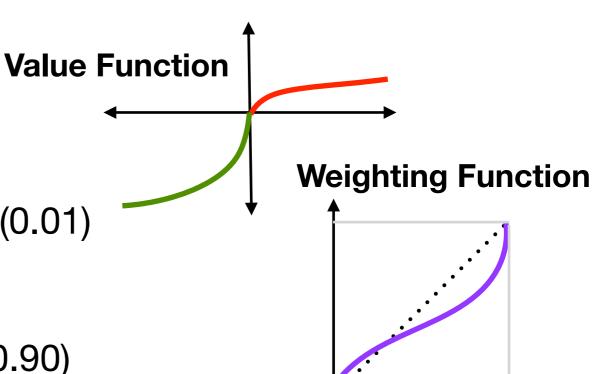
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Experiment 2				
Gamble 2A Gamble 2B				
Winnings Chance		Winnings	Chance	
Nothing	89%	Nothing	90%	
\$1 million	11%			
		\$5 million	10%	

Suppose r = \$1 million

V(Lottery 1A) = 0 $V(\text{Lottery 1B}) = v(4)w^{+}(0.10) + v(-1)w^{-}(0.01)$

 $V(\text{Lottery 2A}) = v(-1)w^{-}(0.89)$ $V(\text{Lottery 2B}) = v(4)w^{+}(0.1) + v(-1)w^{-}(0.90)$



Why CPT?

- Accommodates several empirically observed behavioral features
- Mathematically tractable
- Generalization of EUT

"... there is no good scientific reason why it (prospect theory) should not replace expected utility in current research, and be given prominent space in economics textbooks."

Outline

- Cumulative Prospect Theory (CPT)
- CPT Equilibrium Concepts Nash and Correlated equilibrium
- Results on the Geometry of CPT Equilibrium Notions
- Learning in CPT Games

Game Setup

$$\Gamma = (N, (A_i)_{i \in N}, (x_i)_{i \in N})$$

$$N = \{1, 2, \dots, n\}$$

Set of players

$$a_i \in A_i$$

Actions for player i

$$x_i:\prod_j A_j o \mathbb{R}$$

 $x_i: \prod A_j o \mathbb{R}$ Payoff function for player i

$$a = (a_1, \ldots, a_n)$$

Action profile

$$a \in A = \prod_{i} A_i$$

Set of Action profile

$$a_{-i} \in A_{-i} = \prod_{j \neq i} A_j$$

Set of Action profile of opponents

Game Setup (EUT)

$$N = \{1, 2, \dots, n\}$$
 Set of players
$$a_i \in A_i \qquad \text{Actions for player i}$$

$$x_i:\prod_j A_j o \mathbb{R}$$
 Payoff function for player i

For each player $i \in N$

Utility function
$$u_i(\cdot)$$

$$\downarrow$$

$$\downarrow$$
 Expected Utility $U_i(L)$

Game Setup (CPT)

$$N = \{1, 2, \dots, n\}$$

$$a_i \in A_i$$

$$x_i: \prod A_j \to \mathbb{R}$$

Set of players

Actions for player i

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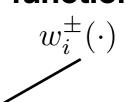
Reference point

 r_i

Value function $v_i(\cdot)$



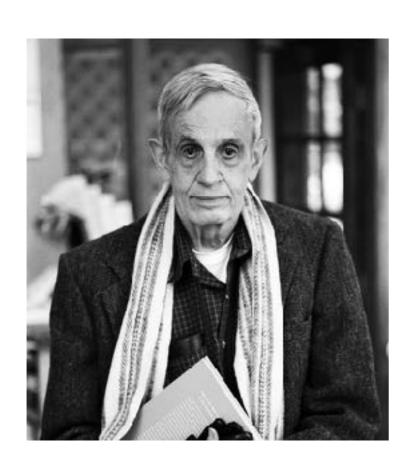
Probability weighting function



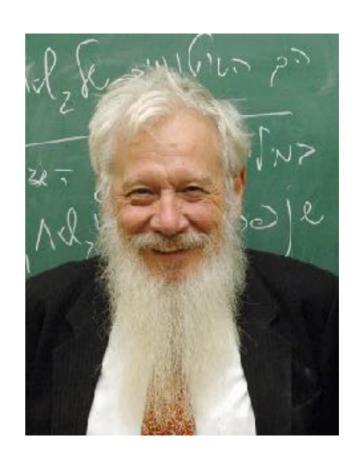
CPT value function

$$V_i(L)$$

Strategic Behavior in Games



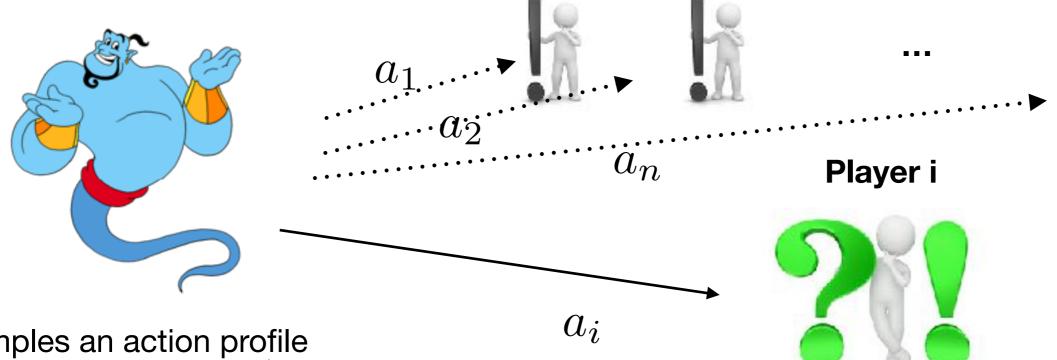
John Nash (1928–2015) Nash Equilibrium



Robert Aumann (b. 1930) Correlated Equilibrium

Correlated Equilibrium (CE)

Mediator

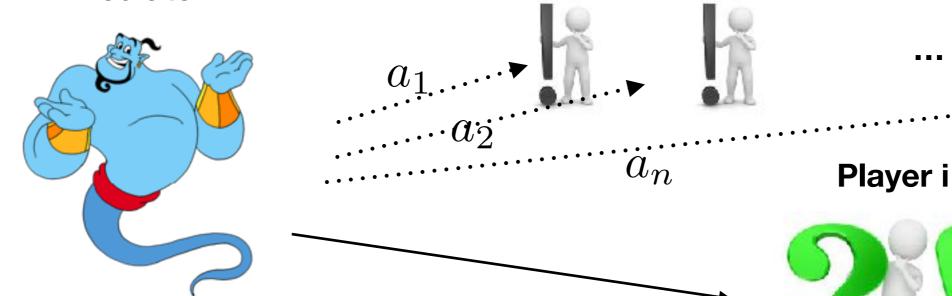


Samples an action profile $(a_1,\ldots,a_i,\ldots,a_n)$ from a distribution $\mu\in\Delta(A)$

Correlated Equilibrium (CE)

 a_i

Mediator



Samples an action profile $(a_1,\ldots,a_i,\ldots,a_n)$ from a distribution $\mu\in\Delta(A)$

Faces a lottery corresponding to each action

$$L_{i}(\mu, a_{i}, \tilde{a}_{i}) := \left\{ \underbrace{\left(\nu(a_{-i}), x_{i}(\tilde{a}_{i}, a_{-i})\right)}_{\mu(a_{i}, a_{-i})} \right\}_{a_{-i} \in A_{-i}} \nu(a_{-i}) = \underbrace{\frac{\mu(a_{i}, a_{-i})}{\sum_{a_{-i} \in A_{-i}} \mu(a_{i}, a_{-i})}}_{\mu(a_{i}, a_{-i})}$$

EUT Correlated Equilibrium (EUT CE)

Definition (Aumann 1987)

A distribution $\mu \in \Delta(A)$ is an **(EUT) Correlated Equilibrium** if no player with EUT preferences has an incentive to deviate from his signaled action, i.e.

$$U_i(L_i(\mu, a_i, a_i)) \ge U_i(L_i(\mu, a_i, \tilde{a}_i))$$

for all i, a_i, \tilde{a}_i such that $\mu_i(a_i) > 0$

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Incentive Constraints

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Incentive Constraints

Denote the set of all correlated equilibria by $C_{EUT}(\Gamma)$

CPT Correlated Equilibrium (CPT CE)

Definition (Keskin 2017)

A distribution $\mu \in \Delta(A)$ is a **CPT Correlated Equilibrium** if no player with CPT preferences has an incentive to deviate from his signaled action, i.e.

$$V_i(L_i(\mu, a_i, a_i)) \ge V_i(L_i(\mu, a_i, \tilde{a}_i))$$

for all i, a_i, \tilde{a}_i such that $\mu_i(a_i) > 0$

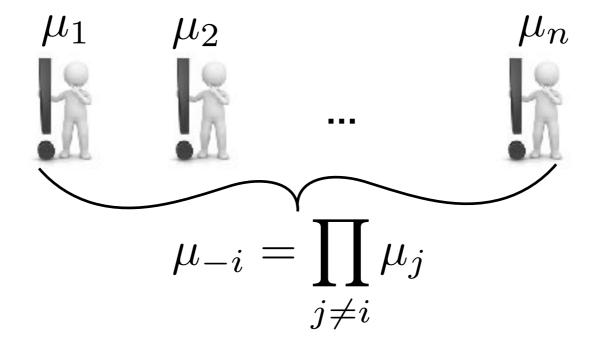
Incentive Constraints

Denote the set of all CPT correlated equilibria by $\ C(\Gamma)$

Nash Equilibrium (NE)

$$\mu \in \Delta^*(A) = \{ \mu \in \Delta(A) : \mu(a) = \mu_1(a_1) \times \cdots \times \mu_n(a_n) \quad \forall \ a \in A \}$$

Product form



Player i



$$L_i(\mu_{-i}, a_i) := \{(\mu_{-i}(a_{-i}), x_i(a_i, a_{-i}))\}_{a_{-i} \in A_{-i}}$$

EUT Nash Equilibrium (EUT NE)

Best response of player i to a product distribution $\mu \in \Delta^*(A)$

$$BR_i(\mu) := \left\{ \mu_i^* \in \Delta(A_i) | \operatorname{supp}(\mu_i^*) \subset \arg \max_{a_i \in A_i} U_i(L_i(\mu_{-i}, a_i)) \right\}$$

 μ_i^* Assigns positive probability only to optimal actions

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Definition (Nash 1951)

A product distribution $\mu \in \Delta^*(A)$ is EUT Nash equilibrium if

$$\mu^* \in BR_i(\mu^*)$$
 for all i

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Existence guaranteed by Kakutani fixed point theorem

CPT Nash Equilibrium (CPT NE)

Best response of player *i* to a product distribution $\mu \in \Delta^*(A)$

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A product distribution $\mu \in \Delta^*(A)$ is CPT Nash equilibrium if

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CPT Nash Equilibrium (CPT NE)

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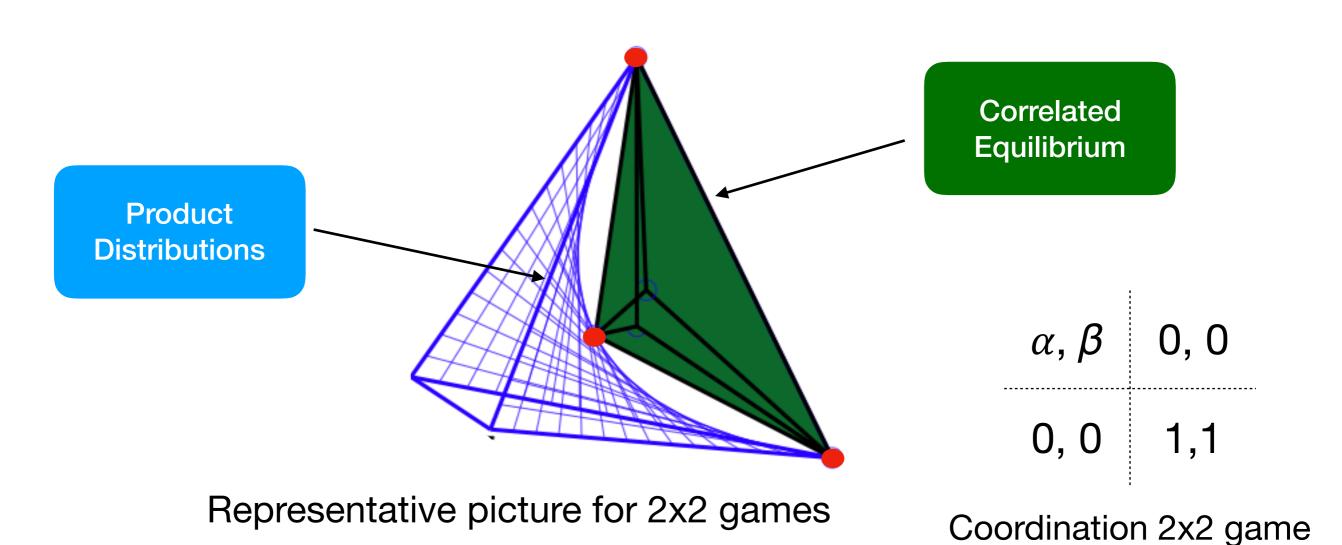
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Outline

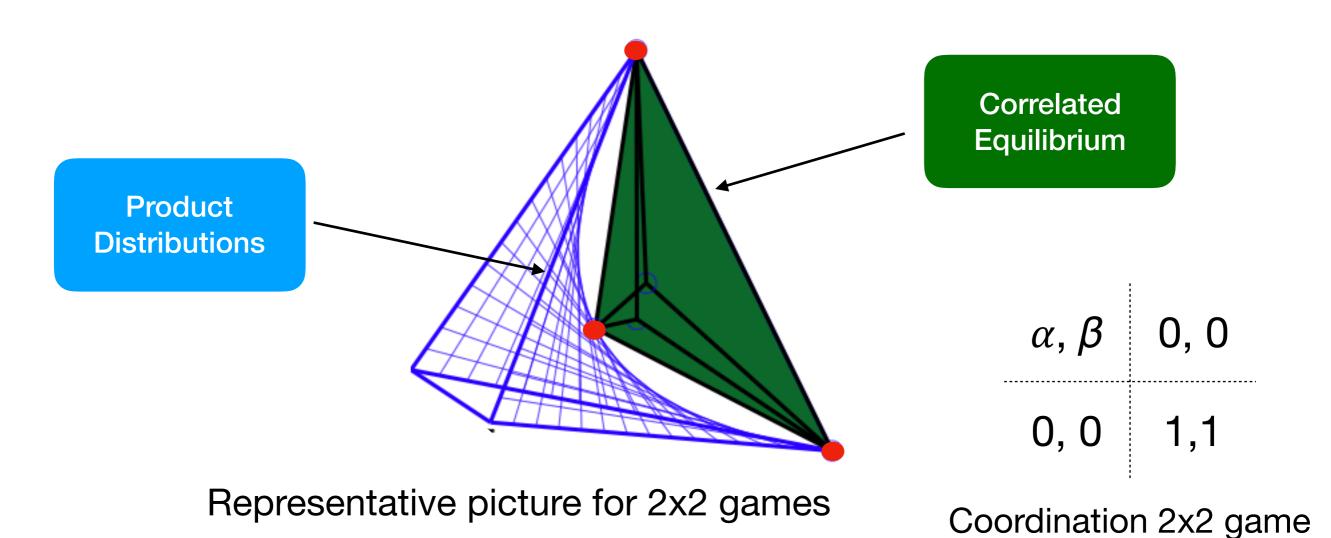
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Geometric Properties of Equilibria



Geometric Properties of Equilibria

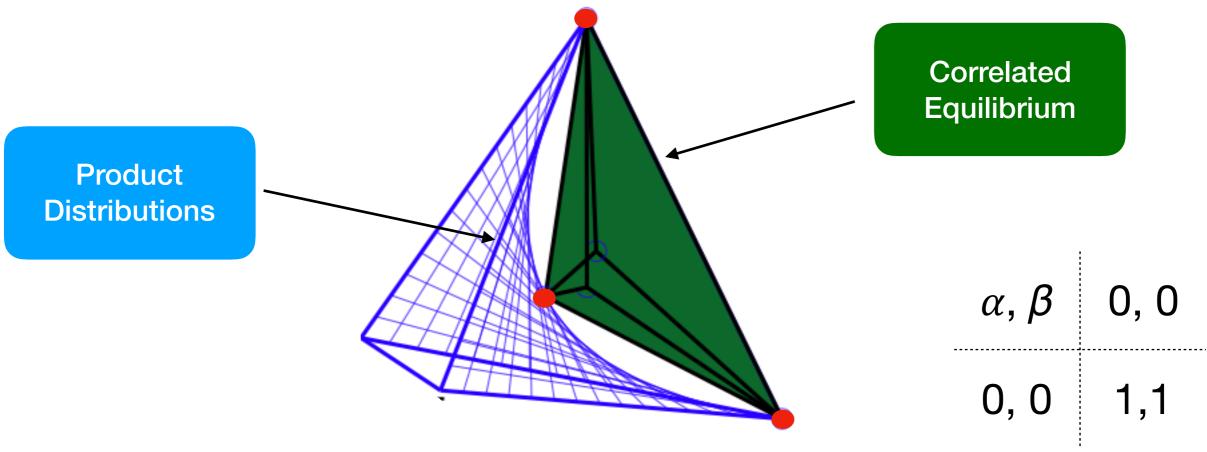
$$NE_{EUT}(\Gamma) = \Delta^*(A) \cap C_{EUT}(\Gamma)$$



Geometric Properties of Equilibria

$$NE_{EUT}(\Gamma) = \Delta^*(A) \cap C_{EUT}(\Gamma)$$

$$NE_{CPT}(\Gamma) = \Delta^*(A) \cap C(\Gamma)$$



Representative picture for 2x2 games

Coordination 2x2 game

Under EUT, Linearity in Probability helps!

Expanding Incentive Constraints using $U_i(L) = \sum u_i(x)p_i$

$$\sum_{a_{-i} \in A_{-i}} \mu(a) \left(u_i \left(x_i(a_i, a_{-i}) \right) - u_i \left(x_i(\tilde{a}_i, a_{-i}) \right) \right) \ge 0$$

for all $i, a_i, \tilde{a}_i \in A_i$

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for all
$$i, a_i, \tilde{a}_i \in A_i$$

The set $C_{EUT}(\Gamma)$ is a convex polytope

Under EUT, Linearity in Probability helps!

Expanding Incentive Constraints using $U_i(L) = \sum u_i(x)p_i$

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The set $C_{EUT}(\Gamma)$ is a convex polytope

What happens under CPT?

CPT Example with Non-convex C(Γ)

Player 2

Player 1

	I	II	III	IV	
0	2β,1	β+1,1	0,1	1,1	
1	1.99,0	1.99,0	1.99,0	1.99,0	

$$\beta = 1/w_1^+(0.5) = 2.299$$

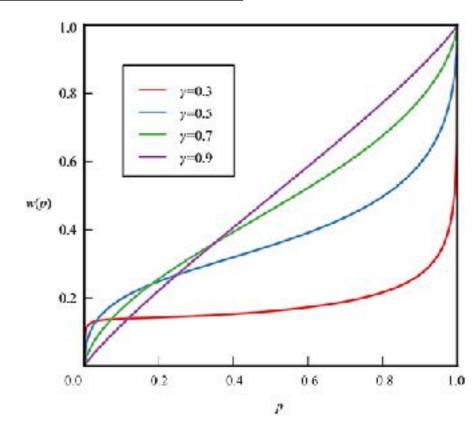
$$r_1 = r_2 = 0$$

$$\gamma_1 = 0.5, \gamma_2 = 1$$

$$v_1(z) = v_2(z) = z$$

Prelec 1998

$$w_i^+(p) = \exp\{-(-\ln p)^{\gamma_i}\}$$



CPT Example with Non-convex C(Γ)

Player 2

Player 1

	I	II	III	IV	
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	I	II	III	IV	
0	2β,1	β+1,1	0,1	1,1	
1	1.99,0	1.99,0	1.99,0	1.99,0	

Action 0 is player 1's Best Response to $\begin{cases} \mu_{odd} = (0.5, 0, 0.5, 0) \\ \mu_{over} = (0.0, 5, 0, 0.5) \end{cases}$

$$\begin{cases} \mu_{odd} = (0.5, 0, 0.5, 0) \\ \mu_{even} = (0, 0.5, 0, 0.5) \end{cases}$$

CPT Example with Non-convex C(Γ)

Player 2

Player 1

	I	II	III	IV	
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Action 1 is player 1's Best Response to $\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$

$$\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$$

CPT Example with Non-convex C(Γ)

Player 2

Player 1

	I	II	III	IV	
0	2β,1	β+1,1	0,1	1,1	
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Action 1 is player 1's Best Response to $\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$

$$\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$$

The set $C(\Gamma)$ is Non-convex!

CPT Example

Player 2

Player 1

	I	II	III	IV	
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1	1.99,0	1.99,0	1.99,0	1.99,0	

$$eta=1/w_1^+(0.5)=2.299$$
 $r_1=r_2=0$ $\gamma_1=0.5, \gamma_2=1$ $v_1(z)=v_2(z)=z$

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Action 1 is player 1's Best Response to $\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$

$$\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$$

$$V_1(L_1(\mu_{odd}, 0)) = 2\beta w_1^+(0.5) = 2,$$
 $V_1(L_1(\mu_{odd}, 1)) = 1.99$
 $V_1(L_1(\mu_{even}, 0)) = 1 + \beta w_1^+(0.5) = 2,$ $V_1(L_1(\mu_{even}, 1)) = 1.99$

Our results on the structure of CE (CPT)

Result (P., Anantharam 2017)

For any 2x2 game, the set $C(\Gamma)$ is a **convex polytope**.

Our results on the structure of CE (CPT)

Result (P., Anantharam 2017)

For any 2x2 game, the set $C(\Gamma)$ is a **convex polytope**.

Result (P., Anantharam 2017)

We provide an example of a 3x3 game for which the set $C(\Gamma)$ is **disconnected**.

$$NE_{EUT}(\Gamma) = \Delta^*(A) \cap C_{EUT}(\Gamma)$$

Theorem (Nau et al 2003)

The Nash equilibria all lie on the boundary of the correlated equilibria polytope.

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Comments:

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Comments:

1. Boundary of the correlated equilibrium set when it is viewed as a subset of $\Delta(A)$.

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Comments:

- 1. Boundary of the correlated equilibrium set when it is viewed as a subset of $\Delta(A)$.
- 2. If the set of correlated equilibria is not "full dimensional" then the statement is trivial.

$$NE_{EUT}(\Gamma) = \Delta^*(A) \cap C_{EUT}(\Gamma)$$

Theorem (Nau et al 2003)

The Nash equilibria all lie on the boundary of the correlated equilibria polytope.

Comments:

- 1. Boundary of the correlated equilibrium set when it is viewed as a subset of $\Delta(A)$.
- 2. If the set of correlated equilibria is not "full dimensional" then the statement is trivial.
- 3. The statement cannot be strengthened because in the case of less than full dimensional C(Γ), Nash equilibria can lie in the relative interior of this convex polytope.

Theorem (Nau et al 2003)

The Nash equilibria all lie on the boundary of the correlated equilibria polytope.

Proof sketch:

A Nash equilibria renders every player indifferent among all of her own strategies, hence it satisfies all of the incentive constraints with **equality**, at least one of which is non trivial if the game is non trivial, and hence lies on one of the faces of the convex polytope CE_{EUT} .

$$\sum_{s_{-i} \in S_{-i}} \mu(s) (u_i(h_i(s_i, s_{-i})) - u_i(h_i(d_i, s_{-i}))) \ge 0,$$

for all i and for all $s_i, d_i \in S_i$

$$NE_{CPT}(\Gamma) = \Delta^*(A) \cap C(\Gamma)$$

Theorem (P., Anantharam 2017)

The CPT Nash equilibria all lie on the boundary of the CPT correlated equilibria set.

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Theorem (P., Anantharam 2017)

The CPT Nash equilibria all lie on the boundary of the CPT correlated equilibria set.

Proof sketch

1. Only need to focus on completely mixed CPT NE.

$$NE_{CPT}(\Gamma) = \Delta^*(A) \cap C(\Gamma)$$

Theorem (P., Anantharam 2017)

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- 3. There exists a direction in the probability space along which at least one incentive constraint is violated.
- 4. Thus CPT Nash equilibria cannot have a ball around it that is completely contained inside the set of CPT correlated equilibria.

A Useful Lemma

Two distinct lotteries (let $p_i > 0$ for j=1,...,t)

$L_1 =$	probability	p_1	p_2	-	-	-	p_t
	outcome	x_1	x_2	-	-		x_t
$L_2 =$	probability	p_1	p_2				$\mid p_t \mid$
L_2 —	outcome	y_1	y_2				y_t

that satisfy either of the two properties

- 1. they are not similarly ranked or
- 2. neither of them dominates the other

then $\boldsymbol{\exists}$ a direction $\boldsymbol{\delta}=(\delta_1,\;\dots\;,\;\delta_n),\;\boldsymbol{\Sigma}\;\delta_i=0$ such that

$$V(p + \epsilon \delta, x) - V(p + \epsilon \delta, y) < V(p, x) - V(p, y)$$

Outline

- Cumulative Prospect Theory (CPT)
- CPT Equilibrium Concepts Nash and Correlated equilibrium
- Results on the Geometry of CPT Equilibrium Notions
- Learning in CPT Games

- Neoclassical economics:
 - hyper-rational players,
 - completely understand the structure of the game,
 - have a coherent model of others' behavior,
 - make rational calculations of infinite complexity,
 - and all of this is common knowledge

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- Learning in Games:
 - boundedly-rational players,
 - players make decisions on limited data,
 - and use simple predictive models

 Classical solution concepts (like NE and CE) can be "recovered" via this route.

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Examples:

- Fictitious play: Nash Equilibrium in zero sum games, potential games, 2x2 games
- Calibrated learning, Adaptive heuristics, Follow the perturbed leader: Correlated Equilibrium

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Examples:

- Fictitious play: Nash Equilibrium in zero sum games, potential games, 2x2 games
- Calibrated learning, Adaptive heuristics, Follow the perturbed leader: Correlated Equilibrium
- Question: What if the players behave according to CPT?

Repeated games

- A Game Γ is played repeatedly at each step t = 1,2,...
- Player i's action sequence: a_i^1, a_i^2, \dots
- Action profile at step t: $a^t = (a_1^t, \dots, a_n^t)$
- History at step t: $H^t = (a^1, a^2, \dots, a^{t-1})$
- Randomized strategy sequence for player i at step t:

$$\sigma_i^t: H^t \to \Delta(A_i)$$

• Empirical distribution at step t: $\boldsymbol{\xi}^t$

$$\xi^{t}(a) = \frac{\text{\#action profile } a \text{ appears in } H^{t}}{t-1}$$

Foster Vohra result







Vohra

- At every step t, each player i predicts a distribution $\mu_{-i}^t \in \Delta(A_{-i})$ on the action profile of the other players.
- Based on this prediction she plays a best EUT response $\,a_i^t\,$

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Theorem (Foster, Vohra 1997)

If each players' predictions are calibrated, then the Empirical Distribution of play convergence to the set of correlated equilibria.

$$\lim_{t \to \infty} d(\xi^t, C_{EUT}(\Gamma)) = 0$$

Nature vs Forecaster

Step	1	2	3	4	5	6	7	8	
Nature	0	1	1	0	0	0	1	0	
Forecaster	10%	70%	80%	30%	10%	10%	80%	30%	

Nature vs Forecaster

Step	1	2	3	4	5	6	7	8	
Nature	0	1	1	0	0	0	1	0	
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Nature: $y^1, y^2, \ldots \in S$.

Forecaster: $q^1, q^2, \ldots \in \Delta(S)$.

Nature vs Forecaster

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 $\rho(q,y,t) = \frac{\# \text{Forecaster predicts } q \text{ and Nature plays } y \text{ up to step } t}{N(q,t)}$

Nature vs Forecaster

Step	1	2	3	4	5	6	7	8	
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$$\lim_{t\to\infty}\sum_{q\in Q^t}|\rho(q,y,t)-q(y)|\frac{N(q,t)}{t}=0 \text{ for all } y\in S,$$

Example of CPT calibrated learning

Player 2/Player 3

	l	II	III	IV
0	2β,1	β+1,1	0,1	1,1
1	1.99,0	1.99,0	1.99,0	1.99,0

- If Player 2 Action ≠ Player 3 Action then all players receive a payoff of -1.
- If Player 2 action = player 3 action then Player 1 receives first payoff shown in table and players 2 and 3 each receive the second payoff.

Non convergence of Calibrated learning to CPT correlated equilibrium

Action 0 is player 1's Best Response to $\begin{cases} \mu_{odd} = (0.5, 0, 0.5, 0) \\ \mu_{even} = (0, 0.5, 0, 0.5) \end{cases}$

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Action 1 is player 1's Best Response to $\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$

$$\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$$

Step	1	2	3	4	5	6
Player 1 Action	0	0	0	0	0	0
Player 2 Action	I	II	III	IV	I	II
Player 3 Action	I	II	III	IV	I	II
Player 1 Forecast	μ _{odd}	Џеvеn	µ odd	μ _{even}	µ odd	µeven

	I	Ш	III	IV
0	0.25	0.25	0.25	0.25
1	0	0	0	0

 $C(\Gamma,i,a_i)\subset \Delta(A_{-i})$: distributions for which action a_i is player i's Best Response

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$$C(\Gamma, i) := \{ \mu \in \Delta(A) | \mu(\cdot | a_i) \in C(\Gamma, i, a_i), \forall a_i \in \text{supp}(\mu_i) \}$$

Distributions for which player i has no incentive to deviate

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$$C(\Gamma) = \bigcap_{i \in N} C(\Gamma, i)$$

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Distributions for which player i has no incentive to deviate

$$C(\Gamma) = \bigcap_{i \in N} C(\Gamma, i)$$

$$D(\Gamma) = \bigcap_{i \in N} co(C(\Gamma, i))$$

$$N = \{1, 2, \dots, n\}$$

$$a_i \in A_i$$

$$x_i:\prod_i A_j \to \mathbb{R}$$

Set of players

Actions for player i

 $x_i: \prod A_j o \mathbb{R}$ Payoff function for player i

$$N = \{1, 2, \dots, n\}$$
$$a_i \in A_i$$

$$x_i: \prod_j A_j \to \mathbb{R}$$

$$b_i \in B_i$$

$$b=(b_1,\ldots,b_n)$$

$$B = \prod_{i} B_{i}$$

$$\sigma_i: B_i \to \Delta(A_i)$$

Set of players

Actions for player i

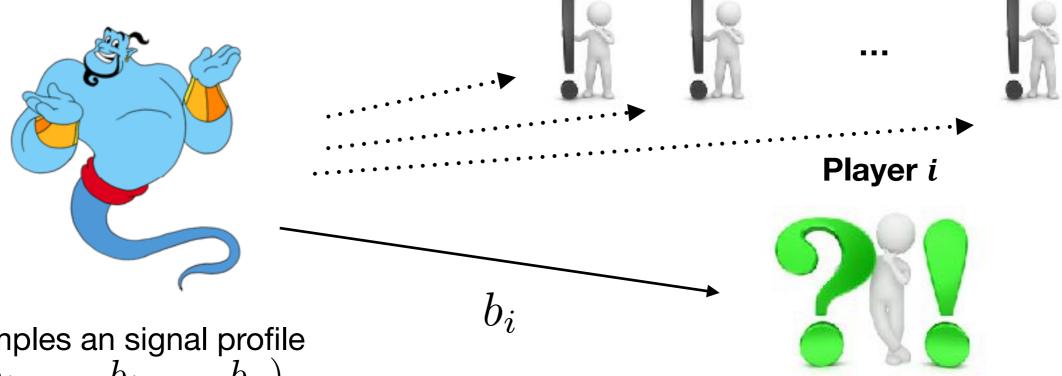
 $x_i: \prod A_j \to \mathbb{R}$ Payoff function for player i

Signal set for player i Signal profile

Set of Signal profiles

Strategy of player i

Mediator



Samples an signal profile $(b_1,\ldots,b_i,\ldots,b_n)$ from a distribution

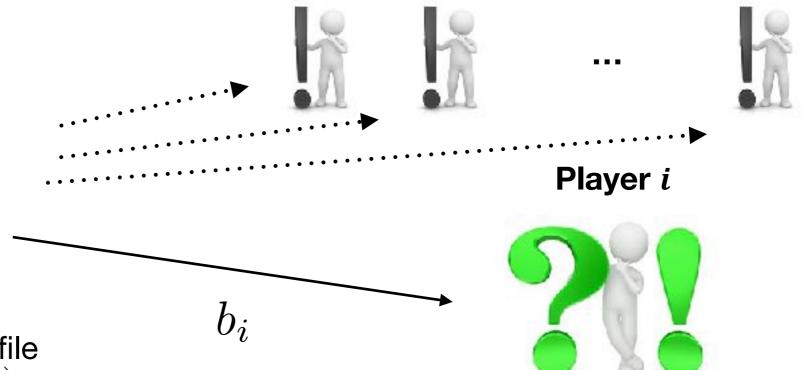
$$\psi \in \Delta(B)$$

Mediator



Samples an signal profile $(b_1,\ldots,b_i,\ldots,b_n)$ from a distribution

$$\psi \in \Delta(B)$$



Faces a lottery corresponding to each action

$$L_i(\psi, b_i, a_i) := \left\{ \underbrace{\tilde{\mu}_{-i}(a_{-i}|b_i)}_{\downarrow}, x_i(a_i, a_{-i}) \right\}_{a_{-i} \in A_{-i}}$$

$$\tilde{\mu}_{-i}(a_{-i}|b_i) := \sum_{b_{-i} \in B_{-i}} \psi_{-i}(b_{-i}|b_i) \prod_{j \in N \setminus i} \sigma_j(b_j)(a_j)$$

Mediated CPT NE

Best response of player i

$$BR_{i}(\psi, \sigma) := \left\{ \sigma_{i}^{*} : B^{i} \to \Delta(A_{i}) \middle| \text{ for all } b_{i} \in B_{i} \right.$$

$$\sup(\sigma_{i}^{*}(b_{i})) \subset \arg\max_{a_{i} \in A_{i}} V_{i}\left(L_{i}(\psi, b_{i}, a_{i})\right) \right\}$$

 $\sigma_i^*(b_i)$ Assigns positive probability only to optimal actions

Mediated CPT NE

Best response of player i

$$BR_i(\psi, \sigma) := \left\{ \sigma_i^* : B^i \to \Delta(A_i) \middle| \text{ for all } b_i \in B_i \right\}$$

$$\operatorname{supp}(\sigma_i^*(b_i)) \subset \operatorname{arg} \max_{a_i \in A_i} V_i \left(L_i(\psi, b_i, a_i) \right) \right\}$$

Definition (P., Anantharam 2018)

A randomized strategy profile σ is CPT Mediated Nash equilibrium if

$$\sigma_i \in BR_i(\psi, \sigma)$$
 for all i

Mediated CPT CE

Distribution induced by ψ and strategy profile σ on the set of action profiles

$$\eta(\psi,\sigma)(a) := \sum_{b \in B} \psi(b) \prod_{i \in N} \sigma_i(b_i)(a_i)$$

Mediated CPT CE

Distribution induced by ψ and strategy profile σ on the set of action profiles

$$\eta(\psi,\sigma)(a) := \sum_{b \in B} \psi(b) \prod_{i \in N} \sigma_i(b_i)(a_i)$$

Definition (P., Anantharam 2018)

A distribution $\mu \in \Delta(A)$ is a **mediated CPT correlated equilibrium** iff there exists a signal system B_i , a mediator distribution ψ and a mediated CPT Nash equilibrium σ with respect to them such that

$$\eta(\psi,\sigma) = \mu$$

Mediated CPT CE

Theorem (P., Anantharam 2018)

The set of all Mediated CPT correlated equilibria is

$$D(\Gamma) = \bigcap_{i \in N} co(C(\Gamma, i))$$

Corollary 1

For EUT games, $D(\Gamma) = C(\Gamma)$.

Corollary 2

For 2x2 CPT games, $D(\Gamma) = C(\Gamma)$.

Convergence of Calibrated learning

- At every step t, each player i predicts a distribution $\mu_{-i}^t \in \Delta(A_{-i})$ on the action profile of the other players.
- Based on this prediction she plays a best CPT response $\,a_i^t\,$

Theorem (P., Anantharam 2018)

If each players' predictions are calibrated, then the Empirical Distribution of play convergence to the set of mediated CPT correlated equilibria.

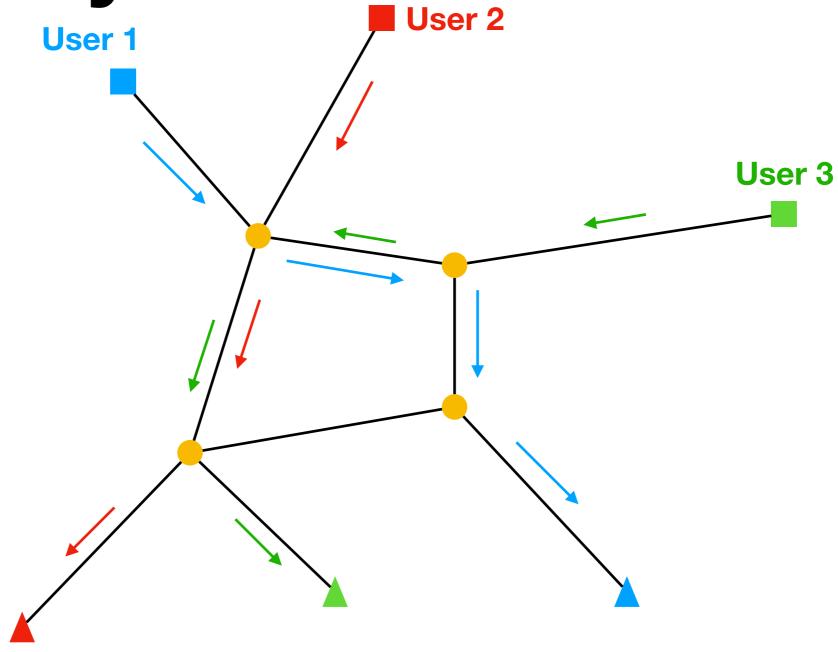
$$\lim_{t \to \infty} d(\xi^t, D(\Gamma)) = 0$$

Converse

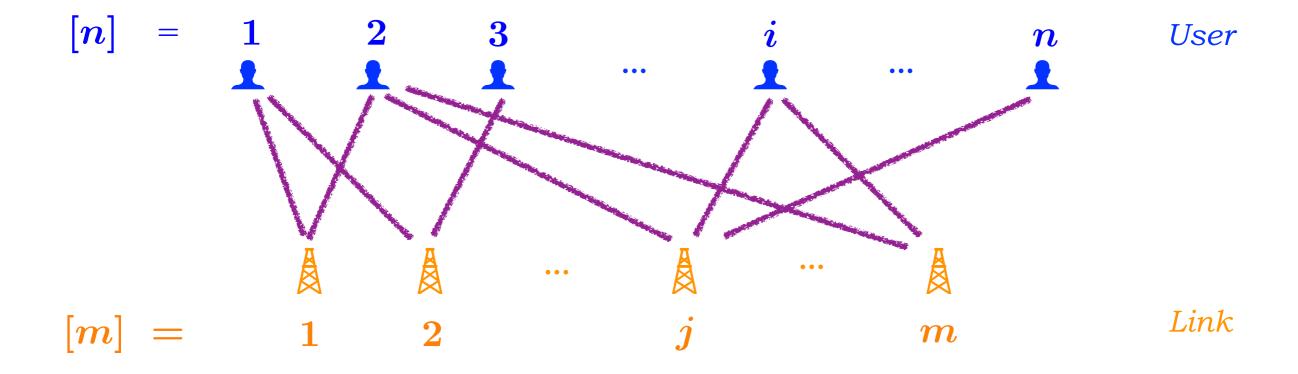
Theorem

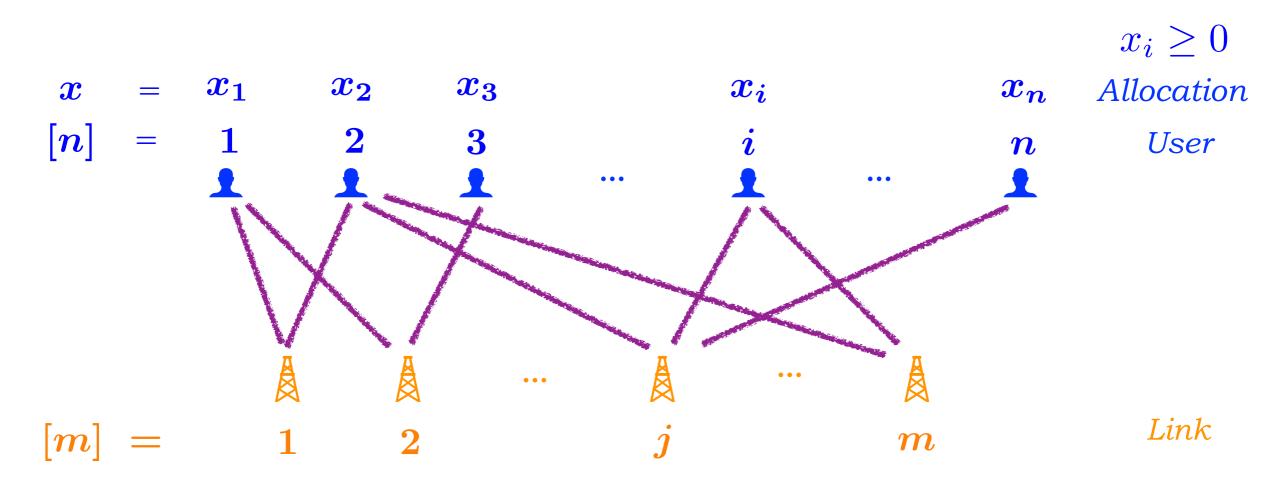
If the sets $C(\Gamma, i, a_i)$ do not have any isolated points, then for any $\mu \in D(\Gamma)$ there exists a sequence of play and corresponding assessments that are calibrated such that the Empirical Distribution converges to μ .

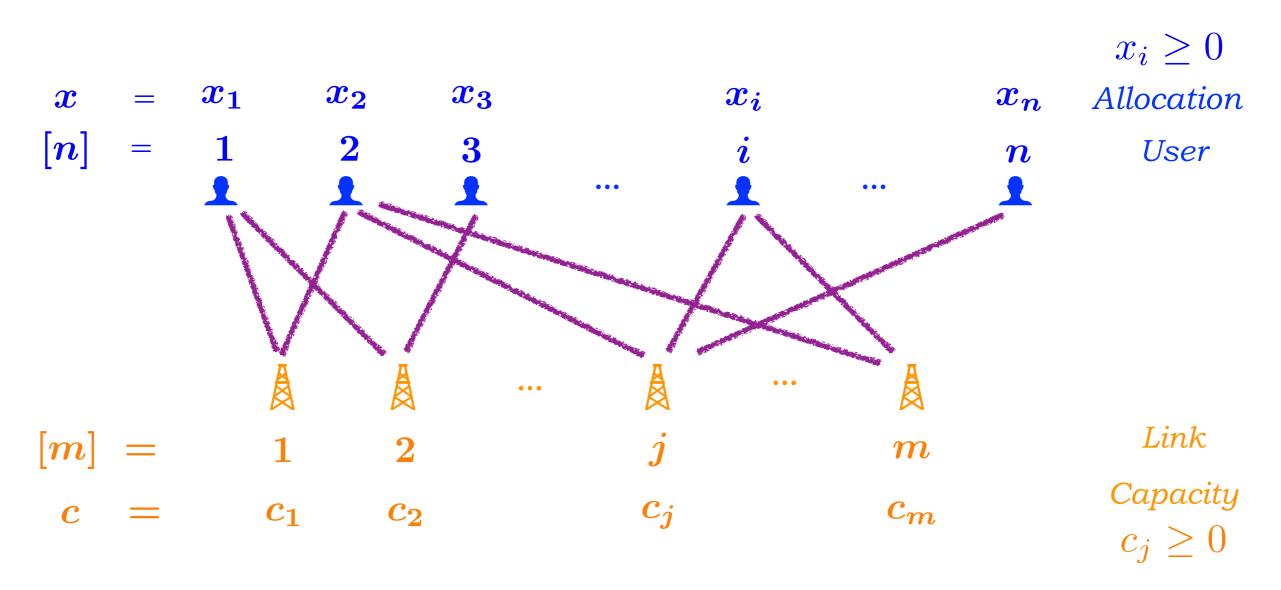
Such games are Generic

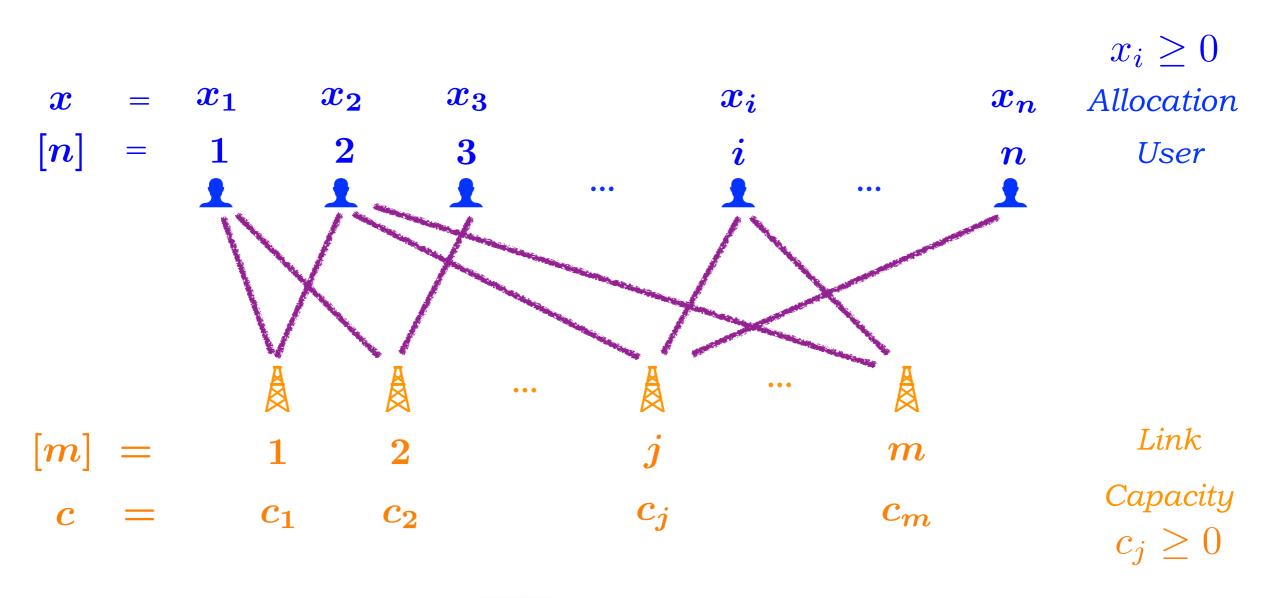


Kelly (1997) - Charging and rate control for elastic traffic









$$\sum_{i \in R_j} x_i \le c_j, \forall j$$

Link Constraints

System problem in the Kelly model

SYSTEM(U, A, C)

Maximize

$$\sum_{i=1}^{n} U_i(x_i)$$

subject to

$$\sum_{i \in R_j} x_i \le c_j, \ \forall j,$$

$$x_i \geq 0, \forall i.$$

User problem in the Kelly model

$$USER_i(U_i, \lambda_i)$$

Maximize

$$U_i(\frac{w_i}{\lambda_i}) - w_i$$

subject to

$$w_i \ge 0$$
.

 w_i : Amount per unit time that user i is willing to pay

 λ_i : charge per unit flow that the network presents to user i

Network problem in the Kelly model

Maximize

$$\sum_{i} w_i \log x_i$$

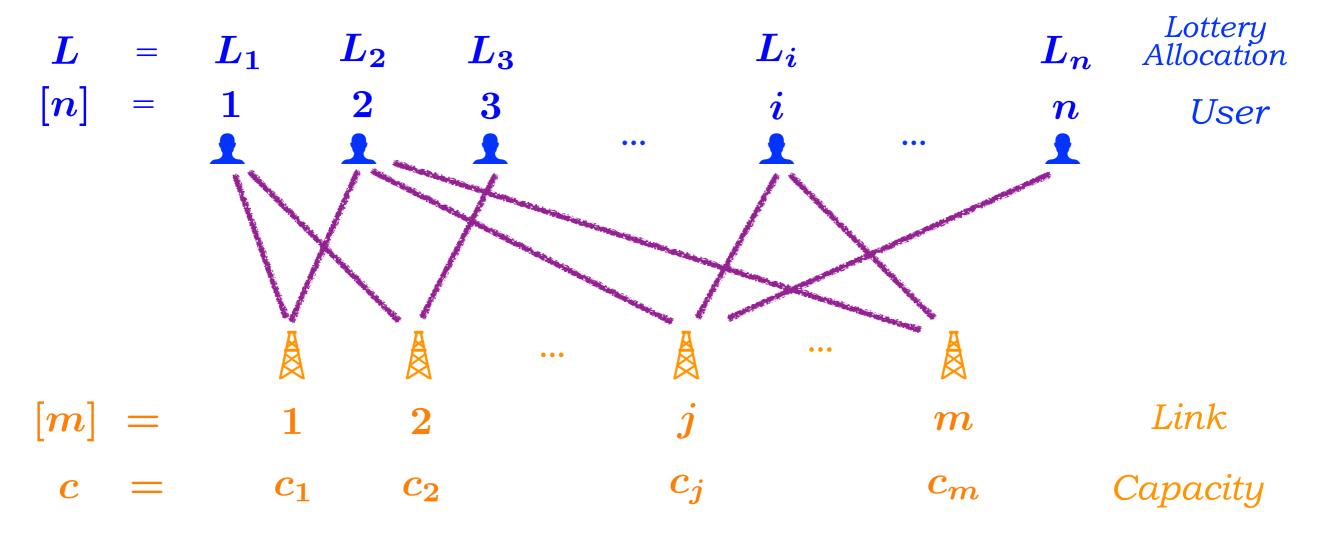
subject to

$$\sum_{i \in R_j} x_i \le c_j, \ \forall j,$$

$$x_i \geq 0, \forall i.$$

Since $w_i = \lambda_i x_i$ having found x_i the network can present λ_i to user i

Lottery Allocation



<i>T</i> —	probability	0.1	0.3	0.25	0.35
L_i —	allocation	10	5	3	6

Implementable Allocation Schemes

User i

\boldsymbol{l}	

$y_i(l)$	1	2	3
1	10	3	2
2	5	5	5
3	3	10	2
4	0	5	10

$$\sum_{i \in R_j} y_i(l) \le c_j, \forall j, l$$

Implementable Allocation Schemes

User i

$y_i(l)$	1	2	3
1	10	3	2
2	5	5	5
3	3	10	2
4	0	5	10

p(l)
0.25
0.25
0.25
0.25

Uniformly Distributed

$$\sum_{i \in R_j} y_i(l) \le c_j, \forall j, l$$

Alternative

permutations and decision weights

User i

 $y_i(l)$ 1
 2
 3

 1
 10
 3
 2

 2
 5
 5
 5

 3
 3
 10
 2

 4
 0
 5
 10

Alternative

User *i*

$z_i(l)$	1	2	3
1	10	10	10
2	5	5	5
3	3	5	2
4	0	3	2

User i

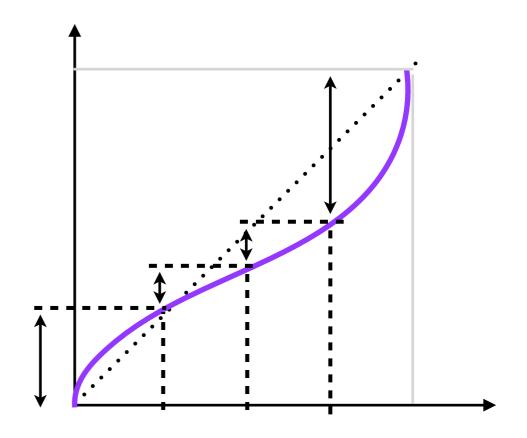
$\pi_i(l)$	1	2	3
1	1	4	3
2	2	2	2
3	3	1	4
4	4	3	1

$$y_i(l) = z_i(\pi_i(l))$$

$$V_{i}(\mathbf{L}_{i}) = \sum_{l \in [k]} h_{i}(l)v_{i}(z_{i}(l))$$

$$w_{i}(\frac{l}{k}) - w_{i}(\frac{l-1}{k})$$

Alternative *1*



Alternative 1

System Problem

$$SYS[z, \pi; h, v, A, c]$$

Maximize
$$\sum_{i}^{n} V_{i}(L_{i})$$

Subject to
$$\sum_{i \in R_j} y_i(l) \leq c_j, \forall j, l$$

System Problem

$$SYS[z, \pi; h, v, A, c]$$

Maximize
$$\sum_{i=1}^{n} \sum_{l=1}^{k} h_i(l) v_i(z_i(l))$$

Subject to
$$\sum_{i \in R_j} y_i(l) \leq c_j, \forall j, l$$

System Problem

$$SYS[z, \pi; h, v, A, c]$$

Maximize

$$\sum_{i=1}^{n} \sum_{l=1}^{k} h_i(l) v_i(z_i(l))$$

Subject to

$$\sum_{i \in R_j} z_i(\pi_i(l)) \le c_j, \forall j, l$$

$$z_i(l) \ge z_i(l+1), \forall i, l$$

$$\pi_i \in S_k, \forall i$$

Fixed permutation problem

Maximize
$$\sum_{i=1}^{n} \sum_{l=1}^{k} h_i(l) v_i(z_i(l))$$

Subject to

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$$\pi_i \in S_k, \forall i$$

Fixed permutation problem

$$SYS_FIX[z; \pi, h, v, A, c]$$

Maximize
$$\sum_{i=1}^{n} \sum_{l=1}^{k} h_i(l) v_i(z_i(l))$$

Subject to

$$\sum_{i \in R_j} z_i(\pi_i(l)) \le c_j, \forall j, l$$

$$z_i(l) \ge z_i(l+1), \forall i, l$$

Convex **Optimization**

User pricing through menu of rates

Rates

Budgets

Menu of Vector of Incremental **Allocation**

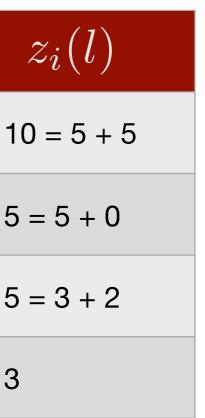
Lottery Allocation

l	
1	
2	
3	
4	

$r_i(l)$
1
2
3
5

$m_i(l)$
5
0
6
15

$\delta_i(l)$
5 = 5/1
0 = 0/2
2 = 6/3
3 = 15/5



User problem

 $\mathrm{USER}[m_i; r_i, h_i, v_i]$

Maximize
$$\sum_{l=1}^k h_i(l) v_i \left(\sum_{s=l}^k \frac{m_i(s)}{r_i(s)} \right) - \sum_{l=1}^k m_i(l)$$

Subject to

$$m_i(l) \geq 0, \forall l$$

Network Problem

$$NET[\delta; m, \pi, A, c]$$

Maximize

$$\sum_{i=1}^{n} \sum_{l=1}^{k} m_i(l) \log(\delta_i(l))$$

Subject to
$$\sum_{i \in R_j} \sum_{s=\pi_i(l)}^k \delta_i(s) \leq c_j, orall j, l$$

$$\delta_i(l) \geq 0, \forall i, l$$

- Eisenberg, Gale (1959) Consensus of subjective probabilities: the parimutuel method
- Kelly (1998) Rate control for communication networks: shadow prices, proportional fairness and stability
- Jain, Vazirani (2010) Eisenberg-Gale markets: Algorithms and Game Theoretic properties

Equilibrium

Theorem

For any fixed permutation, there exist equilibrium

rates budgets incremental allocation lottery allocations $r_i^*(l)$ $m_i^*(l)$ $\delta_i^*(l)$ $z_i^*(l)$

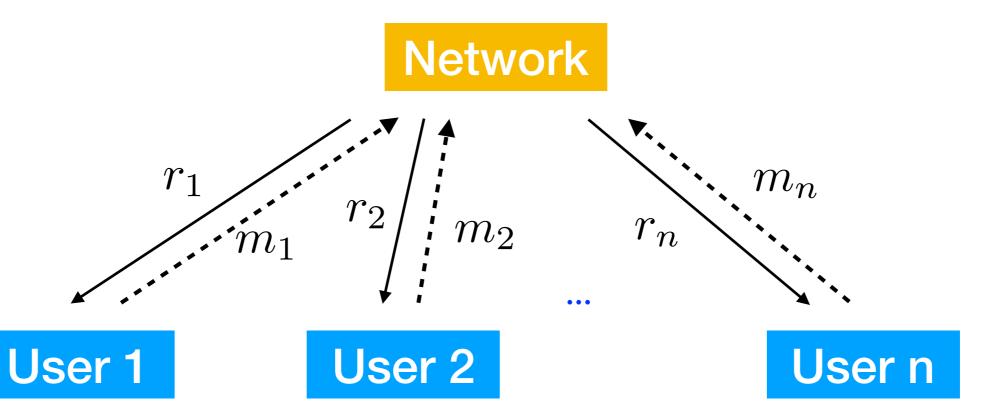
such that

- (i) $m_i^*(l)$ solves $\mathrm{USER}[m_i; r_i^*, h_i, v_i]$
- (ii) $\delta_i^*(l)$ solves $\operatorname{NET}[\delta;m^*,\pi,A,c]$
- (iii) $m_i^*(l) = \delta_i^*(l)r_i^*(l)$
- (iv) $\delta_i^*(l) = z_i^*(l) z_i^*(l+1)$
- (v) $z_i^*(l)$ solves the fixed permutation system problem

 $SYS_FIX[z; \pi, h, v, A, c]$

Iterative Process

 $NET[\delta; m, \pi, A, c]$



 $\mathrm{USER}[m_i; r_i, h_i, v_i]$

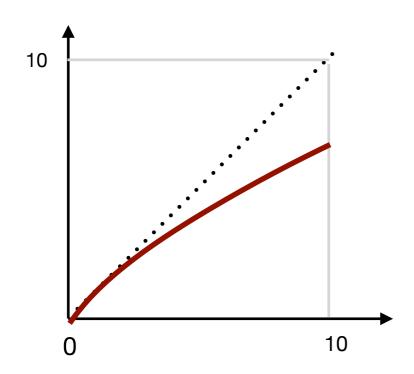
Example

n = 10 Players

$$m = 1 link$$

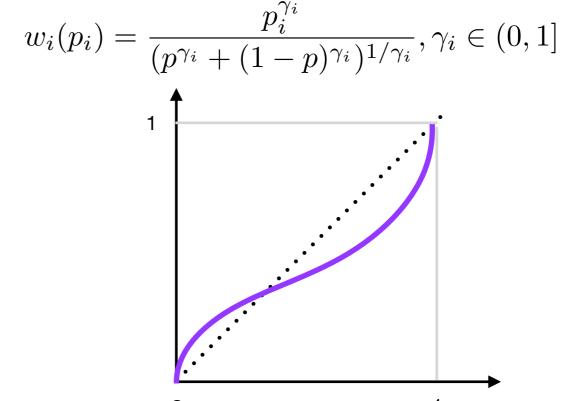
Link capacity c = 10

$$v_i(x_i) = x_i^{\beta_i}, \beta_i \in [0, 1]$$



Value function

$$\beta_i = 0.88$$



Probability weighting function

$$\gamma_i = 0.61$$

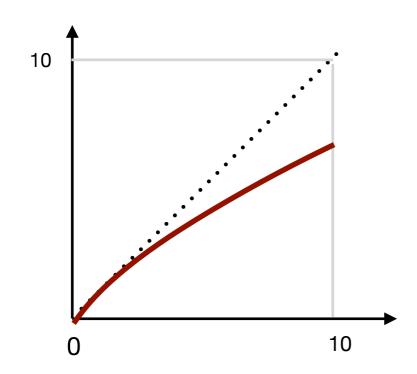
Example

n = 10 Players

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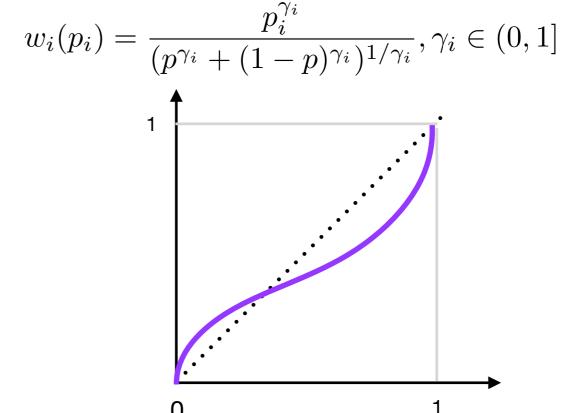
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Deterministic Allocation10

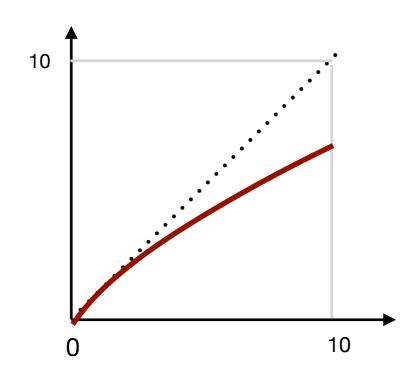
Example

n = 10 Players

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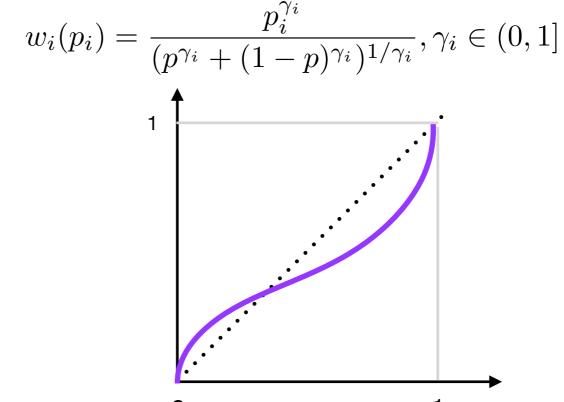
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Value function

$$\beta_i = 0.88$$

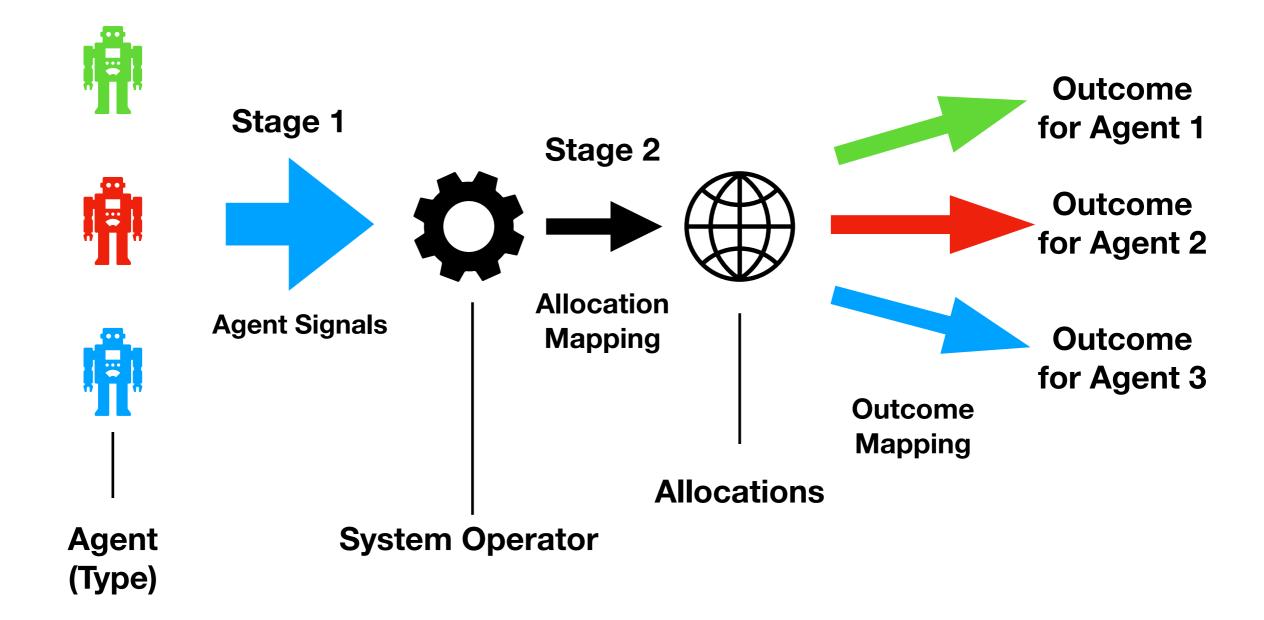


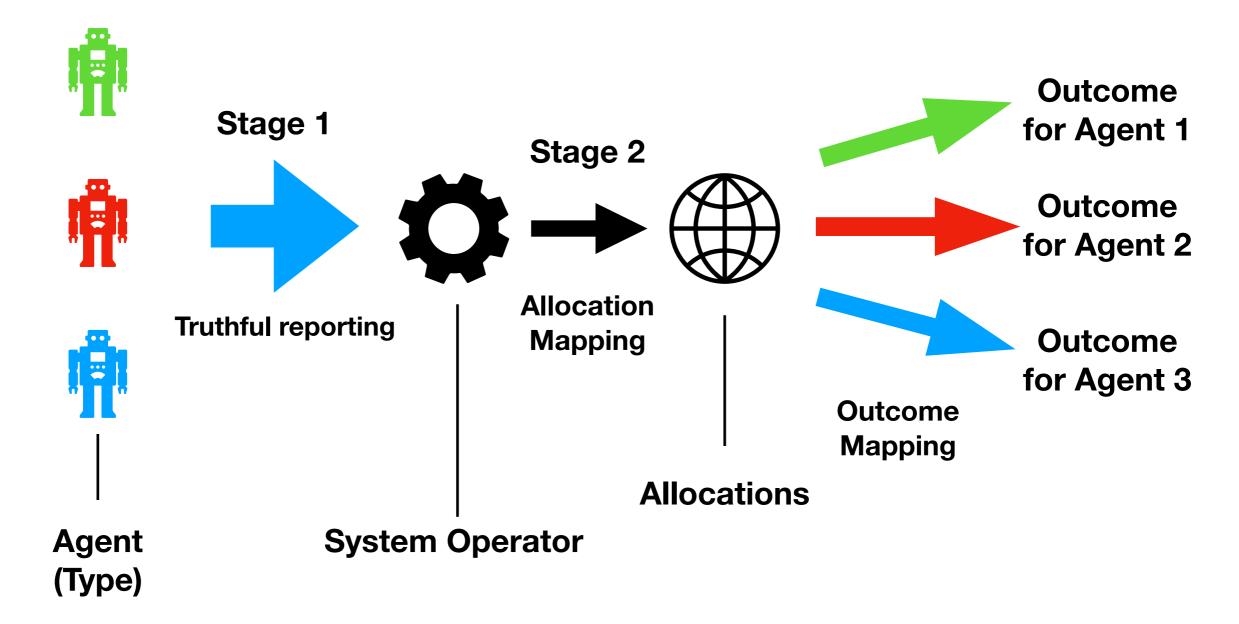
Probability weighting function

$$\gamma_i = 0.61$$

Deterministic Allocation10

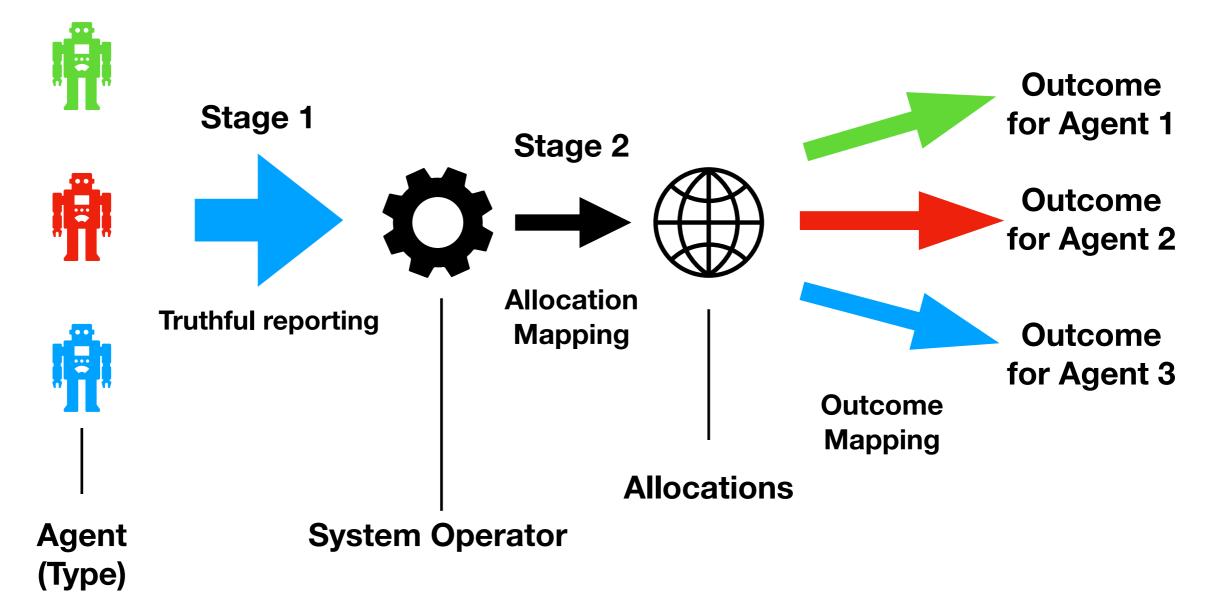
Lottery Allocation 14.17





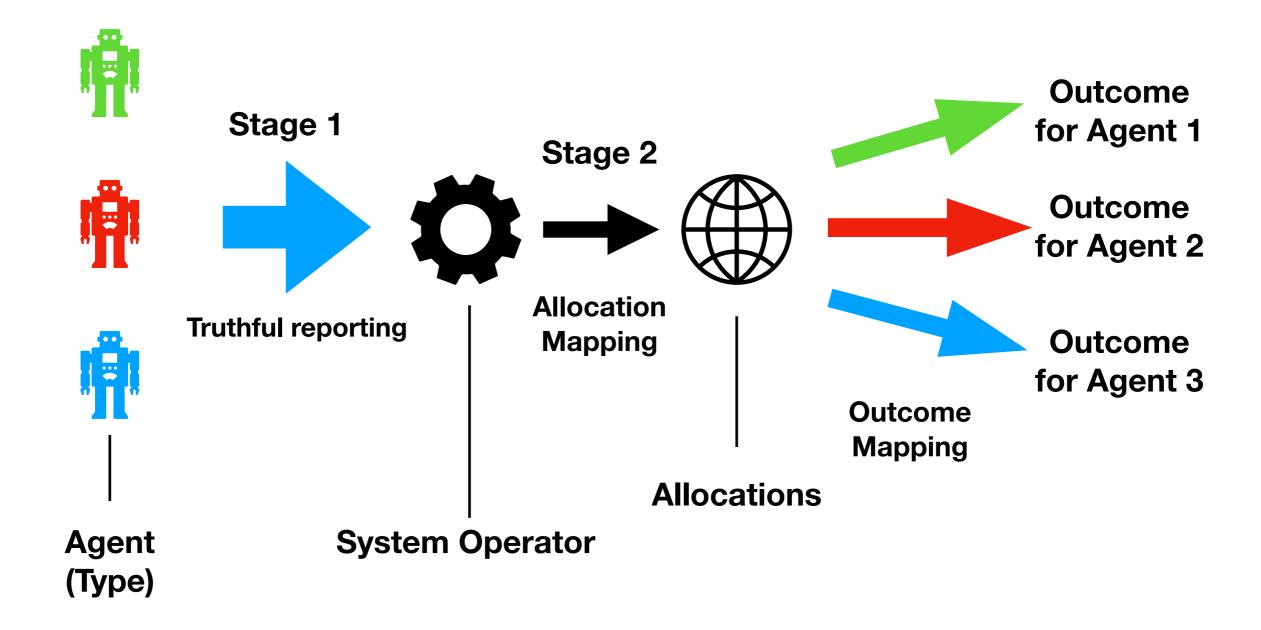
Revelation Principle (under EUT)

- WLOG assume signal set = type set for each player
- restrict attention to direct truthful mechanisms



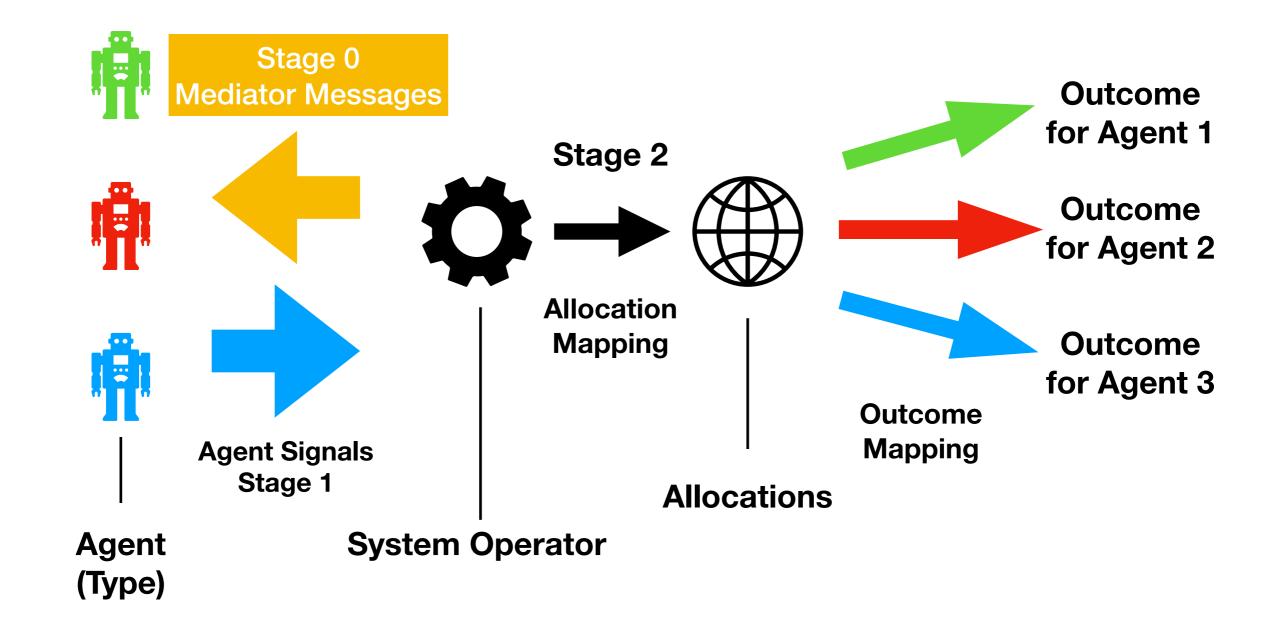
Importance of truthful strategies

- Limits on information availability
- Computational and cognitive limitations
- Users with different levels of access to information and computation.

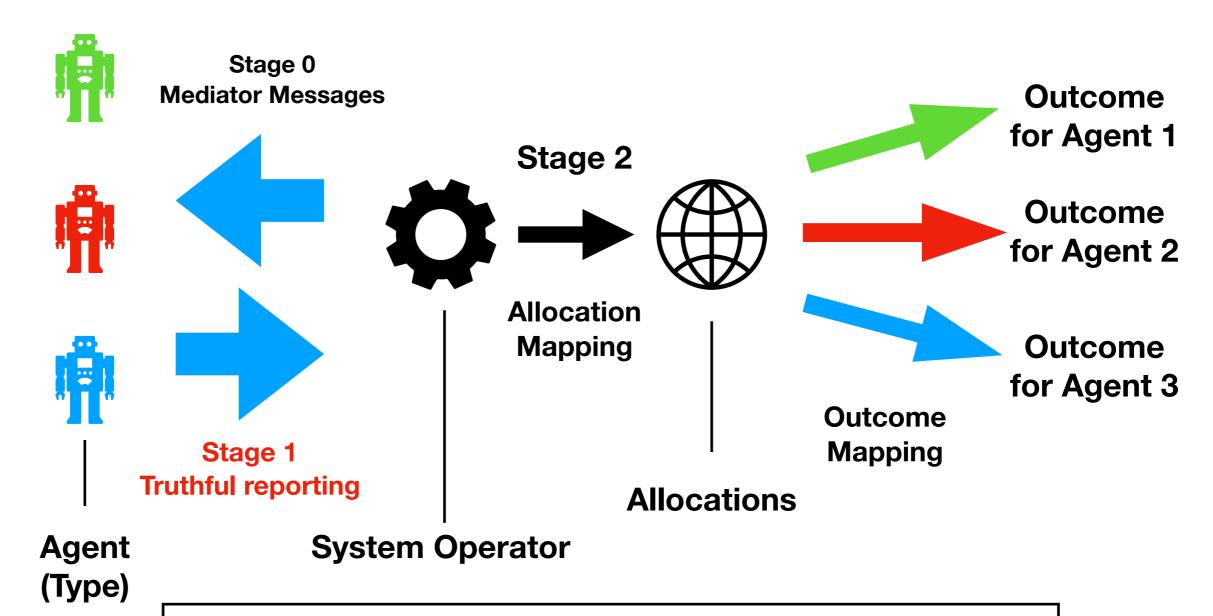


Does not hold under CPT in second-price sealed-bid auctions (Karni and Safra 1989)

Mediated Mechanism Design



Mediated Mechanism Design



Revelation Principle (under CPT)

- WLOG assume signal set = type set for each player
- restrict attention to direct truthful mediated mechanisms

Concluding remarks

- CPT provides a more general framework than EUT.
- CPT seems to more accurately model human agents...
- CPT based designs seem to have tangible benefits.
- Some structural results in EUT continue to hold under CPT in modified form (calibrated learning, mechanism design).
- CPT models provide stronger robustness guarantees relative to the classical techniques of EUT.

Thank you!



Question: Does there exist a learning strategy that does converge to C(Γ)?

Related to the notion of No regret learning

Player 1 imagines replacing his action 0 by action 1

Step	1	2	3	4	5	6
Player 1	0	0	1	1	1	0
Player 2	I	II	I	III	IV	

L1 =

Player 2	I	II	III	IV
Outcome	2β	β+1	0	1
Probability	2/3	1/3	0	0

Regret = $(1/2)[V(L_2) - V(L_1)]$

L₂ =

Player 2	I	II	III	IV
Outcome	1.99	1.99	1.99	1.99
Probability	2/3	1/3	0	0



Player i imagines replacing his action a_i by \tilde{a}_i

$$K_i^t(a_i, \tilde{a}_i) := \xi_i^t(a_i) \left[V_i \left(\left\{ (\xi_{-i}^t(a_{-i}|a_i), x_i(\tilde{a}_i, a_{-i})) \right\}_{l=1}^m \right) - V_i \left(\left\{ (\xi_{-i}^t(a_{-i}|a_i), x_i(a_i, a_{-i})) \right\}_{l=1}^m \right) \right]$$

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 Player i has no regret learning strategy if her regrets tend to be arbitrarily small almost surely, irrespective of other players' strategies.

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- Player i has no regret learning strategy if her regrets tend to be arbitrarily small almost surely, irrespective of other players' strategies.
- No regret learning is equivalent to convergence to empirical distribution

$$\limsup_{t \to \infty} K_i^t(a_i, \tilde{a}_i) \le 0 \forall a_i, \tilde{a}_i \in A_i \Leftrightarrow \xi^t \to C(\Gamma, i)$$

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Question: Does there exists a no regret learning strategy?

Answer - No! Example

Player 2

Player 1

	l	II	III	IV
0	2β,1	β+1,1	0,1	1,1
1	1.99,0	1.99,0	1.99,0	1.99,0

Action 0 is player 1's Best Response to $\begin{cases} \mu_{odd} = (0.5, 0, 0.5, 0) \\ \mu_{even} = (0, 0.5, 0, 0.5) \end{cases}$

$$egin{aligned} egin{aligned} \mu_{odd} &= (0.5, 0, 0.5, 0) \ \mu_{even} &= (0, 0.5, 0, 0.5) \end{aligned}$$

Action 1 is player 1's Best Response to $\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$

$$\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$$

Answer - No! Example

Action 0 is player 1's Best Response to $\begin{cases} \mu_{odd} = (0.5, 0, 0.5, 0) \\ \mu_{even} = (0, 0.5, 0, 0.5) \end{cases}$

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Action 1 is player 1's Best Response to $\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$

$$\mu_{unif} = (0.25, 0.25, 0.25, 0.25)$$

Strategy for player 2

- Play randomized strategy μ_{odd} at step 1,
- Play randomized strategy µ_{even} at step 2,
- Play randomized strategy μ_{odd} at step $2T^k < t \le T^{k+1}$,
- Play randomized strategy μ_{even} at step $T^{k+1} < t \le 2T^{k+1}$.

Relaxations

System Prob.

$$SYS[z, \pi; h, v, A, c]$$

Maximize
$$\sum_{i=1}^n \sum_{l=1}^k h_i(l) v_i(z_i(l))$$

Subject to
$$\sum_{i \in R_j} z_i(\pi_i(l)) \leq c_j, orall j, l$$
 $z_i(l) \geq z_i(l+1), orall i, l$

$$\pi_i \in S_k, \forall i$$

Relaxed System Prob.

$$SYS_REL[z, M; h, v, A, c]$$

Maximize
$$\sum_{i=1}^{n} \sum_{l=1}^{k} h_i(l) v_i(z_i(l))$$

Subject to
$$\sum_{i \in R_j} M_i z_i \leq c_j \mathbf{1}, orall j$$
 $z_i(l) \geq z_i(l+1), orall i, orall l$

 M_i Doubly Stoc.

Average System Prob.

$$SYS_AVG[z; h, v, A, c]$$

Maximize
$$\sum_{i=1}^n \sum_{l=1}^k h_i(l) v_i(z_i(l))$$

Subject to
$$\sum_{i \in R_j} rac{1}{k} \sum_{l=1}^k z_i(l) \le c_j, orall j$$
 $z_i(l) \ge z_i(l+1), orall i, orall l$

Theorem: For any system problem,

$$W_{ps} \leq W_{pr} = W_{pa} = W_{da} = W_{dr} = W_{ds}$$

Observations

The Relaxed system problem and the Average system problem are convex optimization problems.

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Theorem

The Primal System problem is NP Hard.

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Theorem

The Primal System problem is NP Hard.

Proof idea

Integer partition problem can be reduced to a primal system problem and hence NP hard